The Open Group Standard

Additional APIs for the Base Specifications Issue 8, Part 1



Unapproved Draft, Subject to Change

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# Preface

### The Open Group

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### This Document

This document has been prepared by The Open Group Base Working Group. The Open Group Base Working Group is considering submitting a number of additional APIs to the Austin Group as input to the Issue 8 revision of the Base Specifications.

This document contains the first set of these APIs.

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- The Open Group Base Working Group
- The Austin Group

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## 1 Introduction

### 1.1 Scope

The purpose of this document is to define a set of additional APIs for inclusion in the Issue 8 revision of the Base Specifications of the Single UNIX Specification.

The additional APIs proposed by participants in the Austin Group that The Open Group has agreed to sponsor are as follows:

dladdr() getentropy() getlocalename\_l() memmem() posix\_getdents() ppoll() pthread\_cond\_clockwait() pthread\_mutex\_clocklock() pthread\_rwlock\_clockrdlock() pthread\_rwlock\_clockwrlock() qsort\_r() reallocarray() sem\_clockwait() sig2str() str2sig() strlcat() strlcat() wcslcat() wcslcat()

### **1.2** Relationship to Other Formal Standards

This Standard is being forwarded to the Austin Group for consideration as input to the Issue 8 revision of the Base Specifications.

# 2 Application Program Interfaces

The following pages are extracted from a complete draft of the Base Specifications in which the proposed changes have been applied, with change bars showing the differences from Issue 8 draft 1.1. Only pages with technical changes are included – editorial changes such as additions to SEE ALSO and CHANGE HISTORY sections have been omitted (unless they appear on the same page as a technical change). The complete draft is also being made available for reference.

## 2.1 Change Bars

Changed lines are marked with a '|' in the right-hand margin, new lines with a '+', and deleted lines with a '-'.

Note that sometimes the placement of change bars is slightly inaccurate. In particular, changes may extend into a line following a set of change-barred lines. Also, changes within tables do not have change bars.

## 2.2 Reference Pages

The reference pages for the new functions and related header additions follow.

### 2782 4.13 Memory Synchronization

2783Applications shall ensure that access to any memory location by more than one thread of control2784(threads or processes) is restricted such that no thread of control can read or modify a memory2785location while another thread of control may be modifying it. Such access is restricted using2786functions that synchronize thread execution and also synchronize memory with respect to other2787threads. The following functions synchronize memory with respect to other threads:

2789pthread_barrier_wait()pthread_mutex_unlock()pthread_rwlock_wrlock(2790pthread_cond_broadcast()pthread_spin_lock()sem_clockwait()2791pthread_cond_clockwait()pthread_spin_trylock()sem_post()2792pthread_cond_signal()pthread_spin_unlock()sem_timedwait()2793pthread_cond_timedwait()pthread_rwlock_clockrdlock()sem_trywait()2794pthread_cond_wait()pthread_rwlock_clockwrlock()sem_wait()2795pthread_create()pthread_rwlock_rdlock()semctl()2796pthread_join()pthread_rwlock_timedrdlock()semop()2797pthread_mutex_clocklock()pthread_rwlock_timedrvlock()wait()2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()waitpid()	2788	fork()	pthread_mutex_trylock()	<pre>pthread_rwlock_unlock()</pre>
2790pthread_cond_broadcast()pthread_spin_lock()sem_clockwait()2791pthread_cond_clockwait()pthread_spin_trylock()sem_post()2792pthread_cond_signal()pthread_spin_unlock()sem_timedwait()2793pthread_cond_timedwait()pthread_rwlock_clockrdlock()sem_trywait()2794pthread_cond_wait()pthread_rwlock_clockwrlock()sem_wait()2795pthread_create()pthread_rwlock_rdlock()semctl()2796pthread_join()pthread_rwlock_timedrdlock()semop()2797pthread_mutex_clocklock()pthread_rwlock_timedwrlock()wait()2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()waitpid()	2789	pthread_barrier_wait()	pthread_mutex_unlock()	<i>pthread_rwlock_wrlock()</i>
2791pthread_cond_clockwait()pthread_spin_trylock()sem_post()2792pthread_cond_signal()pthread_spin_unlock()sem_timedwait()2793pthread_cond_timedwait()pthread_rwlock_clockrdlock()sem_trywait()2794pthread_cond_wait()pthread_rwlock_clockwrlock()sem_wait()2795pthread_create()pthread_rwlock_rdlock()semctl()2796pthread_join()pthread_rwlock_timedrdlock()semop()2797pthread_mutex_clocklock()pthread_rwlock_timedwrlock()wait()2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()waitpid()	2790	pthread_cond_broadcast()	pthread_spin_lock()	sem_clockwait()
2792pthread_cond_signal()pthread_spin_unlock()sem_timedwait()2793pthread_cond_timedwait()pthread_rwlock_clockrdlock()sem_trywait()2794pthread_cond_wait()pthread_rwlock_clockwrlock()sem_wait()2795pthread_create()pthread_rwlock_rdlock()semctl()2796pthread_join()pthread_rwlock_timedrdlock()semop()2797pthread_mutex_clocklock()pthread_rwlock_timedwrlock()wait()2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()waitpid()	2791	pthread_cond_clockwait()	pthread_spin_trylock()	sem_post()
2793pthread_cond_timedwait()pthread_rwlock_clockrdlock()sem_trywait()2794pthread_cond_wait()pthread_rwlock_clockwrlock()sem_wait()2795pthread_create()pthread_rwlock_rdlock()semctl()2796pthread_join()pthread_rwlock_timedrdlock()semop()2797pthread_mutex_clocklock()pthread_rwlock_timedwrlock()wait()2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()	2792	pthread_cond_signal()	pthread_spin_unlock()	sem_timedwait()
2794pthread_cond_wait()pthread_rwlock_clockwrlock()sem_wait()2795pthread_create()pthread_rwlock_rdlock()semctl()2796pthread_join()pthread_rwlock_timedrdlock()semop()2797pthread_mutex_clocklock()pthread_rwlock_timedwrlock()wait()2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()	2793	pthread_cond_timedwait()	pthread_rwlock_clockrdlock()	sem_trywait()
2795pthread_create()pthread_rwlock_rdlock()semctl()2796pthread_join()pthread_rwlock_timedrdlock()semop()2797pthread_mutex_clocklock()pthread_rwlock_timedwrlock()wait()2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()	2794	pthread_cond_wait()	pthread_rwlock_clockwrlock()	sem_wait()
2796pthread_join()pthread_rwlock_timedrdlock()semop()2797pthread_mutex_clocklock()pthread_rwlock_timedwrlock()wait()2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()	2795	pthread_create()	pthread_rwlock_rdlock()	semctl()
2797pthread_mutex_clocklock()pthread_rwlock_timedwrlock()wait()2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()	2796	pthread_join()	pthread_rwlock_timedrdlock()	semop()
2798pthread_mutex_lock()pthread_rwlock_tryrdlock()waitpid()2799pthread_mutex_timedlock()pthread_rwlock_trywrlock()	2797	pthread_mutex_clocklock()	pthread_rwlock_timedwrlock()	wait()
2799 pthread_mutex_timedlock() pthread_rwlock_trywrlock()	2798	<i>pthread_mutex_lock()</i>	<pre>pthread_rwlock_tryrdlock()</pre>	waitpid()
	2799	<pre>pthread_mutex_timedlock()</pre>	<pre>pthread_rwlock_trywrlock()</pre>	

- 2800The *pthread\_once(*) function shall synchronize memory for the first call in each thread for a given2801**pthread\_once\_t** object. If the *init\_routine* called by *pthread\_once(*) is a cancellation point and is2802canceled, a call to *pthread\_once(*) for the same **pthread\_once\_t** object made from a cancellation2803cleanup handler shall also synchronize memory.
- 2804The *pthread\_mutex\_lock()* function need not synchronize memory if the mutex type is2805PTHREAD\_MUTEX\_RECURSIVE and the calling thread already owns the mutex. The2806*pthread\_mutex\_unlock()* function need not synchronize memory if the mutex type is2807PTHREAD\_MUTEX\_RECURSIVE and the mutex has a lock count greater than one.
- 2808 Unless explicitly stated otherwise, if one of the above functions returns an error, it is unspecified 2809 whether the invocation causes memory to be synchronized.
- Applications may allow more than one thread of control to read a memory location simultaneously.

### 2812 4.14 Pathname Resolution

Pathname resolution is performed for a process to resolve a pathname to a particular directory 2813 entry for a file in the file hierarchy. There may be multiple pathnames that resolve to the same 2814 directory entry, and multiple directory entries for the same file. When a process resolves a 2815 pathname of an existing directory entry, the entire pathname shall be resolved as described 2816 below. When a process resolves a pathname of a directory entry that is to be created immediately 2817 after the pathname is resolved, pathname resolution terminates when all components of the path 2818 prefix of the last component have been resolved. It is then the responsibility of the process to 2819 create the final component. 2820

Each filename in the pathname is located in the directory specified by its predecessor (for example, in the pathname fragment **a/b**, file **b** is located in directory **a**). Pathname resolution shall fail if this cannot be accomplished. If the pathname begins with a <slash>, the predecessor of the first filename in the pathname shall be taken to be the root directory of the process (such pathnames are referred to as ``absolute pathnames''). If the pathname does not begin with a <slash>, the predecessor of the first filename of the pathname shall be taken to be either the current working directory of the process or for certain interfaces the directory identified by a file

## <dirent.h>

7540	NAME					
7541		dirent.h — f	ormat c	of directory ent	ries	
7542	SYNOP	SIS	<diron< th=""><th>ρ+ h∖</th><th></th><th></th></diron<>	ρ+ h∖		
7543	DECO		<ure la="" la<="" th=""><th>10.11/</th><th></th><th></th></ure>	10.11/		
7544 7545	DESCR	The internal	format	of directories i	s unspecified.	
7546		The <b><dirent< b=""></dirent<></b>	<b>.h&gt;</b> hea	der shall define	e the following type:	
7547		DIR At	ype rep	resenting a dir	ectory stream. The <b>DIR</b> type may be an incomplete type.	
7548		It shall also	define t	he structure <b>di</b>	rent which shall include the following members:	
7549 7550		ino_t d_ char d_	ino name[]	File seria Filename	al number. e string of entry.	-
7551		and the stru	cture <b>p</b>	osix_dent whic	h shall include the following members:	
7552 7553 7554 7555 7556		<pre>ino_t reclen_t unsigned char</pre>	char	d_ino d_reclen d_type d_name[]	File serial number. Length of this entry, including trailing padding if necessary. See <i>posix_getdents()</i> . File type or unknown-file-type indication. Filename string of this entry.	
7557 7558		The array <i>d_</i> at most {NA	_ <i>name</i> ir ME_M	n each of these AX} bytes follo	structures is of unspecified size, but shall contain a filename of wed by a terminating null byte.	
7559 7560		The <b><dirent< b=""> <b><sys b="" types.<=""></sys></b></dirent<></b>	. <b>h&gt;</b> hea 1>.	der shall defin	e the <b>ino_t</b> , <b>reclen_t</b> , <b>size_t</b> , and <b>ssize_t</b> types as described in	+ +
7561 7562 7563		The <b><dirent< b=""> unknown-fil values shall</dirent<></b>	<b>t.h&gt;</b> he le-type be disti	ader shall def indicator retur nct and shall b	ine the following symbolic constants for the file types and rned in the <i>d_type</i> member of the <b>posix_dent</b> structure. The e suitable for use in <b>#if</b> preprocessing directives:	+ + +
7564		DT_BLK	Block	special.		+
7565		DT_CHR	Chara	cter special.		+
7566		DT_DIR	Direct	ory.		+
7567		DT_FIFO	FIFO s	special.		+
7568		DT_LNK	Symb	olic link.		+
7569		DT_REG	Regul	ar.		+
7570		DT_SOCK	Socke	t.		+
7571 7572		DT_UNKN(	OWN Unkne	own file type.		+ +
7573 7574 7575 7576 7577 7578	ТҮМ	The implement typed memory these types. beginning we corresponding The values s	nentation ory obje The val with DT ng cons shall be	n may implem cts as distinct lues shall be di '_, except whe stant shall have suitable for use	ent message queues, semaphores, shared memory objects or file types. The following macros shall be provided to represent istinct from each other and from the above symbolic constants in a distinct file type is not implemented, in which case the e a value that is never returned in <i>d_type</i> by <i>posix_getdents()</i> . e in <b>#if</b> preprocessing directives:	+ + + + + + + +
7579		DI_MQ	Messa	ige queue.		+

7580		DT_SEM S	emaphore.	+
7581		DT_SHM S	hared memory object.	+
7582	TYM	DT_TMO T	yped memory object.	+
7583 7584		The following prototypes sha	shall be declared as functions and may also be defined as macros. Function ll be provided.	
7585		int	<pre>alphasort(const struct dirent **, const struct dirent **)</pre>	;
7586		int	closedir(DIR *);	
7587		int	<pre>dirfd(DIR *);</pre>	
7588		DIR	<pre>*fdopendir(int);</pre>	
7589		DIR	<pre>*opendir(const char *);</pre>	
7590		ssize_t	<pre>posix_getdents(int, void *, size_t, int);</pre>	+
7591		struct dire	nt *readdir(DIR *);	
7592		int	<pre>readdir_r(DIR *restrict, struct dirent *restrict,</pre>	
7593			<pre>struct dirent **restrict);</pre>	
7594		void	<pre>rewinddir(DIR *);</pre>	
7595		int	<pre>scandir(const char *, struct dirent ***,</pre>	
7596			<pre>int (*)(const struct dirent *),</pre>	
7597			int (*)(const struct dirent **,	
7598			<pre>const struct dirent **));</pre>	
7599	XSI	void	<pre>seekdir(DIR *, long);</pre>	
7600		long	<pre>telldir(DIR *);</pre>	

#### 7601 APPLICATION USAGE

None.

7602

### 7603 RATIONALE

7604Information similar to that in the <dirent.h> header is contained in a file <sys/dir.h> in 4.2 BSD7605and 4.3 BSD. The equivalent in these implementations of struct dirent from this volume of7606POSIX.1-202x is struct direct. The filename was changed because the name <sys/dir.h> was also7607used in earlier implementations to refer to definitions related to the older access method; this7608produced name conflicts. The name of the structure was changed because this volume of7609POSIX.1-202x does not completely define what is in the structure, so it could be different on7610some implementations from struct direct.

7611 The **posix\_dent** structure was based on existing structures used by traditional getdents() functions, but the name was changed because the existing structures differed in name and in + 7612 their members. Some used the **dirent** structure but this is not required to include a  $d_{type}$ 7613 member, which is the main advantage of using *posix\_getdents()* over *readdir()*. The *d\_reclen* 7614 + member was included, even though some implementations return fixed-length entries and 7615 + therefore do not need it, as almost all existing code that used *getdents()* used  $d_{reclen}$  to iterate 7616 + through the returned entries. Implementations that return fixed-length entries can simply set + 7617 *d\_reclen* to that length in *posix\_getdents*(). The type **reclen\_t** for *d\_reclen* was introduced, instead 7618 + of using **unsigned short**, so as not to create a requirement that {NAME\_MAX} cannot be greater + 7619 than (a value somewhat smaller than) {SHRT\_MAX}. 7620 +

- 7621Implementations are encouraged to define a DT\_FORCE\_TYPE symbolic constant for use in the+7622flags argument to posix\_getdents().See the RATIONALE for posix\_getdents().+
- The name of an array of **char** of an unspecified size should not be used as an lvalue. Use of:
- 7624 sizeof(d\_name)

7625		is incorrect; use:	
7626		<pre>strlen(d_name)</pre>	
7627		instead.	
7628 7629 7630 7631		The array of <b>char</b> <i>d_name</i> cannot be assumed to have a fixed size. Implementations may define the <i>d_name</i> array in the <b>dirent</b> and <b>posix_dent</b> structures to have size 1, or size greater than {NAME_MAX}, or use a flexible array member, but in all cases the actual number of characters used for <i>d_name</i> is at least the length of the filename string including the terminating NUL byte.	
7632 7633 7634	FUTUR	<b>E DIRECTIONS</b> A future version of this standard may add a DT_FORCE_TYPE symbolic constant for use as described in the RATIONALE for <i>posix_getdents()</i> .	
7635 7636	SEE AL	SO <sys types.h=""></sys>	
7637 7638		XSH alphasort(), closedir(), dirfd(), fdopendir(), posix_getdents(), readdir(), rewinddir(), seekdir(), telldir()	+
7639 7640	CHANG	GE HISTORY First released in Issue 2.	
7641 7642	Issue 5	The DESCRIPTION is updated for alignment with the POSIX Threads Extension.	
7643 7644	Issue 6	The Open Group Corrigendum U026/7 is applied, correcting the prototype for <i>readdir_r(</i> ).	
7645		The <b>restrict</b> keyword is added to the prototype for <i>readdir_r(</i> ).	
7646 7647 7648	Issue 7	The <i>alphasort(), dirfd(),</i> and <i>scandir()</i> functions are added from The Open Group Technical Standard, 2006, Extended API Set Part 1.	
7649 7650		The <i>fdopendir</i> () function is added from The Open Group Technical Standard, 2006, Extended API Set Part 2.	
7651 7652		Austin Group Interpretation 1003.1-2001 #110 is applied, clarifying the definition of the <b>DIR</b> type.	
7653 7654		POSIX.1-2008, Technical Corrigendum 1, XBD/TC1-2008/0039 [291], XBD/TC1-2008/0040 [291], XBD/TC1-2008/0041 [291], and XBD/TC1-2008/0042 [206] are applied.	+
7655 7656	Issue 8	Austin Group Defect 697 is applied, adding <i>posix_getdents()</i> .	+

7657	NAME		
7658		dlfcn.h — dynamic linking	
7659	SYNOP	SIS	
7660		<pre>#include <dlfcn.h></dlfcn.h></pre>	
7661 7662 7663	DESCR	<b>IPTION</b> The <b><dlfcn.h></dlfcn.h></b> header shall define the <b>Dl_info_t</b> structure type, which shall include at least the following members:	+ +
7664 7665 7666 7667		const char *dli_fnamePathname of mapped object file.void*dli_fbaseBase of mapped address range.const char *dli_snameSymbol name or null pointer.void*dli_saddrSymbol address or null pointer.	+ + +
7668 7669		The <b><dlfcn.h></dlfcn.h></b> header shall define at least the following symbolic constants for use in the construction of a <i>dlopen() mode</i> argument:	
7670		RTLD_LAZY Relocations are performed at an implementation-defined time.	
7671		RTLD_NOW Relocations are performed when the object is loaded.	
7672		RTLD_GLOBAL All symbols are available for relocation processing of other modules.	
7673 7674		RTLD_LOCAL All symbols are not made available for relocation processing by other modules.	
7675 7676		The following shall be declared as functions and may also be defined as macros. Function prototypes shall be provided.	
7677 7678 7679 7680 7681		<pre>int dladdr(const void *restrict, Dl_info_t *restrict); int dlclose(void *); char *dlerror(void); void *dlopen(const char *, int); void *dlsym(void *restrict, const char *restrict);</pre>	+
7682 7683	APPLIC	CATION USAGE	
7684 7685	RATIO	NALE None.	
7686 7687	FUTUR	E DIRECTIONS None.	
7688 7689	SEE AL	SO XSH dladdr(), dlclose(), dlerror(), dlopen(), dlsym()	+
7690 7691	CHANG	GE HISTORY First released in Issue 5.	
7692 7693	Issue 6	The <b>restrict</b> keyword is added to the prototype for <i>dlsym</i> ().	
7694 7695	Issue 7	The <b><dlfcn.h></dlfcn.h></b> header is moved from the XSI option to the Base.	
7696		This reference page is clarified with respect to macros and symbolic constants.	+

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### Headers

+ +

9402	Other Invariant Values
9403	The <b><limits.h></limits.h></b> header shall define the following symbolic constants:
9404	{GETENTROPY_MAX}
9405	The maximum value of the <i>length</i> argument in calls to the <i>getentropy()</i> function.
9406	Minimum Acceptable Value: 256
9407	<pre>{NL_ARGMAX}</pre>
9408	Maximum value of n in conversion specifications using the "%n\$" sequence in calls to the
9409	printf() and scanf() families of functions.
9410	Minimum Acceptable Value: 9
9411	{NL_LANGMAX}
9412	Maximum number of bytes in a <i>LANG</i> name.
9413	Minimum Acceptable Value: 14
9414	{NL_MSGMAX}
9415	Maximum message number.
9416	Minimum Acceptable Value: 32 767
9417	{NL_SETMAX}
9418	Maximum set number.
9419	Minimum Acceptable Value: 255
9420	{NL_TEXTMAX}
9421	Maximum number of bytes in a message string.
9422	Minimum Acceptable Value: {_POSIX2_LINE_MAX}
9423	<pre>{NSIG_MAX}</pre>
9424	Maximum possible return value of sysconf(_SC_NSIG). See XSH sysconf(). The value of
9425	{NSIG_MAX} shall be no greater than the number of signals that the sigset_t type (see
9426	<signal.h>) is capable of representing, ignoring any restrictions imposed by sigfillset() or</signal.h>
9427	sigaddset().
9428	{NZERO}
9429	Default process priority.
9430	Minimum Acceptable Value: 20
9431 9432	PPLICATION USAGE
9433 9434 9435 9436	A request was made to reduce the value of {_POSIX_LINK_MAX} from the value of 8 specified for it in the POSIX.1-1990 standard to 2. The standard developers decided to deny this request for several reasons:
9437 9438	<ul> <li>They wanted to avoid making any changes to the standard that could break conforming applications, and the requested change could have that effect.</li> </ul>
9439 9440 9441 9442	• The use of multiple hard links to a file cannot always be replaced with use of symbolic links. Symbolic links are semantically different from hard links in that they associate a pathname with another pathname rather than a pathname with a file. This has implications for access control, file permanence, and transparency.
9443	<ul> <li>The original standard developers had considered the issue of allowing for</li></ul>
9444	implementations that did not in general support hard links, and decided that this would
9445	reduce consensus on the standard.

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9612 9613	СХ	The <b><locale.h></locale.h></b> header shall contain at least the following macros representing bitmasks for use with the <i>newlocale()</i> function for each supported locale category:	
9614 9615		LC_COLLATE_MASK LC_CTYPE_MASK	
9616		LC MESSAGES MASK	
9617		LC MONETARY MASK	
9618		LC NUMERIC MASK	
9619		LC_TIME_MASK	
9620		In addition, a macro to set the bits for all categories set shall be defined:	
9621		LC_ALL_MASK	
9622 9623		The <b><locale.h></locale.h></b> header shall define LC_GLOBAL_LOCALE, a special locale object descriptor used by the <i>duplocale()</i> and <i>uselocale()</i> functions.	
9624		The <b><locale.h></locale.h></b> header shall define the <b>locale_t</b> type, representing a locale object.	
9625 9626		The following shall be declared as functions and may also be defined as macros. Function prototypes shall be provided for use with ISO C standard compilers.	
9627	CX	<pre>locale_t duplocale(locale_t);</pre>	
9628		<pre>void freelocale(locale_t);</pre>	
9629		<pre>const char *getlocalename_l(int, locale_t);</pre>	+
9630		<pre>struct lconv *localeconv(void);</pre>	
9631	CX	<pre>locale_t newlocale(int, const char *, locale_t);</pre>	
9632		<pre>char *setlocale(int, const char *);</pre>	
9633	CX	locale_t uselocale (locale_t);	
9634	APPLIC	ATION USAGE	
9635		None.	
9636	RATIO	NALE	
9637		It is suggested that each category macro name for use in <i>setlocale()</i> have a corresponding macro	
9638		name ending in _MASK for use in newlocale().	
9639	FUTUR	E DIRECTIONS	
9640		None.	
9641	SEE AL	50	
9642	OLL IIL	Chapter 8 (on page 153), <stddef.h></stddef.h>	
9643		XSH duplocale(), freelocale(), getlocalename_l(), localeconv(), newlocale(), setlocale(), uselocale()	+
9644	CHANC	GE HISTORY	
9645		First released in Issue 3.	
9040	T (	included for angument with the 150 C standard.	
9647	Issue 6	The large structure is some ded with a second of the large state of th	
9648		The iconv structure is expanded with new members (int_n_cs_precedes, int_n_sep_by_space,	
9649 9650		with the ISO/IEC 9899: 1999 standard.	
9651		Extensions beyond the ISO C standard are marked.	

<poll.h>

10565	NAME				
10566		poll.h —	- definitions	for the poll() function	
10567	SYNOP	SIS	do cooll i		
10568	DECO		de (poir.		
10569 10570 10571	DESCK	The <b><po< b=""> member</po<></b>	<b>ll.h&gt;</b> header s:	shall define the <b>pollfd</b> structure, which shall include at least the following	
10572 10573 10574		int short short	fd events revents	The following descriptor being polled. The input event flags (see below). The output event flags (see below).	
10575		The <b><po< b=""></po<></b>	ll.h> header	shall define the following type through <b>typedef</b> :	
10576		nfds_t		An unsigned integer type used for the number of file descriptors.	
10577 10578 10579		The imp of <b>nfds</b> environ	lementation _t is no gre nents can be	shall support one or more programming environments in which the width eater than the width of type <b>long</b> . The names of these programming obtained using the <i>confstr()</i> function or the <i>getconf</i> utility.	
10580		The <b><po< b=""></po<></b>	ll.h> header	shall define the <b>sigset_t</b> type as described in <b><signal.h></signal.h></b> .	+
10581		The <b><po< b=""></po<></b>	ll.h> header	shall define the <b>timespec</b> structure as described in <b><time.h></time.h></b> .	+
10582 10583		The <b><po< b=""> be OR'ed</po<></b>	oll.h> header d together to	shall define the following symbolic constants, zero or more of which may form the <i>events</i> or <i>revents</i> members in the <b>pollfd</b> structure:	
10584		POLLIN	ſ	Data other than high-priority data may be read without blocking.	
10585		POLLRI	DNORM	Normal data may be read without blocking.	
10586		POLLRI	DBAND	Priority data may be read without blocking.	
10587		POLLPR	RI	High priority data may be read without blocking.	
10588		POLLO	ĴΤ	Normal data may be written without blocking.	
10589		POLLW	RNORM	Equivalent to POLLOUT.	
10590		POLLW	RBAND	Priority data may be written.	
10591		POLLER	RR	An error has occurred ( <i>revents</i> only).	
10592		POLLH	UP	Device has been disconnected ( <i>revents</i> only).	
10593		POLLN	VAL	Invalid <i>fd</i> member ( <i>revents</i> only).	
10594 10595		The sigr specific.	nificance and	l semantics of normal, priority, and high-priority data are file and device-	
10596 10597		The following shall prototypes shall be p		be declared as functions and may also be defined as macros. Function rovided.	
10598 10599 10600		int int	poll(stru ppoll(str const	ct pollfd [], nfds_t, int); uct pollfd [], nfds_t, const struct timespec *restrict, sigset_t *restrict);	+ +
10601 10602		Inclusion	n of the <b><po< b="">l &gt;.</po<></b>	ll.h> header may make visible all symbols from the headers <signal.h> and</signal.h>	+

# <pthread.h>

### Headers

10671		int	<pre>pthread_atfork(void (*)(void), void (*)(void),</pre>	
10672		4 m F	Vold(^)(Vold));	
10673		int	<pre>pthread_attr_destroy(pthread_attr_t ^); </pre>	
10674		int	pthread_attr_getdetachstate(const pthread_attr_t ^, int ^);	
10675		TUC	punread_attr_getguardsize(const punread_attr_t ^restrict,	
10676	TDC	2	size_t ^restrict);	
10677	TPS	int	<pre>ptnread_attr_getinneritsched(const ptnread_attr_t *restrict,</pre>	
10678		1	<pre>int *restrict);</pre>	
10679		int	pthread_attr_getschedparam(const pthread_attr_t ^restrict,	
10680			<pre>struct sched_param *restrict);</pre>	
10681	TPS	int	<pre>pthread_attr_getschedpolicy(const pthread_attr_t *restrict,</pre>	
10682			<pre>int *restrict);</pre>	
10683		int	<pre>pthread_attr_getscope(const pthread_attr_t *restrict,</pre>	
10684			<pre>int *restrict);</pre>	
10685	TSA TSS	int	<pre>pthread_attr_getstack(const pthread_attr_t *restrict,</pre>	
10686			<pre>void **restrict, size_t *restrict);</pre>	
10687	TSS	int	<pre>pthread_attr_getstacksize(const pthread_attr_t *restrict,</pre>	
10688			<pre>size_t *restrict);</pre>	
10689		int	<pre>pthread_attr_init(pthread_attr_t *);</pre>	
10690		int	<pre>pthread_attr_setdetachstate(pthread_attr_t *, int);</pre>	
10691		int	<pre>pthread_attr_setguardsize(pthread_attr_t *, size_t);</pre>	
10692	TPS	int	<pre>pthread_attr_setinheritsched(pthread_attr_t *, int);</pre>	
10693		int	<pre>pthread_attr_setschedparam(pthread_attr_t *restrict,</pre>	
10694			<pre>const struct sched_param *restrict);</pre>	
10695	TPS	int	<pre>pthread_attr_setschedpolicy(pthread_attr_t *, int);</pre>	
10696		int	<pre>pthread_attr_setscope(pthread_attr_t *, int);</pre>	
10697	TSA TSS	int	<pre>pthread_attr_setstack(pthread_attr_t *, void *, size_t);</pre>	
10698	TSS	int	<pre>pthread_attr_setstacksize(pthread_attr_t *, size_t);</pre>	
10699		int	<pre>pthread_barrier_destroy(pthread_barrier_t *);</pre>	
10700		int	<pre>pthread_barrier_init(pthread_barrier_t *restrict,</pre>	
10701			const pthread_barrierattr_t *restrict, unsigned);	
10702		int	pthread_barrier_wait(pthread_barrier_t *);	
10703		int	<pre>pthread_barrierattr_destroy(pthread_barrierattr_t *);</pre>	
10704	TSH	int	pthread_barrierattr_getpshared(	
10705			<pre>const pthread_barrierattr_t *restrict, int *restrict);</pre>	
10706		int	<pre>pthread_barrierattr_init(pthread_barrierattr_t *);</pre>	
10707	TSH	int	<pre>pthread_barrierattr_setpshared(pthread_barrierattr_t *, int);</pre>	
10708		int	<pre>pthread_cancel(pthread_t);</pre>	
10709		int	<pre>pthread_cond_broadcast(pthread_cond_t *);</pre>	
10710		int	<pre>pthread_cond_clockwait(pthread_cond_t *restrict,</pre>	+
10711			<pre>pthread_mutex_t *restrict, clockid_t,</pre>	+
10712			<pre>const struct timespec *restrict);</pre>	+
10713		int	pthread_cond_destroy(pthread_cond_t *);	
10714		int	<pre>pthread_cond_init(pthread_cond_t *restrict,</pre>	
10715			<pre>const pthread_condattr_t *restrict);</pre>	
10716		int	<pre>pthread_cond_signal(pthread_cond_t *);</pre>	
10717		int	<pre>pthread_cond_timedwait(pthread_cond_t *restrict,</pre>	
10718			<pre>pthread_mutex_t *restrict, const struct timespec *restrict);</pre>	
10719		int	<pre>pthread_cond_wait(pthread_cond_t *restrict,</pre>	
10720			<pre>pthread_mutex_t *restrict);</pre>	
10721		int	<pre>pthread_condattr_destroy(pthread_condattr_t *);</pre>	
10722		int	<pre>pthread_condattr_getclock(const pthread_condattr_t *restrict,</pre>	

10723			clockid t *restrict):	
10724	TSH	int	pthread condattr getpshared(const pthread condattr t *restrict.	
10725	1011	1110	int *restrict):	
10726		int	pthread condattr init(pthread condattr t *):	
10727		int	pthread condattr setclock(pthread condattr t *. clockid t):	
10728	TSH	int.	pthread condattr setpshared(pthread condattr t *, int);	
10729		int.	pthread create (pthread t *restrict, const pthread attr t *restrict)	
10730			void *(*) (void*), void *restrict);	,
10731		int	pthread detach(pthread t);	
10732		int	pthread equal (pthread t, pthread t);	
10733		void	pthread exit (void *);	
10734	TCT	int	pthread getcpuclockid(pthread t, clockid t *);	
10735	TPS	int	pthread_getschedparam(pthread_t, int *restrict,	
10736			struct sched_param *restrict);	
10737		void	<pre>*pthread_getspecific(pthread_key_t);</pre>	
10738		int	pthread_join(pthread_t, void **);	
10739		int	<pre>pthread_key_create(pthread_key_t *, void (*)(void*));</pre>	
10740		int	<pre>pthread_key_delete(pthread_key_t);</pre>	
10741		int	pthread_mutex_clocklock(pthread_mutex_t *restrict, clockid_t,	+
10742			const struct timespec *restrict);	+
10743		int	<pre>pthread_mutex_consistent(pthread_mutex_t *);</pre>	
10744		int	<pre>pthread_mutex_destroy(pthread_mutex_t *);</pre>	
10745	RPP   TPP	int	<pre>pthread_mutex_getprioceiling(const pthread_mutex_t *restrict,</pre>	
10746			<pre>int *restrict);</pre>	
10747		int	<pre>pthread_mutex_init(pthread_mutex_t *restrict,</pre>	
10748			<pre>const pthread_mutexattr_t *restrict);</pre>	
10749		int	<pre>pthread_mutex_lock(pthread_mutex_t *);</pre>	
10750	RPP   TPP	int	<pre>pthread_mutex_setprioceiling(pthread_mutex_t *restrict, int,</pre>	
10751			<pre>int *restrict);</pre>	
10752		int	<pre>pthread_mutex_timedlock(pthread_mutex_t *restrict,</pre>	
10753			<pre>const struct timespec *restrict);</pre>	
10754		int	<pre>pthread_mutex_trylock(pthread_mutex_t *);</pre>	
10755		int	pthread_mutex_unlock(pthread_mutex_t *);	
10756		int	<pre>pthread_mutexattr_destroy(pthread_mutexattr_t *);</pre>	
10757	RPP   TPP	int	<pre>pthread_mutexattr_getprioceiling(</pre>	
10758			<pre>const pthread_mutexattr_t *restrict, int *restrict);</pre>	
10759	MC1	int	<pre>pthread_mutexattr_getprotocol(const pthread_mutexattr_t *restrict,</pre>	
10760			<pre>int *restrict);</pre>	
10761	TSH	int	<pre>pthread_mutexattr_getpshared(const pthread_mutexattr_t *restrict,</pre>	
10762		1 to	<pre>int *restrict);</pre>	
10763		int	<pre>ptnread_mutexattr_getrobust(const ptnread_mutexattr_t *restrict, int tuestmist);</pre>	
10764		int	int ^restrict);	
10765		IIIC	<pre>int_treat_int);</pre>	
10766		int	The "resurred);	
10767	ррр   тор	int	nthread mutevattr setprioceiling/nthread mutevattr t * int).	
10760	MC1	int	pthread mutexattr setprotocol (nthread mutexattr t * int);	
10770	TSH	int	pthread mutexattr setpshared (nthread mutexattr + * int).	
10771	1011	int	pthread mutexattr setrobust (pthread mutexattr t * int).	
10772		int	pthread mutexattr settype(pthread mutexattr t *, int):	
10773		int	pthread once (pthread once t *, void (*) (void)):	
10774		int	<pre>pthread_rwlock_destroy(pthread_rwlock t *);</pre>	

10775		int	<pre>pthread_rwlock_init(pthread_rwlock_t *restrict,</pre>	
10776			<pre>const pthread_rwlockattr_t *restrict);</pre>	
10777		int	<pre>pthread_rwlock_clockrdlock(pthread_rwlock_t *restrict,</pre>	+
10778			<pre>clockid_t, const struct timespec *restrict);</pre>	+
10779		int	<pre>pthread_rwlock_clockwrlock(pthread_rwlock_t *restrict,</pre>	+
10780			<pre>clockid_t, const struct timespec *restrict);</pre>	+
10781		int	<pre>pthread_rwlock_rdlock(pthread_rwlock_t *);</pre>	
10782		int	<pre>pthread_rwlock_timedrdlock(pthread_rwlock_t *restrict,</pre>	
10783			<pre>const struct timespec *restrict);</pre>	
10784		int	<pre>pthread_rwlock_timedwrlock(pthread_rwlock_t *restrict,</pre>	
10785			<pre>const struct timespec *restrict);</pre>	
10786		int	<pre>pthread_rwlock_tryrdlock(pthread_rwlock_t *);</pre>	
10787		int	<pre>pthread_rwlock_trywrlock(pthread_rwlock_t *);</pre>	
10788		int	<pre>pthread_rwlock_unlock(pthread_rwlock_t *);</pre>	
10789		int	<pre>pthread_rwlock_wrlock(pthread_rwlock_t *);</pre>	
10790		int	<pre>pthread_rwlockattr_destroy(pthread_rwlockattr_t *);</pre>	
10791	TSH	int	pthread_rwlockattr_getpshared(	
10792			<pre>const pthread_rwlockattr_t *restrict, int *restrict);</pre>	
10793		int	<pre>pthread_rwlockattr_init(pthread_rwlockattr_t *);</pre>	
10794	TSH	int	<pre>pthread_rwlockattr_setpshared(pthread_rwlockattr_t *, int);</pre>	
10795		pthre	ad_t	
10796			<pre>pthread_self(void);</pre>	
10797		int	<pre>pthread_setcancelstate(int, int *);</pre>	
10798		int	<pre>pthread_setcanceltype(int, int *);</pre>	
10799	TPS	int	<pre>pthread_setschedparam(pthread_t, int,</pre>	
10800			<pre>const struct sched_param *);</pre>	
10801		int	<pre>pthread_setschedprio(pthread_t, int);</pre>	
10802		int	<pre>pthread_setspecific(pthread_key_t, const void *);</pre>	
10803		int	<pre>pthread_spin_destroy(pthread_spinlock_t *);</pre>	
10804		int	<pre>pthread_spin_init(pthread_spinlock_t *, int);</pre>	
10805		int	<pre>pthread_spin_lock(pthread_spinlock_t *);</pre>	
10806		int	<pre>pthread_spin_trylock(pthread_spinlock_t *);</pre>	
10807		int	<pre>pthread_spin_unlock(pthread_spinlock_t *);</pre>	
10808		void	<pre>pthread_testcancel(void);</pre>	
10809		The fo	llowing may be declared as functions, or defined as macros, or both. If functions are	
10810		declare	ed, function prototypes shall be provided.	
10811		р	thread_cleanup_pop()	
10812		р	thread_cleanup_push()	

10813Inclusion of the **<pthread.h>** header shall make symbols defined in the headers **<sched.h>** and10814**<time.h>** visible.

11165	NAME						
11166	semaphore.n — semaphores						
11167	SYNOPSIS	YNOPSIS					
11168	<pre>#include <semaphore.h></semaphore.h></pre>	<pre>#include <semaphore.h></semaphore.h></pre>					
11169	DESCRIPTION						
11170	The <b><semaphore.h></semaphore.h></b> header shall define the <b>sem t</b> type, used in performing semaphore						
11171	operations. The semaphore may be implemented using a file descriptor, in which case						
11172	applications are able to open up at least a total of {OPEN MAX} files and semaphores.	applications are able to open up at least a total of {OPEN_MAX} files and semaphores.					
11173	The <b><semaphore.h></semaphore.h></b> header shall define the <b>timespec</b> structure as described in <b><time.h></time.h></b> .						
11174	The <b><semaphore.h></semaphore.h></b> header shall define the symbolic constant SEM_FAILED which shall have						
11175	type sem t*.						
11176	The <b><semaphore.h></semaphore.h></b> header shall define O_CREAT and O_EXCL as described in <b><fcnt1.h></fcnt1.h></b> .						
11177	The following shall be declared as functions and may also be defined as macros. Function						
11178	prototypes shall be provided.						
11170	int sem clockwait (sem t *restrict clockid t	+					
11180	const struct timespec *restrict).	+					
11181	int sem close (sem t *):						
11182	int sem destroy(sem t *).						
11102	int sem_descroy(sem_t );						
11184	int sem init (sem t * int unsigned):						
11185	sem t *sem open(const char * int ):						
11186	int sem post (sem t *):						
11187	int sem timedwait (sem t *restrict const struct timespec *restrict):						
11188	int sem trywait (sem t *):						
11189	int sem unlink (const char *):						
11190	int sem wait(sem t *);						
11191	Inclusion of the <b><semaphore.b></semaphore.b></b> header may make visible symbols defined in the <b><fcntl.b></fcntl.b></b> and						
11192	<time.h> headers.</time.h>						
11103	APPLICATION USAGE						
11194	None.						
11195	RATIONALE						
11196	None.						
11197	<b>FUTURE DIRECTIONS</b>						
11198	None.						
	SEE ALCO						
11199	SEE ALSO						
11200	<re>creations, <sys types.n="">, <time.n></time.n></sys></re>						
11201 11202	XSH sem_close(), sem_destroy(), sem_getvalue(), sem_init(), sem_open(), sem_post(), sem_timedwait(), sem_trywait(), sem_unlink()						
11000	CHANCE HISTORY						
11203	First released in Issue 5. Included for alignment with the POSIX Realtime Extension						
11204	This released in issue 5. included for angriment with the root Reathing Extension.						
11205	Issue 6						
11206	The <b><semaphore.h></semaphore.h></b> header is marked as part of the Semaphores option.						
11207	The Open Group Corrigendum U021/3 is applied, adding a description of SEM FAILED.						
	The same time drugit() for a time is a dial for all a new set with TEEE (14,1000,14,1000)						
11208	The sem_timeawait() function is added for alignment with IEEE Sta 1003.1a-1999.						
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11291	The <b>sigval</b> union shall be defined as:					
11292	int sival_int	Integer signal value.				
11293	void *sival_ptr	Pointer signal value.				
11294	The <b><signal.h></signal.h></b> header	shall declare the SIGRTMIN and SIGRTMAX macros, which shall expand				
11295	to positive integer exp	essions with type <b>int</b> , but which need not be constant expressions. These				
11296	macros specify a range	of signal numbers that are reserved for application use and for which the				
11297	realtime signal behavi	or specified in this volume of POSIX.1-202x is supported. The signal				
11298	numbers in this range of	lo not overlap any of the signals specified in the following table.				
11299 11300	The range SIGRTMIN through SIGRTMAX inclusive shall include at least {RTSIG_MAX} signal numbers. The value of SIGRTMAX shall be less than the value returned by <i>sysconf</i> (SC NSIG).					
11301	It is implementation-de	It is implementation-defined whether realtime signal behavior is supported for other signals.				
11302	The <b><signal.h></signal.h></b> header	shall define the following symbolic constant. The value shall be suitable	+			
11303	for use in <b>#if</b> preprocessing directives:					
11304	SIG2STR_MAX	Maximum size of a signal name returned by sig2str(), including the	4			
11305		erminating null byte.	+			

11306The **<signal.h>** header shall define the following macros that are used to refer to the signals that11307occur in the system. Signals defined here begin with the letters SIG followed by an uppercase11308letter. The macros shall expand to positive integer constant expressions with type **int** and11309CX11310reserved for use as the null signal (see *kill()*). Additional implementation-defined signals may11311occur in the system.

11312The ISO C standard only requires the signal names SIGABRT, SIGFPE, SIGILL, SIGINT,11313SIGSEGV, and SIGTERM to be defined. An implementation need not generate any of these six11314CX11315signals, except as a result of explicit use of interfaces that generate signals, such as *raise(), kill()*,11315the General Terminal Interface (see Section 11.1.9, on page 185), and the *kill* utility, unless11316otherwise stated (see, for example, XSH Section 2.8.3.3, on page 491).

11317The following signals shall be supported on all implementations (default actions are explained11318below the table):

+

+

# <signal.h>

11401	CA	in duard	on, the following sign	ar speenle mornation stan be avalable.	
11462		Signa	l Member	Value	
11463		SIGILL	<b>void</b> * <i>si_addr</i>	Address of faulting instruction.	
11464		SIGFPE	2		
11465		SIGSEC	GV <b>void</b> * <i>si_addr</i>	Address of faulting memory reference.	
11466		SIGBUS	5		
11467		SIGCH	LD <b>pid_t</b> si_pid	Child process ID.	
11468			int si_status	If <i>si_code</i> is equal to CLD_EXITED, then <i>si_status</i> holds the exit	
11469				value of the process; otherwise, it is equal to the signal that	
11470				caused the process to change state. The exit value in <i>si_status</i>	
11471				shall be equal to the full exit value (that is, the value passed to	
11472				_exit(), _Exit(), or exit(), or returned from main()); it shall not	
11473				be limited to the least significant eight bits of the value.	
11474			uid_t si_uid	Real user ID of the process that sent the signal.	
11475		For some	e implementations, the	e value of <i>si_addr</i> may be inaccurate.	
11476		The follo	owing shall be decla	red as functions and may also be defined as macros. Function	
11477		prototyp	es shall be provided.		
11478	CX	int	kill(pid_t, int)	);	
11479	XSI	int	killpg(pid_t, in	nt);	
11480	CX	void	<pre>psiginfo(const siginfo_t *, const char *);</pre>		
11481		void	<pre>psignal(int, const char *);</pre>		
11482		int	<pre>pthread_kill(pthread_t, int);</pre>		
11483		int	<pre>pthread_sigmask(int, const sigset_t *restrict,</pre>		
11484			<pre>sigset_t *restrict);</pre>		
11485		int	<pre>raise(int);</pre>		
11486	CX	int	sig2str(int, cha	ar *);	
11487		int	sigaction(int,	const struct sigaction *restrict,	
11488		1 m 4	struct siga	ction *restrict);	
11489	VCI	int	sigadoset (sigse	L_t ^, INC);	
11490	XSI	int	sigaltstack (cons	<pre>st stack_t ^restrict, stack_t ^restrict); </pre>	
11491	CX	int	sigemptyset (sige	$L_t = 1$	
11492		int	signillent (sign	$s = c_{1}$	
11495		int	sigismember(con	st signed t * int).	
11495		void (	*signal(int, voi	d (*) (int))) (int):	
11496	СХ	int.	signer(inc, ici	et t *):	
11497		int	sigprocmask(int	, const sigset t *restrict, sigset t *restrict);	
11498		int	siqqueue(pid t,	int, union sigval);	
11499		int	sigsuspend (const	t sigset_t *);	
11500		int	sigtimedwait (con	nst sigset_t *restrict, siginfo_t *restrict,	
11501			const struct	t timespec *restrict);	
11502		int	sigwait(const s:	igset_t *restrict, int *restrict);	
11503		int	sigwaitinfo(con	<pre>st sigset_t *restrict, siginfo_t *restrict);</pre>	
11504		int	str2sig(const cl	har *restrict, int *restrict);	
		<b>T</b> 1 •			

11461 CX In addition, the following signal-specific information shall be available:

11505 CX Inclusion of the **<signal.h>** header may make visible all symbols from the **<time.h>** header.

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12398		int	<pre>putenv(char *);</pre>
12399		void	<pre>qsort(void *, size_t, size_t, int (*)(const void *,</pre>
12400			<pre>const void *));</pre>
12401	CX	void	<pre>qsort_r(void *, size_t, size_t, int (*)(const void *, +</pre>
12402			<pre>const void *, void *), void *);</pre> +
12403		int	rand(void);
12404	XSI	long	<pre>random(void);</pre>
12405		void	<pre>*realloc(void *, size_t);</pre>
12406	CX	void	<pre>*reallocarray(void *, size_t, size_t);</pre>
12407	XSI	char	<pre>*realpath(const char *restrict, char *restrict);</pre>
12408		unsigned short	t *seed48(unsigned short [3]);
12409	CX	int	<pre>setenv(const char *, const char *, int);</pre>
12410	OB XSI	void	<pre>setkey(const char *);</pre>
12411	XSI	char	<pre>*setstate(char *);</pre>
12412		void	<pre>srand(unsigned);</pre>
12413	XSI	void	<pre>srand48(long);</pre>
12414		void	<pre>srandom(unsigned);</pre>
12415		double	<pre>strtod(const char *restrict, char **restrict);</pre>
12416		float	<pre>strtof(const char *restrict, char **restrict);</pre>
12417		long	<pre>strtol(const char *restrict, char **restrict, int);</pre>
12418		long double	<pre>strtold(const char *restrict, char **restrict);</pre>
12419		long long	<pre>strtoll(const char *restrict, char **restrict, int);</pre>
12420		unsigned long	<pre>strtoul(const char *restrict, char *restrict, int);</pre>
12421		unsigned long	long
12422			<pre>strtoull(const char *restrict, char *restrict, int);</pre>
12423		int	<pre>system(const char *);</pre>
12424	XSI	int	unlockpt(int);
12425	CX	int	<pre>unsetenv(const char *);</pre>
12426		size_t	<pre>wcstombs(char *restrict, const wchar_t *restrict, size_t);</pre>
12427		int	<pre>wctomb(char *, wchar_t);</pre>
10400	CY	Inclusion of the	etdlib h> header may also make visible all symbols from <i>stantl</i> h> slimits h>
12428	CX		f h> and cove/wait h>
12429		<main.mz, <stude<="" td=""><td>inter, and systwart.</td></main.mz,>	inter, and systwart.
12430	APPLIC	ATION USAGE	
12431		None.	
12/22	RATION	JAIF	
12432	KAHOP	The ISO C stands	requires that exit (EXIT FAILURE) returns "unsuccessful termination
12455		statue" to the host	any ironmont. In a POSIX host any ironmont this maps that the lower 8 bits of
12434		FXIT FAILURE n	oust have at least one bit set. The standard developers decided to further
12455		rostrict the allowo	d values for the following reasons:
12430		iconici ne anowe	a variace for the following reasons.
12437		• Exit statuses	s of 126, 127, and greater than 128 are ambiguous in certain circumstances
12438		because they	have special meanings in the shell (see XCU Section 2.8.2, on page 2321).
12439		• The <i>xargs</i> uti	ility quits when a command execution exits with status 255 (see XCU <i>xargs</i> ).

Calling *exit*() with a value greater than 255 or less than 0 is something that only programs which are specifically designed to have their exit status obtained by *waitid*() should do (since it does not truncate the exit status to 8 bits). "Pure ISO C" programs that call exit (EXIT\_FAILURE) do not meet this design criterion.

12501 12502	NAME	VAME string.h — string operations				
12502						
12503 12504	4 #include <string.h></string.h>					
12505 12506 12507 12508	DESCRIPTION CX Some of the functionality described on this reference page extends the ISO C standard. Applications shall define the appropriate feature test macro (see XSH Section 2.2, on page 460) to enable the visibility of these symbols in this header.					
12509		The <b><strin< b=""></strin<></b>	<b>g.h&gt;</b> header shall define NULL and <b>size_t</b> as described in <b><stddef.h></stddef.h></b> .			
12510	СХ	The <b><strin< b=""></strin<></b>	<b>g.h&gt;</b> header shall define the <b>locale_t</b> type as described in <b><locale.h></locale.h></b> .			
12511 12512		The follow prototypes	ving shall be declared as functions and may also be defined as macros. Function s shall be provided for use with ISO C standard compilers.			
12513 12514 12515	XSI	void void int	<pre>*memccpy(void *restrict, const void *restrict, int, size_t); *memchr(const void *, int, size_t); memcmp(const void *, const void *, size_t);</pre>			
12516 12517 12518 12519	СХ	void void void void	<pre>*memcpy(void *restrict, const void *restrict, size_t); *memmem(const void *, size_t, const void *, size_t); *memmove(void *, const void *, size_t); *memset(void *, int, size t);</pre>	+		
12520 12521 12522	СХ	char char char	<pre>*stpcpy(char *restrict, const char *restrict); *stpncpy(char *restrict, const char *restrict, size_t); *strcat(char *restrict, const char *restrict);</pre>			
12523 12524 12525		char int int	<pre>*strchr(const char *, int); strcmp(const char *, const char *); strcoll(const char *, const char *);</pre>			
12526 12527 12528	СХ	int char size_t	<pre>strcoll_l(const char *, const char *, locale_t); *strcpy(char *restrict, const char *restrict); strcspn(const char *, const char *);</pre>			
12529 12530	CX	char char	<pre>*strdup(const char *); *strerror(int);</pre>			
12531 12532 12533	СХ	char int size t	<pre>*strerror_l(int, locale_t); strerror_r(int, char *, size_t); strlcat(char *restrict, const char *restrict, size t);</pre>	+		
12534 12535 12536 12537		size_t size_t char int	<pre>strlcpy(char *restrict, const char *restrict, size_t); strlen(const char *); *strncat(char *restrict, const char *restrict, size_t); strncmp(const char *, const char *, size_t);</pre>	+		
12538 12539 12540 12541	СХ	char char size_t char	<pre>*strncpy(char *restrict, const char *restrict, size_t); *strndup(const char *, size_t); strnlen(const char *, size_t); *strpbrk(const char *, const char *);</pre>			
12542 12543 12544	СХ	char char size_t	<pre>*strrchr(const char *, int); *strsignal(int); strspn(const char *, const char *); *strstr(const char *, const char *);</pre>			
12545 12546 12547	СХ	char char	<pre>*strict(const char *, const char *); *strtok(char *restrict, const char *restrict); *strtok_r(char *restrict, const char *restrict, char *restrict); </pre>	;		
12548		sıze_t	<pre>stixiim(char ^restrict, const char ^restrict, size_t);</pre>			

13725		pthread_t	Used to identify a thread.	+
13726		reclen_t	Used for directory entry lengths.	
13727		size_t	Used for sizes of objects.	
13728		ssize_t	Used for a count of bytes or an error indication.	
13729		suseconds_t	Used for time in microseconds.	
13730		time_t	Used for time in seconds.	
13731		timer_t	Used for timer ID returned by <i>timer_create(</i> ).	
13732		uid_t	Used for user IDs.	
13733 13734		All of the types shall be de exceptions:	efined as arithmetic types of an appropriate length, with the following	
13735 13736 13737 13738 13740 13741 13742 13743 13744 13745 13746 13747 13748		pthread_attr_t pthread_barrier_t pthread_barrierattr_t pthread_cond_t pthread_condattr_t pthread_key_t pthread_mutex_t pthread_mutexattr_t pthread_mutexattr_t pthread_rwlock_t pthread_rwlock_tt pthread_spinlock_t pthread_t timer_t		
13749		Additionally:		
13750		<ul> <li>mode_t shall be an in</li> </ul>	nteger type.	
13751		• <b>dev_t</b> shall be an inte	eger type.	
13752		<ul> <li>nlink_t, uid_t, gid_t</li> </ul>	, and <b>id_t</b> shall be integer types.	
13753		<ul> <li>blkcnt_t and off_t sh</li> </ul>	hall be signed integer types.	
13754		<ul> <li>fsblkcnt_t, fsfilcnt_t</li> </ul>	t, <b>reclen_t</b> , and <b>ino_t</b> shall be defined as unsigned integer types.	+
13755		<ul> <li>size_t shall be an un</li> </ul>	signed integer type.	
13756		<ul> <li>blksize_t, pid_t, and</li> </ul>	ssize_t shall be signed integer types.	
13757	СХ	<ul> <li>clock_t shall be an ir</li> </ul>	teger or real-floating type. <b>time_t</b> shall be an integer type.	
13758		The type <b>ssize_t</b> shall be c	apable of storing values at least in the range [-1, {SSIZE_MAX}].	
13759 13760	XSI	The type <b>suseconds_t</b> sha range [–1, 1 000 000].	all be a signed integer type capable of storing values at least in the	
13761 13762 13763 13764		The implementation shall of <b>blksize_t</b> , <b>pid_t</b> , <b>size_t</b> . The names of these progra the <i>getconf</i> utility.	support one or more programming environments in which the widths , <b>ssize_t</b> , and <b>suseconds_t</b> are no greater than the width of type <b>long</b> . amming environments can be obtained using the <i>confstr()</i> function or	
13765		There are no defined comp	parison or assignment operators for the following types:	

+

15234		int	<pre>dup2(int, int);</pre>
15235		int	<pre>dup3(int, int, int);</pre>
15236		void	_exit(int);
15237	OB XSI	void	<pre>encrypt(char [64], int);</pre>
15238		int	<pre>execl(const char *, const char *,);</pre>
15239		int	<pre>execle(const char *, const char *,);</pre>
15240		int	<pre>execlp(const char *, const char *,);</pre>
15241		int	<pre>execv(const char *, char *const []);</pre>
15242		int	<pre>execve(const char *, char *const [], char *const []);</pre>
15243		int	<pre>execvp(const char *, char *const []);</pre>
15244		int	<pre>faccessat(int, const char *, int, int);</pre>
15245		int	<pre>fchdir(int);</pre>
15246		int	<pre>fchown(int, uid_t, gid_t);</pre>
15247		int	<pre>fchownat(int, const char *, uid_t, gid_t, int);</pre>
15248	SIO	int	<pre>fdatasync(int);</pre>
15249		int	<pre>fexecve(int, char *const [], char *const []);</pre>
15250		pid_t	_Fork(void);
15251		pid_t	fork(void);
15252		long	<pre>fpathconf(int, int);</pre>
15253	FSC	int	fsync(int);
15254		int	<pre>ftruncate(int, off_t);</pre>
15255		char	<pre>*getcwd(char *, size_t);</pre>
15256		gid_t	getegid(void);
15257		int	<pre>getentropy(void *, size_t);</pre>
15258		uid_t	geteuid(void);
15259		gid_t	<pre>getgid(void);</pre>
15260		int	<pre>getgroups(int, gid_t []);</pre>
15261	XSI	long	gethostid(void);
15262		int	<pre>gethostname(char *, size_t);</pre>
15263		char	*getlogin(void);
15264		int	<pre>getlogin_r(char *, size_t);</pre>
15265		int	<pre>getopt(int, char * const [], const char *);</pre>
15266		pid_t	<pre>getpgid(pid_t);</pre>
15267		pid_t	getpgrp(void);
15268		pid_t	<pre>getpid(void);</pre>
15269		pid_t	getppid(void);
15270		pid_t	<pre>getsid(pid_t);</pre>
15271		uid_t	<pre>getuid(void);</pre>
15272		int	<pre>isatty(int);</pre>
15273		int	<pre>lchown(const char *, uid_t, gid_t);</pre>
15274		int	<pre>link(const char *, const char *);</pre>
15275		int	linkat(int, const char *, int, const char *, int);
15276	XSI	int	<pre>lockf(int, int, off_t);</pre>
15277		off_t	<pre>lseek(int, off_t, int);</pre>
15278	XSI	int	<pre>nice(int);</pre>
15279		long	<pre>pathconf(const char *, int);</pre>
15280		int	pause(void);
15281		int	pipe(int [2]);
15282		int	<pre>pipe2(int [2], int);</pre>
15283		int	<pre>posix_close(int, int);</pre>
15284		ssize_t	<pre>pread(int, void *, size_t, off_t);</pre>
15285		ssize_t	<pre>pwrite(int, const void *, size_t, off_t);</pre>

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		int	fputws(const wchar_t *restrict, FILE *restrict);	
15665		int	<pre>fwide(FILE *, int);</pre>	
15666		int	<pre>fwprintf(FILE *restrict, const wchar_t *restrict,);</pre>	
15667		int	<pre>fwscanf(FILE *restrict, const wchar_t *restrict,);</pre>	
15668		wint_t	getwc(FILE *);	
15669		wint_t	getwchar(void);	
15670		size_t	<pre>mbrlen(const char *restrict, size_t, mbstate_t *restrict);</pre>	
15671		size_t	<pre>mbrtowc(wchar_t *restrict, const char *restrict, size_t,</pre>	
15672			<pre>mbstate_t *restrict);</pre>	
15673		int	<pre>mbsinit(const mbstate_t *);</pre>	
15674	CX	size_t	<pre>mbsnrtowcs(wchar_t *restrict, const char **restrict,</pre>	
15675			<pre>size_t, size_t, mbstate_t *restrict);</pre>	
15676		size_t	<pre>mbsrtowcs(wchar_t *restrict, const char **restrict, size_t,</pre>	
15677			<pre>mbstate_t *restrict);</pre>	
15678	CX	FILE	<pre>*open_wmemstream(wchar_t **, size_t *);</pre>	
15679		wint_t	<pre>putwc(wchar_t, FILE *);</pre>	
15680		wint_t	<pre>putwchar(wchar_t);</pre>	
15681		int	<pre>swprintf(wchar_t *restrict, size_t,</pre>	
15682			<pre>const wchar_t *restrict,);</pre>	
15683		int	<pre>swscanf(const wchar_t *restrict,</pre>	
15684			<pre>const wchar_t *restrict,);</pre>	
15685		wint_t	<pre>ungetwc(wint_t, FILE *);</pre>	
15686		int	<pre>vfwprintf(FILE *restrict, const wchar_t *restrict, va_list);</pre>	
15687		int	<pre>vfwscanf(FILE *restrict, const wchar_t *restrict, va_list);</pre>	
15688		int	<pre>vswprintf(wchar_t *restrict, size_t,</pre>	
15689			<pre>const wchar_t *restrict, va_list);</pre>	
15690		int	<pre>vswscanf(const wchar_t *restrict, const wchar_t *restrict,</pre>	
15691			<pre>va_list);</pre>	
15692		int	<pre>vwprintf(const wchar_t *restrict, va_list);</pre>	
15693		int	<pre>vwscanf(const wchar_t *restrict, va_list);</pre>	
15694	CX	wchar_t	<pre>*wcpcpy(wchar_t *restrict, const wchar_t *restrict);</pre>	
15695		wchar_t	<pre>*wcpncpy(wchar_t *restrict, const wchar_t *restrict, size_t);</pre>	
15696		size_t	wortomb(char *restrict, wohar t, mbstate t *restrict):	
15697	СХ		wore on a fondi fondi footilioe, wondi_e, mobeliee_e footilioe,,	
15698		int	<pre>wcscasecmp(const wchar_t *, const wchar_t *);</pre>	
10070		int int	<pre>wcscasecmp_l(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t);</pre>	
15699		int int wchar_t	<pre>wcscasecmp_l(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict);</pre>	
15699 15700		int int wchar_t wchar_t	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t);</pre>	
15699 15700 15701		int int wchar_t wchar_t int	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *);</pre>	
15699 15700 15701 15702		int wchar_t wchar_t int int	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *);</pre>	
15699 15700 15701 15702 15703	СХ	<pre>int int wchar_t wchar_t int int int</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *);</pre>	
15699 15700 15701 15702 15703 15704	СХ	<pre>int int wchar_t wchar_t int int wchar_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *restrict);</pre>	
15699 15700 15701 15702 15703 15704 15705	СХ	<pre>int int wchar_t wchar_t int int wchar_t size_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *, locale_t); wcscspn(const wchar_t *, const wchar_t *);</pre>	
15699 15700 15701 15702 15703 15704 15705 15706	СХ	<pre>int int wchar_t wchar_t int int int wchar_t size_t wchar_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *restrict); wcscspn(const wchar_t *, const wchar_t *); *wcscup(const wchar_t *, const wchar_t *); *wcscup(const wchar_t *, const wchar_t *);</pre>	
15699 15700 15701 15702 15703 15703 15704 15705 15706 15707	сх	<pre>int int wchar_t wchar_t int int int wchar_t size_t wchar_t size_t size_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *restrict); wcscup(const wchar_t *, const wchar_t *, strict); wcscup(const wchar_t *, const wchar_t *); *wcsdup(const wchar_t *); wcsftime(wchar_t *restrict, size_t,</pre>	
15699 15700 15701 15702 15703 15704 15705 15706 15707 15708	сх	<pre>int int wchar_t wchar_t int int int wchar_t size_t wchar_t size_t size_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *); wcscspn(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, size_t, const wchar_t *restrict, size_t, const wchar_t *restrict, const struct tm *restrict);</pre>	
15699 15700 15701 15702 15703 15704 15705 15706 15707 15708 15709	CX CX CX	<pre>int int wchar_t wchar_t int int int wchar_t size_t wchar_t size_t size_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *); wcscdup(const wchar_t *, const wchar_t *); *wcsdup(const wchar_t *, const wchar_t *); *wcsdup(const wchar_t *, const wchar_t *); *wcsdup(const wchar_t *); wcsftime(wchar_t *restrict, size_t,</pre>	
15699 15700 15701 15702 15703 15704 15705 15706 15707 15708 15709 15710	CX CX CX	<pre>int int wchar_t wchar_t int int int wchar_t size_t wchar_t size_t size_t size_t size_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *restrict); wcscspn(const wchar_t *, const wchar_t *restrict); wcsclup(const wchar_t *, const wchar_t *); *wcsdup(const wchar_t *, const wchar_t *); *wcsdup(const wchar_t *, const wchar_t *); *wcsdup(const wchar_t *); wcsftime(wchar_t *restrict, size_t, const wchar_t *restrict, const struct tm *restrict); wcslcat(wchar_t *restrict, const wchar_t *restrict, + size_t); *</pre>	
15699 15700 15701 15702 15703 15704 15705 15706 15707 15708 15709 15710 15711	CX CX CX	<pre>int int wchar_t wchar_t int int int wchar_t size_t wchar_t size_t size_t size_t size_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *restrict); wcscspn(const wchar_t *, const wchar_t *restrict); wcsclup(const wchar_t *, const wchar_t *restrict); wcscftime(wchar_t *restrict, size_t, const wchar_t *restrict, const struct tm *restrict); wcslcat(wchar_t *restrict, const wchar_t *restrict, + size_t); wcslcpy(wchar_t *restrict, const wchar_t *restrict, + wcslcpy(wchar_t *restrict, const wchar_t *restrict,</pre>	
15699 15700 15701 15702 15703 15704 15705 15706 15707 15708 15709 15710 15711 15712	CX CX CX	<pre>int int wchar_t wchar_t int int int wchar_t size_t wchar_t size_t size_t size_t size_t size_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *restrict); wcscspn(const wchar_t *, const wchar_t *restrict); wcscspn(const wchar_t *, const wchar_t *); *wcsdup(const wchar_t *); wcsftime(wchar_t *restrict, size_t,</pre>	
15699 15700 15701 15702 15703 15704 15705 15706 15707 15708 15709 15710 15711 15712 15713	cx cx cx	<pre>int int wchar_t wchar_t int int int wchar_t size_t wchar_t size_t size_t size_t size_t size_t size_t</pre>	<pre>wcscasecmp(const wchar_t *, const wchar_t *); wcscasecmp_l(const wchar_t *, const wchar_t *, locale_t); *wcscat(wchar_t *restrict, const wchar_t *restrict); *wcschr(const wchar_t *, wchar_t); wcscmp(const wchar_t *, const wchar_t *); wcscoll(const wchar_t *, const wchar_t *); wcscoll_l(const wchar_t *, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *, locale_t); *wcscpy(wchar_t *restrict, const wchar_t *); wcscspn(const wchar_t *, const wchar_t *); *wcsdup(const wchar_t *, const wchar_t *); *wcslcat(wchar_t *restrict, size_t,</pre>	
15699 15700 15701 15702 15703 15704 15705 15706 15707 15708 15709 15710 15711 15712 15713 15714	CX CX CX	<pre>int int wchar_t wchar_t int int int wchar_t size_t wchar_t size_t size_t size_t size_t size_t int</pre>	<pre>worked worked and a second and a second and a second and a second a se</pre>	

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16135					Complete
16136		Header	Prefix	Suffix	Name
16137		<aio.h></aio.h>	aio_, lio_, AIO_, LIO_		
16138		<arpa inet.h=""></arpa>	inet_		
16139		<ctype.h></ctype.h>	to[a-z], is[a-z]		
16140		<dlfcn.h></dlfcn.h>	RTLD_, dli_		
16141		<dirent.h></dirent.h>	d_, DT_		
16142		<fcntl.h></fcntl.h>	1_		
16143	XSI	<fmtmsg.h></fmtmsg.h>	MM_		
16144		<fnmatch.h></fnmatch.h>	FNM_		
16145	XSI	<ftw.h></ftw.h>	FTW		
16146		<glob.h></glob.h>	gl_, GLOB_		
16147		<grp.h></grp.h>	gr_		
16148		<limits.h></limits.h>		_MAX, _MIN	
16149	XSI	<math.h></math.h>	M_		
16150	MSG	<mqueue.h></mqueue.h>	mq_, MQ_		
16151	XSI	<ndbm.h></ndbm.h>	dbm_, DBM_		
16152		<netdb.h></netdb.h>	ai_, h_, n_, p_, s_		
16153		<net if.h=""></net>	if_, IF_		
16154		<netinet in.h=""></netinet>	in_, ip_, s_, sin_, INADDR_, IPPROTO_		
16155	IP6		in6_, in6addr_, s6_, sin6_, IPV6_		
16156		<netinet tcp.h=""></netinet>	TCP_		
16157		<nl_types.h></nl_types.h>	NL_		
16158		<poll.h></poll.h>	pd_, ph_, ps_, POLL		
16159		<pthread.h></pthread.h>	pthread_, PTHREAD_		
16160		<pwd.h></pwd.h>	pw_		
16161		<regex.h></regex.h>	re_, rm_, REG_		
16162		<sched.h></sched.h>	sched_, SCHED_		
16163		<semaphore.h></semaphore.h>	sem_, SEM_		
16164	СХ	<signal.h></signal.h>	sa_, si_, sigev_, sival_, uc_, BUS_, CLD_,		
16165			FPE_, ILL_, SA_, SEGV_, SI_, SIGEV_,		
16166	XSI		ss_, sv_, SS_, TRAP_		
16167		<stdlib.h></stdlib.h>	str[a-z]		
16168		<string.h></string.h>	str[a-z], mem[a-z], wcs[a-z]		
16169	XSI	<sys ipc.h=""></sys>	ipc_, IPC_		key, pad, seq
16170		<sys mman.h=""></sys>	shm_, MAP_, MCL_, MS_,		
16171			PROT_		
16172	XSI	<sys msg.h=""></sys>	msg, MSG_[A-Z]		msg
16173	XSI	<sys resource.h=""></sys>	rlim_, ru_, PRIO_, RLIMIT_, RUSAGE_		
16174		<sys select.h=""></sys>	fd_, fds_, FD_		

16894	_Exit()	getpgrp()	read()	strncmp()
16895	_Fork()	getpid()	readlink()	strncpy()
16896	_ <i>exit</i> ()	getppid()	readlinkat()	strnlen()
16897	abort()	getsockname()	recv()	strpbrk()
16898	accept()	getsockopt()	recvfrom()	strrchr()
16899	accept4()	getuid()	recvmsg()	strspn()
16900	access()	htobe16()	rename()	strstr()
16901	aio_error()	htobe32()	renameat ( )	<pre>strtok_r()</pre>
16902	aio_return()	htobe64()	rmdir()	symlink()
16903	aio_suspend()	htole16()	select()	symlinkat()
16904	alarm()	htole32()	sem_post()	tcdrain()
16905	be16toh()	htole64()	send()	tcflow()
16906	be32toh()	htonl()	sendmsg()	tcflush()
16907	be64toh()	htons()	sendto()	tcgetattr()
16908	bind()	kill()	setegid()	tcgetpgrp()
16909	cfgetispeed()	le16toh()	seteuid()	tcgetwinsize()
16910	cfgetospeed()	le32toh()	setgid()	tcsendbreak()
16911	cfsetispeed()	le64toh()	setpgid()	tcsetattr()
16912	cfsetospeed()	link()	setregid()	tcsetpgrp()
16913	chdir()	linkat()	setreuid()	tcsetwinsize()
16914	chmod()	listen()	setsid()	time()
16915	chown()	longjmp()	setsockopt()	timer_getoverrun()
16916	clock_gettime()	lseek()	setuid()	timer_gettime()
16917	close()	lstat()	shutdown()	timer_settime()
16918	connect()	memccpy()	sig2str()	times()
16919	creat()	memchr()	sigaction()	umask()
16920	dup()	memcmp()	sigaddset()	uname()
16921	dup2()	memcpy()	sigdelset ()	unlink()
16922	dup3()	memmove()	sigemptyset()	unlinkat()
16923	execl()	memset()	sigfillset()	utimensat()
16924	execle()	mkdir()	sigismember()	utimes()
16925	execv()	mkdirat()	siglongjmp()	va_arg()
16926	execve()	mkfifo()	signal()	va_copy()
16927	faccessat()	mkfifoat()	sigpending()	va_end()
16928	fchdir()	mknod()	sigprocmask()	va_start()
16929	fchmod()	mknodat()	sigqueue()	wait()
16930	fchmodat()	ntohl()	sigsuspend()	waitpid()
16931	fchown()	ntohs()	sleep()	wcpcpy()
16932	fchownat()	open()	sockatmark()	wcpncpy()
16933	fcntl()	openat()	socket()	wcscat()
16934	fdatasync()	pause()	socketpair()	wcschr()
16935	fexecve()	pipe()	stat()	wcscmp()
16936	ffs()	pipe2()	stpcpy()	wcscpy()
16937	fstat()	poll()	stpncpy()	wcscspn()
16938	fstatat()	ppoll()	strcat()	wcslcat()
16939	fsync()	pread()	strchr()	wcslcpy()
16940	ftruncate()	pselect()	strcmp()	wcslen()
16941	futimens()	pthread_kill()	strcpy()	wcsncat()
16942	getegid()	pthread_self()	strcspn()	wcsncmp()
16943	geteuid()	<i>pthread_setcancelstate()</i>	strlcat()	wcsncpy()
16944	getgid()	<i>pthread_sigmask()</i>	strlcpy()	wcsnlen()
16945	getgroups()	pwrite()	strlen()	wcspbrk()
16946	getpeername()	raise()	strncat()	wcsrchr()

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17706 17707	<ul><li>If a thread is detached, its thread ID is invalid for use as an argument in a call to <i>pthread_de</i> or <i>pthread_join()</i>.</li></ul>		
17708	2.9.3	Thread Mutexes	
17709 17710 17711		A thread that has blocked shall not prevent any unblocked thread that is eligible to use the same processing resources from eventually making forward progress in its execution. Eligibility for processing resources is determined by the scheduling policy.	
17712		A thread shall become the owner of a mutex, <i>m</i> , when one of the following occurs:	
17713 17714 17715		• It calls <i>pthread_mutex_clocklock()</i> , <i>pthread_mutex_lock()</i> , <i>pthread_mutex_timedlock()</i> , or <i>pthread_mutex_trylock()</i> with <i>m</i> as the <i>mutex</i> argument and the call returns zero or [EOWNERDEAD].	
17716 17717		• It calls <i>pthread_mutex_setprioceiling()</i> with <i>m</i> as the <i>mutex</i> argument and the call returns [EOWNERDEAD].	-
17718 17719 17720		• It calls <i>pthread_cond_clockwait()</i> , <i>pthread_cond_timedwait()</i> , or <i>pthread_cond_wait()</i> with <i>m</i> as the <i>mutex</i> argument and the call returns zero or certain error numbers (see <i>pthread_cond_timedwait()</i> ).	+
17721		The thread shall remain the owner of <i>m</i> until one of the following occurs:	
17722		• It executes <i>pthread_mutex_unlock()</i> with <i>m</i> as the <i>mutex</i> argument	
17723 17724		• It blocks in a call to <i>pthread_cond_clockwait()</i> , <i>pthread_cond_timedwait()</i> , or <i>pthread_cond_wait()</i> with <i>m</i> as the <i>mutex</i> argument.	+ -
17725		The implementation shall behave as if at all times there is at most one owner of any mutex.	
17726 17727 17728		A thread that becomes the owner of a mutex is said to have ``acquired" the mutex and the mutex is said to have become ``locked"; when a thread gives up ownership of a mutex it is said to have ``released" the mutex and the mutex is said to have become ``unlocked".	
17729 17730 17731		A problem can occur if a process terminates while one of its threads holds a mutex lock. Depending on the mutex type, it might be possible for another thread to unlock the mutex and recover the state of the mutex. However, it is difficult to perform this recovery reliably.	
17732 17733 17734 17735 17736		Robust mutexes provide a means to enable the implementation to notify other threads in the event of a process terminating while one of its threads holds a mutex lock. The next thread that acquires the mutex is notified about the termination by the return value [EOWNERDEAD] from the locking function. The notified thread can then attempt to recover the state protected by the mutex, and if successful mark the state protected by the mutex as consistent by a call to	
17737		<i>pthread_mutex_consistent()</i> . If the notified thread is unable to recover the state, it can declare the	

Whether or not the state protected by a mutex can be recovered is dependent solely on the 17740 application using robust mutexes. The robust mutex support provided in the implementation 17741 provides notification only that a mutex owner has terminated while holding a lock, or that the 17742 state of the mutex is not recoverable. 17743

state as not recoverable by a call to *pthread\_mutex\_unlock()* without a prior call to

17738

17739

pthread\_mutex\_consistent().

17870

#### 17869 2.9.5.2 Cancellation Points

Cancellation points shall occur when a thread is executing the following functions:

17871	accept()	nanosleep()	recvmsg()
17872	accept4()	open()	select()
17873	aio_suspend()	openat()	send()
17874	clock_nanosleep()	pause()	sendmsg()
17875	close()	poll()	sendto()
17876	connect()	ppoll()	sigsuspend()
17877	creat()	pread()	sigtimedwait()
17878	fcntl()+	pselect()	sigwait()
17879	fdatasync()	pthread_cond_clockwait()	sigwaitinfo()
17880	fsync()	pthread_cond_timedwait()	sleep()
17881	lockf()++	pthread_cond_wait()	tcdrain()
17882	mq_receive()	pthread_join()	wait()
17883	$mq\_send()$	pthread_testcancel()	waitid()
17884	mq_timedreceive()	pwrite()	waitpid()
17885	$mq_timedsend()$	read()	write()
17886	msgrcv()	readv()	writev()
17887	msgsnd()	recv()	
17888	msync()	recvfrom()	
17889	A cancellation point ma	y also occur when a thread is execu	ting the following functions:
17890	access()	fchownat()	fseeko()
17891	asctime r()	fclose()	fsetnos()
17892	catclose()	fcntl()+++	fstat()
17893	catopen()	fflush()	fstatat()
17894	chmod()	fgetc()	ftell()
17895	chown()	faetnos()	ftello()
17896	closedir()	foets()	futimens()
17897	closelog()	fortup()	funrintf()
17898	ctermid()	fgetws()	fwrite()
17899	ctime r()	fmtmsg()	fwscanf()
17900	dlclose()	fopen()	getaddrinfo()
17901	dloven()	fpathconf()	getc()
17902	dprintf()	fprintf()	getc unlocked()
17903	endhostent()	fputc()	getchar()
17904	endnetent()	fputs()	getchar unlocked()
17905	endprotoent()	fputwc()	getcwd()
17906	endservent()	fputws()	getdelim()
17907	faccessat()	fread()	getgrgid_r()
17908	fchmod()	freopen()	getgrnam_r()
17909	fchmodat()	fscanf()	gethostid()
17910	fchown()	fseek()	gethostname()

17911 † When the *cmd* argument is F\_SETLKW.

17913 ††† For any value of the *cmd* argument.

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<sup>17912</sup>  $\dagger$ † When the *function* argument is F\_LOCK.

17914	getline()	posix_openpt()	sem_wait()
17915	getlogin_r()	posix_spawn()	semop()
17916	getnameinfo()	posix_spawnp()	sethostent()
17917	getpwnam_r()	<pre>posix_typed_mem_open()</pre>	setnetent()
17918	getpwuid_r()	printf()	setprotoent()
17919	getwc()	psiginfo()	setservent()
17920	getwchar()	psignal()	stat()
17921	glob()	<pre>pthread_rwlock_clockrdlock()</pre>	strerror_l()
17922	iconv_close()	<pre>pthread_rwlock_clockwrlock()</pre>	<pre>strerror_r()</pre>
17923	iconv_open()	pthread_rwlock_rdlock()	strftime()
17924	link()	<pre>pthread_rwlock_timedrdlock()</pre>	strftime_l()
17925	linkat()	<pre>pthread_rwlock_timedwrlock()</pre>	symlink()
17926	lio_listio()	<pre>pthread_rwlock_wrlock()</pre>	symlinkat()
17927	<i>localtime_r()</i>	ptsname()	sync()
17928	lockf()	ptsname_r()	syslog()
17929	lseek()	putc()	tmpfile()
17930	lstat()	putc_unlocked()	tmpnam()
17931	mkdir()	putchar()	ttyname_r()
17932	mkdirat()	putchar_unlocked()	tzset()
17933	mkdtemp()	puts()	ungetc()
17934	mkfifo()	putwc()	ungetwc()
17935	mkfifoat()	putwchar()	unlink()
17936	mknod()	readdir_r()	unlinkat()
17937	mknodat()	readlink()	utimensat()
17938	mkstemp()	readlinkat()	utimes()
17939	mktime()	remove()	vdprintf()
17940	opendir()	rename()	vfprintf()
17941	openlog()	renameat()	vfwprintf()
17942	pathconf()	rewind()	vprintf()
17943	perror()	rewinddir()	vwprintf()
17944	popen()	scandir()	wcsftime()
17945	<pre>posix_fadvise()</pre>	scanf()	wordexp()
17946	posix_fallocate()	seekdir()	wprintf()
17947	posix_getdents()	sem_clockwait()	wscanf()
17948	posix_madvise()	sem_timedwait()	

- 17949In addition, a cancellation point may occur when a thread is executing any function that this17950standard does not require to be thread-safe but the implementation documents as being thread-17951safe. If a thread is cancelled while executing a non-thread-safe function, the behavior is17952undefined.
- 17953An implementation shall not introduce cancellation points into any other functions specified in17954this volume of POSIX.1-202x.
- 17955The side-effects of acting upon a cancellation request while suspended during a call of a function17956are the same as the side-effects that may be seen in a single-threaded program when a call to a17957function is interrupted by a signal and the given function returns [EINTR]. Any such side-17958effects occur before any cancellation cleanup handlers are called. For functions that are explicitly17959required not to return when interrupted (for example, *pclose()*), if a thread is canceled while17960executing the function, the behavior is undefined.
- 17961Whenever a thread has cancelability enabled and a cancellation request has been made with that17962thread as the target, and the thread then calls any function that is a cancellation point (such as17963*pthread\_testcancel()* or *read()*), the cancellation request shall be acted upon before the function17964returns. If a thread has cancelability enabled and a cancellation request is made with the thread

### 18595 **2.11.1 Defined Types**

18596 18597 18598 All of the data types used by various functions are defined by the implementation. The following table describes some of these types. Other types referenced in the description of a function, not mentioned here, can be found in the appropriate header for that function.

18599	Defined Type	Description
18600	cc_t	Type used for terminal special characters.
18601	clock_t	Integer or real-floating type used for processor times, as defined in
18602		the ISO C standard.
18603	clockid_t	Used for clock ID type in some timer functions.
18604	dev_t	Integer type used for device numbers.
18605	DIR	Type representing a directory stream.
18606	div_t	Structure type returned by the $div()$ function.
18607	FILE	Structure containing information about a file.
18608	glob_t	Structure type used in pathname pattern matching.
18609	fpos_t	Type containing all information needed to specify uniquely every
18610		position within a file.
18611	gid_t	Integer type used for group IDs.
18612	iconv_t	Type used for conversion descriptors.
18613	id_t	Integer type used as a general identifier; can be used to contain
18614		at least the largest of a <b>pid_t</b> , <b>uid_t</b> , or <b>gid_t</b> .
18615	ino_t	Unsigned integer type used for file serial numbers.
18616	key_t	Arithmetic type used for XSI interprocess communication.
18617	ldiv_t	Structure type returned by the <i>ldiv</i> () function.
18618	mode_t	Integer type used for file attributes.
18619	mqd_t	Used for message queue descriptors.
18620	nfds_t	Integer type used for the number of file descriptors.
18621	nlink_t	Integer type used for link counts.
18622	off_t	Signed integer type used for file sizes.
18623	pid_t	Signed integer type used for process and process group IDs.
18624	pthread_attr_t	Used to identify a thread attribute object.
18625	pthread_cond_t	Used for condition variables.
18626	pthread_condattr_t	Used to identify a condition attribute object.
18627	pthread_key_t	Used for thread-specific data keys.
18628	pthread_mutex_t	Used for mutexes.
18629	pthread_mutexattr_t	Used to identify a mutex attribute object.
18630	pthread_once_t	Used for dynamic package initialization.
18631	pthread_rwlock_t	Used for read-write locks.
18632	pthread_rwlockattr_t	Used for read-write lock attributes.
18633	pthread_t	Used to identify a thread.
18634	ptrdiff_t	Signed integer type of the result of subtracting two pointers.
18635	reclen_t	Unsigned integer type used for directory entry lengths.
18636	regex_t	Structure type used in regular expression matching.
18637	regmatch_t	Structure type used in regular expression matching.
18638	rlim_t	Unsigned integer type used for limit values, to which objects of
18639		type int and off_t can be cast without loss of value.
18640	sem_t	Type used in performing semaphore operations.
18641	sig_atomic_t	Possibly volatile-qualified integer type of an object that can be
18642		accessed as an atomic entity, even in the presence of asynchronous
18643	•	interrupts.
18644	sigset_t	Integer or structure type of an object used to represent sets

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21037	NAME	hind — hind a na	me to a socket	
21030	SVNOD			
21039 21040	51101	#include <sys< td=""><td>s/socket.h&gt;</td></sys<>	s/socket.h>	
21041 21042		int bind(int socklen_t	socket, const struct sockaddr *address, address_len);	
21043 21044 21045 21046	DESCR	<b>DESCRIPTION</b> The <i>bind()</i> function shall assign a local socket address <i>address</i> to a socket identified by descriptor <i>socket</i> that has no local socket address assigned. Sockets created with the <i>socket()</i> function are initially unnamed; they are identified only by their address family.		
21047		The <i>bind</i> () functi	on takes the following arguments:	
21048		socket	Specifies the file descriptor of the socket to be bound.	
21049 21050 21051		address	Points to a <b>sockaddr</b> structure containing the address to be bound to the socket. The length and format of the address depend on the address family of the socket.	
21052 21053		address_len	Specifies the length of the <b>sockaddr</b> structure pointed to by the <i>address</i> argument.	
21054 21055		The socket specif <i>bind()</i> function.	Tied by <i>socket</i> may require the process to have appropriate privileges to use the	
21056 21057		If the address family of the socket is AF_UNIX and the pathname in <i>address</i> names a symbolic link, <i>bind()</i> shall fail and set <i>errno</i> to [EADDRINUSE].		
21058 21059 21060 21061 21062		If the socket address cannot be assigned immediately and O_NONBLOCK is set for the file descriptor for the socket, <i>bind()</i> shall fail and set <i>errno</i> to [EINPROGRESS], but the assignment request shall not be aborted, and the assignment shall be completed asynchronously. Subsequent calls to <i>bind()</i> for the same socket, before the assignment is completed, shall fail and set <i>errno</i> to [EALREADY].		
21063 21064		When the assign shall indicate tha	ment has been performed asynchronously, <i>pselect(), select(), poll(),</i> and <i>ppoll()</i> t the file descriptor for the socket is ready for reading and writing.	
21065 21066 21067	<ul> <li><b>RETURN VALUE</b></li> <li>Upon successful completion, <i>bind</i>() shall return 0; otherwise, -1 shall be returned and <i>errno</i> set to indicate the error.</li> </ul>			
21068 21069	ERROR	. <b>S</b> The <i>bind</i> () functi	on shall fail if:	
21070		[EADDRINUSE]	The specified address is already in use.	
21071 21072		[EADDRNOTAV	AIL] The specified address is not available from the local machine.	
21073 21074 21075		[EAFNOSUPPOI	The specified address is not a valid address for the address family of the specified socket.	
21076		[EALREADY]	An assignment request is already in progress for the specified socket.	
21077		[EBADF]	The <i>socket</i> argument is not a valid file descriptor.	

2357       connect — connect a socket         2558       #include <sys socket.b="">         2560       int connect (int socket, const struct sockaddr *address, socklen_t address_len);         2561       DESCRIPTION         2562       The connect() function shall attempt to make a connection on a connection-mode socket or to set or reset the peer address of a connectionless-mode socket. The function takes the following arguments:         2566       socket       Specifies the file descriptor associated with the socket.         2567       address       Points to a sockaddr structure containing the peer address. The length and format of the address depend on the address family of the socket.         2568       address_len       Specifies the length of the sockaddr structure pointed to by the address argument.         2571       If the socket has not already been bound to a local address, connect() shall bind it to an address.         2572       which, unless the socket's address family is AF_UNIX, is an unused local address.         2573       If the socket has not already been bound to a local address. (be peer address. fand no connection is made. For SOCK_DGRAM sockets, the peer address is dentifies where all datagrams are sent on subsequent sod() functions, and limits the remote sender for subsequent corcl) shall be reset. Note that despite no connection being made, the term ``connected'' is used to describe a connectionelises is AF UNSPEC. (the socket's peer address as hall be reset. Note that despite no connect() shall fail and the connection in the sout put to a unspecified timout interval aurit the con</sys>	23556	NAME			
2359       ¥include <sys socket.h="">         2359       #include <sys socket.h="">         2350       int connect(int socket, const struct sockaddr *address, socklen_t address_len);         2350       DESCRIPTION         2356       The connect() function shall attempt to make a connection on a connection-mode socket or to set or reset the peer address of a connectionless-mode socket. The function takes the following arguments:         2356       socket       Specifies the file descriptor associated with the socket.         2357       address       Points to a sockaddr structure containing the peer address. The length and format of the address depend on the address family of the socket.         2358       address_len       Specifies the length of the sockaddr structure pointed to by the address argument.         2357       If the socket has not already been bound to a local address, connect() shall bind it to an address which, unless the socket's address family is AF_UNIX, is an unused local address.         2357       If the initiating socket is not connection-mode, then connect() shall set the socket's peer address shall be reset. Note that despite no connect() shall set the socket's peer address shall be reset. Note that despite no connect() shall attempt to establish a connection is stablished.         2357       If the initiating socket is not secontection-mode, then connect() shall attempt to establish a connection in the valdress is AF_UNSPEC, the socket's peer address shall be reset. Note that despite no connection is made, the term `connected'' is used to describe a connectioneless-mode socket for</sys></sys>	23557		connect — conne	ct a socket	
<ul> <li>#include <sys socket.h=""></sys></li> <li>int connect(int socket, const struct sockaddr *address, socklen_t address_len);</li> <li>DESCRIPTION</li> <li>DESCRIPTION</li> <li>The connect() function shall attempt to make a connection on a connection-mode socket or to set or reset the peer address of a connectionless-mode socket. The function takes the following arguments:</li> <li>socket Specifies the file descriptor associated with the socket.</li> <li>address Points to a sockaddr structure containing the peer address. The length and format of the address depend on the address family of the socket.</li> <li>address Int Specifies the length of the sockaddr structure pointed to by the address argument.</li> <li>If the socket has not already been bound to a local address, connect() shall bind it to an address which, unless the socket's address family is AF_UNIX, is an unused local address.</li> <li>If the initiating socket is not connection-mode, then connect() shall set the socket's peer address, and no connection is made. For SOCK_DCRAM sockets, the peer address identifies where all tadagrams are sent on subsequent send() functions, and limits the remote sender for subsequent rect() functions. If the safamily member of address is AF_UNSPEC, the socket's peer address is all be reset. Note that despite no connection being made, the term `connected'' is used to describe a connectionless-mode socket for which a peer address has been set.</li> <li>If the initiating socket is connection-mode, then connect() shall attempt to establish a connection is established, connect() shall fail and the connection is setablished immediately and O_NONBLOCK is not set for the file descriptor for the socket, connect() shall fail and set errno to EINTRI, but the connection set to the address argument.</li> <li>If the initiating to established asynchronously. Subsequent interval expires before the connection shall be established asynchronously.</li> <li>If the connection cannot be estab</li></ul>	23558	SYNOP	SIS		
<ul> <li>int connect (int socket, const struct sockaddr *address, socklen_t address_len);</li> <li>DESCRIPTION</li> <li>The connect() function shall attempt to make a connection on a connection-mode socket or to set or reset the peer address of a connectionless-mode socket. The function takes the following arguments:</li> <li>socket Specifies the file descriptor associated with the socket.</li> <li>address Points to a sockaddr structure containing the peer address. The length and format of the address depend on the address family of the socket.</li> <li>address_len Specifies the length of the sockaddr structure pointed to by the address argument.</li> <li>If the socket has not already been bound to a local address, connect() shall bind it to an address which, unless the socket's address family is AF_UNIX, is an unused local address.</li> <li>If the initiating socket is not connection-mode, then connect() shall set the socket's peer address, and no connection is made. For SOCK_DCRAM sockets, the peer address identifies where all datagrams are sent on subsequent send() functions, and limits the remote sender for subsequent recc() functions. If the sa_fmily member of address is AF_UNSPEC, the socket's peer address shall be reset. Note that despite no connection being made, the term `connected' is used to describe a connectionless-mode socket for which a peer address has been set.</li> <li>If the initiating socket is connection-mode, then connect() shall atial and the connection interval expires before the connection is established. If the timeout interval expires before the connection is established. If the timeout interval expires before the connection is established. If the connection attempt shall be aborted, and the connection shall be established asynchronously.</li> <li>If the connection cannot be established and the connection attempt shall be aborted, and the connection shall be established, shall fail and set errno to [ELNPROCRESS], but the connection request shall</li></ul>	23559		<pre>#include <sys< pre=""></sys<></pre>	s/socket.h>	
2252 <b>DESCRIPTION</b> 2253       The connect() function shall attempt to make a connection-mode socket or to set or reset the peer address of a connectionless-mode socket. The function takes the following arguments:         22566       socket       Specifies the file descriptor associated with the socket.         22567       address       Points to a sockaddr structure containing the peer address. The length and format of the address depend on the address family of the socket.         22569       address_len       Specifies the length of the sockaddr structure pointed to by the address argument.         22571       If the socket has not already been bound to a local address, connect() shall bind it to an address which, unless the socket's address family is AF_UNIX, is an unused local address.         22573       If the initiating socket is not connection-mode, then connect() shall set the socket's peer address.         22574       and no connection is made. For SOCK_DGRAM sockets, the peer address identifies where all datagrams are sent on subsequent send() functions, and limits the remote sender for subsequent recv() functions. If the sa_family member of address is AF_UNSPEC, the socket's peer address shall be reset. Note that despite no connection being made, the term ``connected'' is used to describe a connectionless-mode socket for which a peer address has been set.         22573       If the initiating socket is connection interval until the connection is established connection interval expires before the connection is established, connect() shall fail and the connection interval expires before the connection shall be established. If the timeout interval	23560 23561		int connect(i socklen_t	.nt <i>socket,</i> const struct sockaddr * <i>address,</i> : address_len);	
2353       The connect() function shall attempt to make a connection on a connection-mode socket or to set         2354       or reset the peer address of a connectionless-mode socket. The function takes the following         2355       arguments:         2356       socket       Specifies the file descriptor associated with the socket.         2357       address       Points to a sockaddr structure containing the peer address. The length and format of the address depend on the address family of the socket.         2358       address_len       Specifies the length of the sockaddr structure pointed to by the address argument.         2357       If the socket has not already been bound to a local address, connect() shall bind it to an address.         2357       If the initiating socket's address family is AF_UNIX, is an unused local address.         2357       and no connection is made. For SOCK_DGRAM sockets, the peer address identifies where all datagrams are sent on subsequent send() functions, and limits the emote sender for subsequent recr() functions. If the sa_family member of address is AF_UNSPEC, the socket's seer address specified by the address argument.         2358       the initiating socket is connection-mode, then connect() shall stempt to establish a connection to the address specified by the address recr() functions. If the sa_family member of address is AF_UNSPEC, the socket's ener address is able be active address specified by the address argument.         2357       If the initiating socket is not connection-mode, then connect() shall stempt to establish a connection is established.	23562	DESCR	IPTION		
2354       or reset the peer address of a connectionless-mode socket. The function takes the following         2356       socket       Specifies the file descriptor associated with the socket.         2356       address       Points to a sockaddr structure containing the peer address. The length and         2356       address       Points to a sockaddr structure containing the peer address. The length and         2356       address       Points to a sockaddr structure pointed to by the address         2357       address_len       Specifies the length of the sockaddr structure pointed to by the address         2357       address_len       Specifies the socked's and unused local address.         2357       If the socket has not already been bound to a local address, connect() shall bind it to an address         2357       If the initiating socket is not connection-mode, then connect() shall set the socket's peer address,         2357       and no connection is made. For SOCK_DGRAM sockets, the peer address identifies where all         2357       the initiating socket is connection being made, the term ``connected'' is used to         2358       describe a connection-mode, then connect() shall attempt to establish a connection         2359       If the initiating socket is connection-mode, then connect() shall attempt to establish a connection         2357       shall be reset. Note that despite no connect() shall attempt to establish a connection	23563		The connect() fur	action shall attempt to make a connection on a connection-mode socket or to set	
2356       arguments:         2356       socket       Specifies the file descriptor associated with the socket.         2356       address       Points to a sockaddr structure containing the peer address. The length and format of the address depend on the address family of the socket.         2359       address_len       Specifies the length of the sockaddr structure pointed to by the address argument.         2357       If the socket has not already been bound to a local address, connect() shall bind it to an address, which, unless the socket's address family is AF_UNIX, is an unused local address.         2357       If the initiating socket is not connection-mode, then connect() shall set the socket's peer address, and no connection is made. For SOCK_DGRAM sockets, the peer address is shall be reset. Note that despite no connection being made, the term `connected" is used to describe a connectionless-mode socket for which a peer address has been set.         2357       If the initiating socket is connection-mode, then connect() shall attempt to establish a connection being made, the term `connected" is used to describe a connectionless-mode socket for which a peer address has been set.         2357       If the initiating socket is connection is restablished, connect() shall attempt to establish a connection interval until the connection is established. If the timeout interval any peer before the connection is established, connect() shall fail and the connection interval expires before the connection shall be established, south the connection interval expires before the connection shall be established, connect() shall fail and the connection request shall not be aborted, and the connection shall be esta	23564		or reset the peer	address of a connectionless-mode socket. The function takes the following	
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23569address_lenSpecifies the length of the sockaddr structure pointed to by the address argument.23571If the socket has not already been bound to a local address, connect() shall bind it to an address which, unless the socket's address family is AF_UNIX, is an unused local address.23573If the initiating socket is not connection-mode, then connect() shall set the socket's peer address, and no connection is made. For SOCK_DGRAM sockets, the peer address identifies where all datagrams are sent on subsequent send() functions, and limits the remote sender for subsequent recv() functions. If the sa_family member of address is AF_UNSPEC, the socket's peer address shall be reset. Note that despite no connection being made, the term ``connected'' is used to describe a connectionless-mode socket for which a peer address has been set.23579If the initiating socket is connection-mode, then connect() shall attempt to establish a connection to the address specified by the address argument. If the connection cannot be established immediately and O_NONBLOCK is not set for the file descriptor for the socket, connect() shall the uniterval expires before the connect() shall fail and the connection attempt shall be aborted. If connect() shall fail and set errno to [EINTR], but the connection attempt shall be aborted, and the connection shall be established asynchronously.2358If the connection cannot be established immediately and O_NONBLOCK is set for the file descriptor for the socket, set for the file descriptor for the socket, connection stall the connection attempt shall be aborted, and the connection shall be established asynchronously.2358If the connection cannot be established immediately and O_NONBLOCK is set for the file descriptor for the socket, connect() shall fail and set errno to [EIN	23567 23568		address	Points to a <b>sockaddr</b> structure containing the peer address. The length and format of the address depend on the address family of the socket.	
23571If the socket has not already been bound to a local address, connect() shall bind it to an address23572which, unless the socket's address family is AF_UNIX, is an unused local address.23573If the initiating socket is not connection-mode, then connect() shall set the socket's peer address,23574and no connection is made. For SOCK_DGRAM sockets, the peer address identifies where all23575datagrams are sent on subsequent send() functions, and limits the remote sender for subsequent2576recv() functions. If the sa_family member of address is AF_UNSPEC, the socket's peer address2577shall be reset. Note that despite no connection being made, the term ``connected'' is used to2587describe a connectionless-mode socket for which a peer address has been set.2587lf the initiating socket is connection-mode, then connect() shall attempt to establish a connection2580to the address specified by the address argument. If the connection cannot be established2581immediately and O_NONBLOCK is not set for the file descriptor for the socket, connect() shall2582block for up to an unspecified timeout interval until the connection is established. If the timeout2583interval expires before the connect() shall fail and the connection2584attempt shall be aborted, and the connect() shall fail and set errno to [EINTR], but the connection2585waiting to establish a connect() shall fail and set errno to [EINPROCRESS], but the connection2586request shall not be aborted, and the connection shall be established asynchronously.2587lf the connection cannot be established immediately and O_NONBLOCK is	23569 23570		address_len	Specifies the length of the <b>sockaddr</b> structure pointed to by the <i>address</i> argument.	
<ul> <li>which, unless the socket's address family is AF_UNIX, is an unused local address.</li> <li>If the initiating socket is not connection-mode, then <i>connect()</i> shall set the socket's peer address, and no connection is made. For SOCK_DGRAM sockets, the peer address identifies where all datagrams are sent on subsequent <i>send()</i> functions, and limits the remote sender for subsequent <i>recv()</i> functions. If the <i>sa_family</i> member of <i>address</i> is AF_UNSPEC, the socket's peer address</li> <li>the initiating socket is connection-mode, then <i>connect()</i> shall set the socket's peer address</li> <li>the initiating socket is connection-mode, then <i>connect()</i> shall attempt to establish a connection to the address specified by the <i>address</i> argument. If the connection cannot be established immediately and O_NONBLOCK is not set for the file descriptor for the socket, <i>connect()</i> shall stall and the connection attempt shall be aborted. If <i>connect()</i> shall fail and set <i>errno</i> to [EINTR], but the connection attempt shall be aborted. If <i>connect()</i> shall fail and set <i>errno</i> to [EINTR], but the connection argument attempt to establish a connection attempt shall not be aborted, and the connection shall be established asynchronously.</li> <li>If the connection cannot be established immediately and O_NONBLOCK is set for the file descriptor for the socket attempt shall be aborted. If <i>connect()</i> shall fail and set <i>errno</i> to [EINTR], but the connection establish a connection shall be established asynchronously.</li> <li>If the connection cannot be established immediately and O_NONBLOCK is set for the file descriptor for the socket, <i>connect()</i> shall fail and set <i>errno</i> to [EINTR], but the connection estables descriptor for the socket, <i>connect()</i> shall fail and set <i>errno</i> to [EINTR], but the connection attempt shall not be aborted, and the connection shall be established asynchronously.</li></ul>	23571		If the socket has	not already been bound to a local address. <i>connect()</i> shall bind it to an address	
23573If the initiating socket is not connection-mode, then connect() shall set the socket's peer address,23574and no connection is made. For SOCK_DGRAM sockets, the peer address identifies where all23575datagrams are sent on subsequent send() functions, and limits the remote sender for subsequent23576recv() functions. If the sa_family member of address is AF_UNSPEC, the socket's peer address23577shall be reset. Note that despite no connection being made, the term ``connected'' is used to23578describe a connectionless-mode socket for which a peer address has been set.23579If the initiating socket is connection-mode, then connect() shall attempt to establish a connection23580to the address specified by the address argument. If the connection cannot be established23581immediately and O_NONBLOCK is not set for the file descriptor for the socket, connect() shall23582block for up to an unspecified timeout interval until the connect() shall fail and the connection23583attempt shall be aborted. If connect() is interrupted by a signal that is caught while blocked23584waiting to establish a connection, connect() shall fail and set errno to [EINTR], but the connection23585waiting to establish a connect() shall fail and set errno to [EINTR], but the connection23586request shall not be aborted, and the connection shall be established asynchronously.23587If the connection cannot be established immediately and O_NONBLOCK is set for the file23588descriptor for the socket, connect() shall fail and set errno to [EINTR], but the connection23589request shall not be aborted, and the	23572		which, unless the	socket's address family is AF_UNIX, is an unused local address.	
23574and no connection is made. For SOCK_DGRAM sockets, the peer address identifies where all23575datagrams are sent on subsequent send() functions, and limits the remote sender for subsequent2576recv() functions. If the sa_family member of address is AF_UNSPEC, the socket's peer address2577shall be reset. Note that despite no connection being made, the term ``connected" is used to2578describe a connectionless-mode socket for which a peer address has been set.2579If the initiating socket is connection-mode, then connect() shall attempt to establish a connection2580to the address specified by the address argument. If the connection cannot be established2581immediately and O_NONBLOCK is not set for the file descriptor for the socket, connect() shall2582block for up to an unspecified timeout interval until the connection is established. If the timeout2583interval expires before the connect() shall fail and set errno to [EINTR], but the connection2584attempt shall be aborted, and the connection shall be established asynchronously.2585Uf the connection cannot be established immediately and O_NONBLOCK is set for the file2586descriptor for the socket, connect() shall fail and set errno to [EINPROGRESS], but the connection25872588descriptor for the socket, and the connection shall be established asynchronously.2589connect() for the same socket, before the connection is established, shall fail and set errno2589connect() for the same socket, before the connection is established, shall fail and set errno2589connect() for the same socket, before the connection is es	23573		If the initiating se	ocket is not connection-mode, then <i>connect()</i> shall set the socket's peer address,	
23575datagrams are sent on subsequent send() functions, and limits the remote sender for subsequent23576recv() functions. If the sa_family member of address is AF_UNSPEC, the socket's peer address23577shall be reset. Note that despite no connection being made, the term "connected" is used to23578describe a connectionless-mode socket for which a peer address has been set.23579If the initiating socket is connection-mode, then connect() shall attempt to establish a connection23580to the address specified by the address argument. If the connection cannot be established23581immediately and O_NONBLOCK is not set for the file descriptor for the socket, connect() shall23582block for up to an unspecified timeout interval until the connection is established. If the timeout23583interval expires before the connect() is interrupted by a signal that is caught while blocked23584attempt shall be aborted. If connect() is all fail and set errno to [EINTR], but the connection23585uaiting to establish a connect() shall fail and set errno to [EINTR], but the connection23586request shall not be aborted, and the connection shall be established asynchronously.23587If the connection cannot be established immediately and O_NONBLOCK is set for the file23588descriptor for the socket, connect() shall fail and set errno to [EINTR], but the connection2359request shall not be aborted, and the connection shall be established asynchronously.2358lf the connection cannot be established immediately and O_NONBLOCK is set for the file2358descriptor for the socket, connect() shall fail and set errno to	23574		and no connection	on is made. For SOCK_DGRAM sockets, the peer address identifies where all	
<ul> <li><i>recv</i>() functions. If the <i>sa_family</i> member of <i>address</i> is AF_UNSPEC, the socket's peer address</li> <li>shall be reset. Note that despite no connection being made, the term ``connected" is used to</li> <li>describe a connectionless-mode socket for which a peer address has been set.</li> <li>If the initiating socket is connection-mode, then <i>connect</i>() shall attempt to establish a connection</li> <li>to the address specified by the <i>address</i> argument. If the connection cannot be established</li> <li>immediately and O_NONBLOCK is not set for the file descriptor for the socket, <i>connect</i>() shall</li> <li>block for up to an unspecified timeout interval until the connection is established. If the timeout</li> <li>interval expires before the connect() is interrupted by a signal that is caught while blocked</li> <li>waiting to establish a connection, <i>connect</i>() shall fail and set <i>errno</i> to [EINTR], but the connection</li> <li>request shall not be aborted, and the connection shall be established asynchronously.</li> <li>If the connect() for the same socket, <i>before</i> the connection shall be established asynchronously. Subsequent</li> <li>calls to <i>connect</i>() for the same socket, before the connection is established, shall fail and set <i>errno</i></li> <li>to [EALREADY].</li> <li>When the connection has been established asynchronously, <i>pselect(), select(), poll()</i>, and <i>ppoll()</i></li> <li>shall indicate that the file descriptor for the socket is ready for writing.</li> </ul>	23575		datagrams are sent on subsequent <i>send()</i> functions, and limits the remote sender for subsequent		
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23591to [EALREADY].23592When the connection has been established asynchronously, <i>pselect(), select(), poll(),</i> and <i>ppoll()</i> 23593shall indicate that the file descriptor for the socket is ready for writing.	23590		calls to connect()	for the same socket, before the connection is established, shall fail and set <i>errno</i>	
When the connection has been established asynchronously, <i>pselect()</i> , <i>select()</i> , <i>poll()</i> , and <i>ppoll()</i> shall indicate that the file descriptor for the socket is ready for writing.	23591		to [EALREADY].		
shall indicate that the file descriptor for the socket is ready for writing.	23592		When the connec	tion has been established asynchronously, <i>vselect()</i> , <i>select()</i> , <i>voll()</i> , and <i>vvoll()</i>	
	23593		shall indicate tha	t the file descriptor for the socket is ready for writing.	
The socket in use may require the process to have appropriate privileges to use the <i>connect</i> () function.	23594 23595		The socket in us function	e may require the process to have appropriate privileges to use the <i>connect()</i>	
		DETTIN			
235% KEIUKIN VALUE	23596	KEIUK	IN VALUE	completion connect() shall not um () otherwise 1 shall be not um of and and	
23597 Opon succession completion, <i>connect()</i> shall return 0; otherwise, -1 shall be returned and <i>errno</i>	23597		set to indicate the	completion, connect() shall return 0, otherwise, -1 shall be returned and ermo	

# dladdr()

+

		+
24808 24809 24810 24811	NAME dladdr — get information relating to an address SYNOPSIS #include <dlfcn.h></dlfcn.h>	+ + + + +
24812	<pre>int dladdr(const void *restrict addr, Dl_info_t *restrict dlip);</pre>	+
24813 24814 24815 24816 24817 24818 24819 24820 24821 24822 24823 24824 24825	<b>DESCRIPTION</b> The <i>dladdr</i> () function shall determine whether the address specified by <i>addr</i> is located within the address range occupied by a mapped object. The mapped objects examined shall include any executable object files that have previously been loaded by a call to <i>dlopen</i> () and for which <i>dlclose</i> () has not subsequently been called, and any shared library files that were loaded as dependencies of the executable file from which the current process image was loaded; they may also include any executable object files that have previously been loaded by a call to <i>dlopen</i> () and for which <i>dlclose</i> () has subsequently been called, the executable file from which the current process image was loaded, and implementation-defined additional mapped objects (for example, all regular files mapped using <i>mmap</i> () might be included). If the specified address is within the mapped address range of one of these mapped objects and the object contains a symbol table, the symbol table shall be searched for a symbol (a function identifier or a data object identifier) that has the largest address less than or equal to the specified address.	+ + + + + + + + + + + + + + + + + + +
24826 24827	If the address specified by <i>addr</i> is within the mapped address range of one of the examined mapped objects, the structure pointed to by <i>dlip</i> shall be populated as follows:	+ +
24828 24829 24830	• The value of the <i>dli_fname</i> member shall be set to point to the pathname of the mapped object. (This might no longer resolve to the file that was mapped, for example if it was a link that has subsequently been removed or renamed.)	+ + +
24831 24832	• The value of the <i>dli_fbase</i> member shall be set to the base of the address range occupied by the mapped object.	+ +
24833 24834 24835	• The value of the <i>dli_sname</i> member shall be set to point to the name of the symbol that has the largest address less than or equal to the specified address, or to a null pointer if no such symbol was found.	+ + +
24836 24837 24838 24839 24840	• If <i>dli_sname</i> is set to a null pointer, the value of the <i>dli_saddr</i> member shall also be set to a null pointer. Otherwise, if <i>dli_sname</i> names a function identifier, <i>dli_saddr</i> shall be set to the address of the function converted from type pointer to function to type pointer to <b>void</b> ; otherwise, <i>dli_saddr</i> shall be set to the address of the data object named by <i>dli_sname</i> converted from a pointer to the type of the data object to a pointer to <b>void</b> .	+ + + +
24841 24842 24843 24844 24845	RETURN VALUE Upon successful completion, a non-zero value shall be returned. If the specified address is not located within the address range occupied by an examined mapped object, or if an error occurs, zero shall be returned. More detailed diagnostic information shall be available through <i>dlerror</i> (). ERRORS	+ + + +
41010		

24846 No errors are defined.
## dladdr()

24847	EXAMPLES	+
24848	None.	+
24849	APPLICATION USAGE	+
24850	The Dl_info_t members may point to addresses within the mapped object. These pointers can	+
24851	become invalid if the object is unmapped (for example, loaded executable objects may be	+
24852	unloaded by <i>dlclose()</i> ).	+
24853 24854 24855 24856	If <i>dli_sname</i> names a function identifier, the value of <i>dli_saddr</i> can be converted back to type pointer to function using a cast in the manner shown in the <i>dlsym()</i> EXAMPLES section. Note that this conversion is not defined by the ISO C standard. This standard requires this conversion to work correctly on conforming implementations.	+ + + +
24857	RATIONALE	+
24858	None.	+
24859	FUTURE DIRECTIONS	+
24860	None.	+
24861	SEE ALSO	+
24862	dlclose(), dlerror(), dlopen(), dlsym()	+
24863	XBD <dlfcn.h></dlfcn.h>	+
24864 24865 24866	CHANGE HISTORY First released in Issue 8.	+ + +

738

25353	assert(xsubi[1] == 10728);	
25354	assert(xsubi[2] == 27921);	
25355	assert(nrand48(xsubi) == 754104482);	
25356	assert(xsubi[0] == 6828);	
25357	assert(xsubi[1] == 28997);	
25358	assert(xsub1[2] == 23013);	
25359	assert(nrand48(xsub1) == 609453945);	
25360	assert(xsub1[0] == 58183);	
25361	assert(xsub1[1] == 3826);	
25362	assert(xsub1[2] == 18599);	
25363	assert(mrand40(xsubr) = 1070044500);	
25364	assert (xsubi $[1] = 44304$ );	
25365	assert(xsubi[1] == 57331).	
25367	assert(xsub1[2] = -37331), assert(nrand48(xsub1) == 2114923686).	
25368	assert(minu40(x30b)) = 2114925000),	
25369	assert (xsubi $[1] = 22861$ );	
25370	assert (xsubi[2] == $64542$ ):	
25371	}	
25372		
20072	,	
25373	APPLICATION USAGE	
25374	These functions should be avoided whenever non-trivial requirements (including safety) have to	
25375	be fulfilled, unless seeded using getentropy().	
25376	RATIONALE	
25377	None.	
25378	FUTURE DIRECTIONS	
25379	None.	
25380	SEE ALSO	
25381	getentropy(), initstate(), rand()	+
25282	XBD <stdlib h=""></stdlib>	
23362		
25383	CHANGE HISTORY	
25384	First released in Issue 1. Derived from Issue 1 of the SVID.	
25385	Issue 5	
25386	A note indicating that the $drand48()$ $lrand48()$ and $mrand48()$ functions need not be reentrant is	
25387	added to the DESCRIPTION	
20007		
25388	Issue 6	
25389	The normative text is updated to avoid use of the term "must" for application requirements.	
25390	Issue 7	
25391	Austin Group Interpretation 1003.1-2001 #156 is applied.	
25392	POSIX.1-2008, Technical Corrigendum 2, XSH/TC2-2008/0083 [743] is applied.	
25393	Issue 8	
25394	Austin Group Defect 1107 is applied, clarifying how the return value is calculated from $X_i$ for	
25395	each function.	+
25207	Austin Group Defect 1134 is applied adding <i>activityonu()</i>	I
25396	Austin Group Delect 1154 is applied, adding getentropy().	I

<sup>30202</sup> possible for the system to conform to the intent of this volume of POSIX.1-202x.

The [EAGAIN] error exists to warn applications that such a condition might occur. Whether it occurs or not is not in any practical sense under the control of the application because the condition is usually a consequence of the user's use of the system, not of the application's code. Thus, no application can or should rely upon its occurrence under any circumstances, nor should the exact semantics of what concept of ``user'' is used be of concern to the application developer. Validation writers should be cognizant of this limitation.

30209There are two reasons why POSIX programmers call *fork()*. One reason is to create a new thread30210of control within the same program (which was originally only possible in POSIX by creating a30211new process); the other is to create a new process running a different program. In the latter case,30212the call to *fork()* is soon followed by a call to one of the *exec* functions.

The general problem with making fork() work in a multi-threaded world is what to do with all 30213 30214 of the threads. There are two alternatives. One is to copy all of the threads into the new process. This causes the programmer or implementation to deal with threads that are suspended on 30215 system calls or that might be about to execute system calls that should not be executed in the 30216 new process. The other alternative is to copy only the thread that calls fork(). This creates the 30217 difficulty that the state of process-local resources is usually held in process memory. If a thread 30218 that is not calling fork() holds a resource, that resource is never released in the child process 30219 30220 because the thread whose job it is to release the resource does not exist in the child process.

30221When a programmer is writing a multi-threaded program, the first described use of fork(),30222creating new threads in the same program, is provided by the  $pthread\_create()$  function. The30223fork() function is thus used only to run new programs, and the effects of calling functions that30224require certain resources between the call to fork() and the call to an *exec* function are undefined.

The addition of the *forkall()* function to the standard was considered and rejected. The *forkall()* 30225 function lets all the threads in the parent be duplicated in the child. This essentially duplicates 30226 the state of the parent in the child. This allows threads in the child to continue processing and 30227 allows locks and the state to be preserved without explicit *pthread\_atfork()* code. The calling 30228 30229 process has to ensure that the threads processing state that is shared between the parent and child (that is, file descriptors or MAP\_SHARED memory) behaves properly after forkall(). For 30230 example, if a thread is reading a file descriptor in the parent when *forkall()* is called, then two 30231 threads (one in the parent and one in the child) are reading the file descriptor after the *forkall()*. 30232 If this is not desired behavior, the parent process has to synchronize with such threads before 30233 calling *forkall()*. 30234

When *forkall*() is called, threads, other than the calling thread, that are in functions that can return with an [EINTR] error may have those functions return [EINTR] if the implementation cannot ensure that the function behaves correctly in the parent and child. In particular, *pthread\_cond\_clockwait*(), *pthread\_cond\_timedwait*(), and *pthread\_cond\_wait*() need to return in order to ensure that the condition has not changed. These functions can be awakened by a spurious condition wakeup rather than returning [EINTR].

#### 30241 FUTURE DIRECTIONS

30242 None.

# 30243 SEE ALSO 30244 alarm(), exec, fcntl(), pthread\_atfork(), semop(), signal(), times() 20245 XBD Section 4.13 (on page 91) < sys/types b> < unistd b>

30245 XBD Section 4.13 (on page 91), <**sys/types.h**>, <**unistd.h**>

free()

31485	NAME		
31486		free — free allocated memory	
31487	7 SYNOPSIS		
31488		<pre>#include <stdlib.h></stdlib.h></pre>	
31489		<pre>void free(void *ptr);</pre>	
31490	DESCR	IPTION	
31491	CX	The functionality described on this reference page is aligned with the ISOC standard. Any	
31492 31493		conflict between the requirements described here and the ISO C standard is unintentional. This volume of POSIX.1-202x defers to the ISO C standard.	
31494		The <i>free()</i> function shall cause the space pointed to by <i>ptr</i> to be deallocated; that is, made	
31495		available for further allocation. If <i>ptr</i> is a null pointer, no action shall occur. Otherwise, if the	
31496	<i></i>	argument does not match a pointer earlier returned by a function in POSIX.1-202x that allocates	ı
31497	CX	reallocarrau() the behavior is undefined	I
31499		Any use of a pointer that refers to freed space results in undefined behavior.	
31500	CX	The <i>free()</i> function shall not modify <i>errno</i> if <i>ptr</i> is a null pointer or a pointer previously returned	
31501		as if by <i>malloc()</i> and not yet deallocated.	
31502	RETUR	N VALUE	
31503		The <i>free</i> () function shall not return a value.	
31504	ERROR	S	
31505		No errors are defined.	
31506	EXAMP	PLES	
31506 31507	EXAMP	None.	
31506 31507 31508	EXAMP	PLES None. CATION USAGE	
31506 31507 31508 31509	EXAMP APPLIC	PLES None. CATION USAGE There is now no requirement for the implementation to support the inclusion of <b><malloc.h></malloc.h></b> .	
<ul> <li>31506</li> <li>31507</li> <li>31508</li> <li>31509</li> <li>31510</li> </ul>	EXAMP APPLIC	PLES         None.         CATION USAGE         There is now no requirement for the implementation to support the inclusion of <malloc.h>.         Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup</malloc.h>	
31506 31507 31508 31509 31510 31511	EXAMP	PLES         None.         CATION USAGE         There is now no requirement for the implementation to support the inclusion of <malloc.h>.         Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:</malloc.h>	
31506 31507 31508 31509 31510 31511 31512	EXAMF	None. CATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code: // buf was obtained by malloc(buflen)</malloc.h>	
31506 31507 31508 31509 31510 31511 31512 31513	EXAMF	<pre>None. CATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the free() function does not modify errno for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);</malloc.h></pre>	
31506 31507 31508 31509 31510 31511 31512 31513 31514	EXAMF	<pre>PLES None. PLES None. PLES CATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);     if (ret &lt; 0) {         for (here for the form)     } </malloc.h></pre>	
31506 31507 31508 31509 31510 31511 31512 31513 31514 31515 31516	EXAMF	<pre>None. XATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the free() function does not modify errno for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);     if (ret &lt; 0) {         free(buf);         return ret;     } }</malloc.h></pre>	
31506 31507 31508 31509 31510 31511 31512 31513 31514 31515 31516 31517	EXAMF	<pre>PLES None. CATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:</malloc.h></pre>	
31506 31507 31508 31509 31510 31511 31512 31513 31514 31515 31516 31517	EXAMF	<pre>PLES None. PLES None. PLES None. PLES There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);     if (ret &lt; 0) {         free(buf);         return ret;     } However earlier versions of this standard did not require this, and the same example had to be </malloc.h></pre>	
31506 31507 31508 31509 31510 31511 31512 31513 31514 31515 31516 31517 31518 31519	EXAMF	<pre>None. EXATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:</malloc.h></pre>	
31506 31507 31508 31509 31510 31511 31512 31513 31514 31515 31516 31517 31518 31519	EXAMP	<pre>None. CATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:</malloc.h></pre>	
31506 31507 31508 31509 31510 31511 31512 31513 31514 31515 31516 31517 31518 31519 31520 31521	EXAMF	<pre>None. EXATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the free() function does not modify errno for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);     if (ret &lt; 0) {         free(buf);         return ret;     } However, earlier versions of this standard did not require this, and the same example had to be written as:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);</malloc.h></pre>	
31506 31507 31508 31510 31511 31512 31513 31514 31515 31516 31517 31518 31519 31520 31521 31522	EXAMP	<pre>None. EXATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the free() function does not modify errno for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);     if (ret &lt; 0) {         free(buf);         return ret;     } However, earlier versions of this standard did not require this, and the same example had to be written as:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);     if (ret &lt; 0) {         free(buf);         return ret;     } </malloc.h></pre>	
31506 31507 31508 31509 31510 31511 31512 31513 31514 31516 31516 31517 31518 31519 31520 31521 31522 31523	EXAMP	<pre>None. EXATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);     if (ret &lt; 0) {         free(buf);         return ret;     } However, earlier versions of this standard did not require this, and the same example had to be written as:     // buf was obtained by malloc(buflen)     ret = write(fd, buf, buflen);     if (ret &lt; 0) {         int save = errno;     } }</malloc.h></pre>	
31506 31507 31508 31509 31510 31512 31512 31513 31514 31515 31516 31517 31518 31519 31520 31521 31522 31523 31523	EXAMP	<pre>LES None. CATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:</malloc.h></pre>	
31506 31507 31508 31510 31511 31512 31513 31514 31515 31516 31517 31518 31519 31520 31521 31522 31523 31524 31525	EXAMP	<pre>None. XATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free(</i>) function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:</malloc.h></pre>	
31506 31507 31508 31509 31510 31512 31512 31513 31514 31515 31516 31517 31518 31519 31520 31521 31522 31522 31523 31524 31525 31526 31527	EXAMP	<pre>LES None. ZATION USAGE There is now no requirement for the implementation to support the inclusion of <malloc.h>. Because the <i>free()</i> function does not modify <i>errno</i> for valid pointers, it is safe to use it in cleanup code without corrupting earlier errors, such as in this example code:</malloc.h></pre>	

34913 34914	NAME getentropy — fill a buffer with random bytes
34915 34916	SYNOPSIS #include <unistd.h></unistd.h>
34917	<pre>int getentropy(void *buffer, size_t length);</pre>
34918 34919 34920 34921 34922 34923	DESCRIPTION The <i>getentropy</i> () function shall write <i>length</i> bytes of data starting at the location pointed to by <i>buffer</i> . The output shall be unpredictable high quality random data, generated by a cryptographically secure pseudo-random number generator. The maximum permitted value for the <i>length</i> argument is given by the {GETENTROPY_MAX} symbolic constant defined in <b><limits.h></limits.h></b> .
34924	A successful call to <i>getentropy</i> () shall always provide the requested number of bytes of entropy.
34925 34926 34927	<b>RETURN VALUE</b> Upon successful completion, getentropy() shall return 0; otherwise, -1 shall be returned and errno set to indicate the error.       Image: set of the error
34928 34929	ERRORS The getentropy() function shall fail if:
34930	[EINVAL] The value of <i>length</i> is greater than {GETENTROPY_MAX}.
34931	The <i>getentropy</i> () function may fail if:
34932	[ENOSYS] The system does not provide the necessary source of entropy.
34933 34934	EXAMPLES None.
34935 34936 34937	APPLICATION USAGE The intended use of this function is to create a seed for other pseudo-random number generators.
34938 34939	RATIONALE                 The getentropy() function is not a cancellation point. (See Section 2.9.5.2 (on page 504).)
34940 34941	FUTURE DIRECTIONS None.
34942 34943	SEE ALSO drand48(), initstate(), rand()
34944	XBD <limits.h>, <unistd.h></unistd.h></limits.h>
34945 34946 34947	CHANGE HISTORY First released in Issue 8.

## getlocalename\_l()

	NAME
35541	getlocalename_l — get a locale name from a locale object
35542	SYNOPSIS
35544	<pre>CX #include <locale.h></locale.h></pre>
35545	<pre>const char *getlocalename_l(int category, locale_t locobj);</pre>
35546	
35547	DESCRIPTION
35548	The <i>getlocalename_l()</i> function shall return the locale name for the given locale category of the
35549	locale object locobj, or of the global locale if locobj is the special locale object
35550	LC_GLOBAL_LOCALE.
35551	The <i>category</i> argument specifies the locale category to be queried. If the value is LC_ALL or is
35552	not a supported locale category value (see <i>setlocale()</i> ), <i>getlocalename_l()</i> shall fail.
35553	The behavior is undefined if the <i>locobj</i> argument is neither the special locale object
35554	LC_GLOBAL_LOCALE nor a valid locale object handle.
35555	RETURN VALUE
35556	Upon successful completion, <i>getlocalename_l()</i> shall return a pointer to a string containing the
35557	locale name; otherwise, a null pointer shall be returned.
35558	If locobj is LC_GLOBAL_LOCALE, the returned string pointer might be invalidated or the string
35559	content might be overwritten by a subsequent call in the same thread to <i>getlocalename_l()</i> with
35560	LC_GLOBAL_LOCALE; the returned string pointer might also be invalidated if the calling
35561	the locale object <i>locabi</i> is used in a call to <i>freelocale()</i> or as the <i>base</i> argument in a successful call to
35563	newlocale().
35564	FRRORS
35565	No errors are defined.
35566	EXAMPLES
00000	
35567	Determining the locale name for a category of the current locale
35568	The following example shows how to obtain the locale name for the LC_NUMERIC category of
35569	the current thread-local locale, or of the global locale if no thread-local locale is in use.
35570	<pre>#include <locale.h></locale.h></pre>
35571	
35572	const char *name;
35573	locale_t loc = uselocale(NULL);
55574	Thanke - gettocatemane_t(lc_NOMERIC, toc),
35575	APPLICATION USAGE
35576	
35577	RATIONALE
35578	$\Gamma$ ISTORICAL VERSIONS OF getuccalename_I() and not handle the special locale object   IC CLOBAL LOCALE requiring that applications used setlocale(category NULL) to guary the
35580	global locale if uselocale(NULL) returned LC GLOBAL LOCALE. However, since setlocale() is
35581	not required to be thread-safe (even when the only concurrent calls are ones that query the
35582	locale), this method was problematic for multi-threaded processes. This standard requires that
35583	getlocalename_l(category, LC_GLOBAL_LOCALE) queries the global locale in a thread-safe
35584	manner, for example by returning a pointer to a thread-local internal buffer instead of a process-
35585	
35568 35570 35571 35572 35573 35574 35575 35576 35576 35577 35578 35578 35579 35580 35581 35582 35583 35584 35584	The following example shows how to obtain the locale name for the LC_NUMERIC category of the current thread-local locale, or of the global locale if no thread-local locale is in use. #include <locale.h>  const char *name; locale_t loc = uselocale(NULL); name = getlocalename_l(LC_NUMERIC, loc); APPLICATION USAGE None. RATIONALE Historical versions of getlocalename_l() did not handle the special locale object LC_GLOBAL_LOCALE, requiring that applications used setlocale(category, NULL) to query the global locale if uselocale(NULL) returned LC_GLOBAL_LOCALE. However, since setlocale() is not required to be thread-safe (even when the only concurrent calls are ones that query the locale), this method was problematic for multi-threaded processes. This standard requires that getlocalename_l(category, LC_GLOBAL_LOCALE) queries the global locale in a thread-safe manner, for example by returning a pointer to a thread-local internal buffer instead of a process- wide internal buffer.</locale.h>

## 35586FUTURE DIRECTIONS35587None.

35588 SEE ALSO
 35589 freelocale(), newlocale(), setlocale(), uselocale()
 35590 XBD Chapter 7 (on page 113), <locale.h>
 35591 CHANGE HISTORY

35592 First released in Issue 8.

35593

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38372 38373	ERRORS No errors are defined.	
38374 38375	EXAMPLES None.	
38376 38377	APPLICATION USAGE After initialization, a state array can be restarted at a different point in one of two ways:	
38378 38379	1. The <i>initstate()</i> function can be used, with the desired seed, state array, and size of the array.	
38380 38381 38382	2. The <i>setstate()</i> function, with the desired state, can be used, followed by <i>srandom()</i> with the desired seed. The advantage of using both of these functions is that the size of the state array does not have to be saved once it is initialized.	
38383 38384	Although some implementations of <i>random()</i> have written messages to standard error, such implementations do not conform to POSIX.1-202x.	
38385	Issue 5 restored the historical behavior of this function.	
38386 38387	Threaded applications should use <i>erand48(), nrand48()</i> , or <i>jrand48()</i> instead of <i>random()</i> when an independent random number sequence in multiple threads is required.	
38388 38389	These functions should be avoided whenever non-trivial requirements (including safety) have to be fulfilled, unless seeded using <i>getentropy</i> ().	
38390 38391	RATIONALE None.	
38392 38393	FUTURE DIRECTIONS None.	
38394	SEE ALSO	
38395	drand48(), getentropy(), rand()	+
38396	XBD <stdlib.h></stdlib.h>	
38397 38398	CHANGE HISTORY First released in Issue 4, Version 2.	
38399	Issue 5	
38400	Moved from X/OPEN UNIX extension to BASE.	
38401 38402 38403 38404	In the DESCRIPTION, the phrase "values smaller than 8" is replaced with "values greater than or equal to 8, or less than $32''$ , " <i>size</i> <8" is replaced with " $8 \le size < 32''$ , and a new first paragraph is added to the RETURN VALUE section. A note is added to the APPLICATION USAGE indicating that these changes restore the historical behavior of the function.	
38405	Issue 6	
38406	In the DESCRIPTION, duplicate text "For values greater than or equal to 8" is removed.	
38407 38408	IEEE Std 1003.1-2001/Cor 1-2002, item XSH/TC1/D6/30 is applied, removing $rand_r()$ from the list of suggested functions in the APPLICATION USAGE section.	
38409 38410	<b>Issue 7</b> The type of the first argument to <i>setstate(</i> ) is changed from <b>const char</b> * to <b>char</b> *	
38/11	POSIX 1-2008 Technical Corrigendum 2 XSH/TC2-2008/0179 [743] is applied	+
50411	1 001/11 2000, reclinical configentiant 2, 7011/102-2000/01/2 [745] 15 applica.	'

## memmem()

43075 43076	NAME memmem — find a byte subsequence in a byte sequence
43077	SYNOPSIS
43078	CX #include <string.h></string.h>
43079 43080 43081	<pre>void *memmem(const void *haystack, size_t haystacklen,</pre>
43082	DESCRIPTION
43083 43084	The <i>memmem()</i> function shall locate the first occurrence of byte sequence <i>needle</i> of length <i>needlelen</i> in byte sequence <i>haystack</i> of length <i>haystacklen</i> .
43085	RETURN VALUE
43086 43087	Upon successful completion, <i>memmem()</i> shall return a pointer to the the first byte of the located byte sequence in <i>haystack</i> , or a null pointer if the byte sequence is not found.
43088	If <i>needlelen</i> is zero, the function shall return <i>haystack</i> .
43089	If <i>haystacklen</i> is less than <i>needlelen</i> , the function shall return a null pointer.
43090	ERRORS
43091	No errors are defined.
43092	EXAMPLES
43093	None.
43094 43095	APPLICATION USAGE None.
43096	RATIONALE
43097 43098	This function is similar to <i>strstr()</i> , except that NUL bytes may be included in either <i>needle</i> or <i>haystack</i> .
43099	<b>FUTURE DIRECTIONS</b>
43100	None.
43101	SEE ALSO
43102	memchr(), strstr()
43103	XBD <string.h></string.h>
43104	CHANGE HISTORY
43105	First released in Issue 8.
43106	

poll()

47102	NAME		
47103	poll, ppoll — inp	ut/output multiplexing	+
47104	SYNOPSIS		
47105	<pre>#include <pol< pre=""></pol<></pre>	ll.h>	
47106 47107 47108 47109	int poll(stru int ppoll(stru const stru const sig	<pre>act pollfd fds[], nfds_t nfds, int timeout); cuct pollfd fds[], nfds_t nfds, cuct timespec *restrict timeout, gset_t *restrict sigmask);</pre>	+ + +
47110	DESCRIPTION		
47111 47112 47113 47114 47115 47116	The <i>ppoll</i> () functover a set of file examine the give structures in the descriptors on we occurred.	tion provides applications with a mechanism for multiplexing input/output e descriptors. For each member of the array pointed to by <i>fds</i> , <i>ppoll</i> () shall en file descriptor for the event(s) specified in <i>events</i> . The number of <b>pollfd</b> <i>fds</i> array is specified by <i>nfds</i> . The <i>ppoll</i> () function shall identify those file which an application can read or write data, or on which certain events have	 
47117	The <i>poll</i> () function	on shall be equivalent to the <i>ppoll</i> () function, except as follows:	+
47118 47119 47120 47121	For the <i>poll</i> <b>int</b> , where     nanosecond     shall be equ	() function, the timeout period is given in milliseconds in an argument of type as for the <i>ppoll</i> () function the timeout period is given in seconds and ds via an argument of type pointer to <b>struct timespec</b> . A <i>timeout</i> of $-1$ for <i>poll</i> () uvalent to passing a null pointer for the <i>timeout</i> for <i>ppoll</i> ().	+ + +
47122 47123	<ul> <li>The <i>poll()</i> f a null point</li> </ul>	unction has no <i>sigmask</i> argument; it shall behave as <i>ppoll()</i> does when <i>sigmask</i> is ter.	+ +
47124 47125 47126 47127 47128	The <i>fds</i> argument file descriptor. It interest. The arr descriptor and <i>e</i> following event f	t specifies the file descriptors to be examined and the events of interest for each t is a pointer to an array with one member for each open file descriptor of ray's members are <b>pollfd</b> structures within which <i>fd</i> specifies an open file <i>events</i> and <i>revents</i> are bitmasks constructed by OR'ing a combination of the flags:	
47129	POLLIN	Data other than high-priority data may be read without blocking.	
47130	POLLRDNORM	Normal data may be read without blocking.	
47131	POLLRDBAND	Priority data may be read without blocking.	
47132	POLLPRI	High-priority data may be read without blocking.	
47133	POLLOUT	Normal data may be written without blocking.	
47134	POLLWRNORM	Equivalent to POLLOUT.	
47135	POLLWRBAND	Priority data may be written.	
47136 47137	POLLERR	An error has occurred on the device or stream. This flag is only valid in the <i>revents</i> bitmask; it shall be ignored in the <i>events</i> member.	
47138 47139 47140 47141 47142 47143 47144 47145	POLLHUP	A device has been disconnected, or a pipe or FIFO has been closed by the last process that had it open for writing. Once set, the hangup state of a FIFO shall persist until some process opens the FIFO for writing or until all read-only file descriptors for the FIFO are closed. This event and POLLOUT are mutually-exclusive; a stream can never be writable if a hangup has occurred. However, this event and POLLIN, POLLRDNORM, POLLRDBAND, or POLLPRI are not mutually-exclusive. This flag is only valid in the <i>revents</i> bitmask; it shall be ignored in the <i>events</i> member.	

- POLLNVAL The specified *fd* value is invalid. This flag is only valid in the *revents* member; 47146 it shall ignored in the *events* member. 47147 47148 The significance and semantics of normal, priority, and high-priority data are file and devicespecific. 47149 If the value of *fd* is less than 0, *events* shall be ignored, and *revents* shall be set to 0 in that entry on 47150 return from *poll()* or *ppoll()*. 47151 In each **pollfd** structure, *poll()* or *ppoll()* shall clear the *revents* member, except that where the + 47152 application requested a report on a condition by setting one of the bits of *events* listed above, 47153 *poll()* or *ppoll()* shall set the corresponding bit in *revents* if the requested condition is true. In 47154 + addition, *poll()* or *ppoll()* shall set the POLLHUP, POLLERR, and POLLNVAL flag in *revents* if 47155 the condition is true, even if the application did not set the corresponding bit in *events*. 47156 The *timeout* argument controls how long the *poll()* or *ppoll()* function shall wait before timing 47157 out. If the *timeout* argument is positive for *poll()* or not a null pointer for *ppoll()*, it specifies a 47158 maximum interval to wait for the poll to complete. If the specified time interval expires without 47159 any of the defined events having occurred, the function shall return. If the *timeout* argument is 47160 47161 -1 for *poll(*) or a null pointer for *ppoll(*), then the call shall block indefinitely until at least one descriptor meets the specified criteria or until the call is interrupted. To effect a poll, the 47162 47163 application shall ensure that the *timeout* argument for *poll(*) is 0, or for *ppoll(*) is not a null pointer and points to a zero-valued timespec structure. 47164 Implementations may place limitations on the maximum timeout interval supported. All 47165 implementations shall support a maximum timeout interval of at least 31 days for *ppoll()*. If the 47166 timeout argument specifies a timeout interval greater than the implementation-defined 47167 47168 maximum value, the maximum value shall be used as the actual timeout value. Implementations may also place limitations on the granularity of timeout intervals. If the requested timeout 47169 interval requires a finer granularity than the implementation supports, the actual timeout 47170 interval shall be rounded up to the next supported value. 47171 The *poll()* and *ppoll()* functions shall not be affected by the O\_NONBLOCK flag. 47172 The *poll()* and *ppoll()* functions shall support regular files, terminal and pseudo-terminal 47173 devices, FIFOs, pipes, and sockets. The behavior of *poll()* and *ppoll()* on elements of *fds* that refer 47174 + to other types of file is unspecified. 47175 Regular files shall always poll TRUE for reading and writing. 47176 47177 A file descriptor for a socket that is listening for connections shall indicate that it is ready for reading, once connections are available. A file descriptor for a socket that is connecting 47178 asynchronously shall indicate that it is ready for writing, once a connection has been established. 47179 Provided the application does not perform any action that results in unspecified or undefined 47180 behavior, the value of the *fd* and *events* members of each element of *fds* shall not be modified by 47181 poll() or ppoll(). 47182 If *sigmask* is not a null pointer, the *ppoll()* function shall replace the signal mask of the caller by 47183 the set of signals pointed to by sigmask before examining the descriptors, and shall restore the 47184 47185 signal mask of the calling thread before returning. If a signal is unmasked as a result of the signal mask being altered by *ppoll()*, and a signal-catching function is called for that signal 47186 47187 during the execution of the *ppoll()* function, and SA\_RESTART is clear for the interrupting signal, then 47188 • If none of the defined events have occurred on any selected file descriptor, *ppoll()* shall 47189 immediately fail with the [EINTR] error after the signal-catching function returns. 47190
  - 1414

```
poll()
```

47191 47192 47193 47194	• If one or more of the defined events have occurred, it is unspecified whether <i>ppoll()</i> behaves the same as if none of the events had occurred (failing with [EINTR] as above) or behaves the same as if it was not interrupted (returning the total number of <b>pollfd</b> structures that have selected events).	   
47195 47196 47197	If a thread is canceled during a <i>ppoll()</i> call, it is unspecified whether the signal mask in effect when executing the registered cleanup functions is the original signal mask or the signal mask installed as part of the <i>ppoll()</i> call.	   
47198 47199 47200 47201 47202 47203	<b>RETURN VALUE</b> Upon successful completion, a non-negative value shall be returned. A positive value shall indicate the total number of <b>pollfd</b> structures that have selected events (that is, those for which the <i>revents</i> member is non-zero). A value of 0 shall indicate that the call timed out and no file descriptors have been selected. Upon failure, -1 shall be returned and <i>errno</i> set to indicate the error.	
47204	ERRORS	I
47205 47206 47207	[EAGAIN] The allocation of internal data structures failed but a subsequent request may succeed.	I
47208	[EINTR] A signal was caught during <i>poll()</i> or <i>ppoll()</i> .	
47209	[EINVAL] The <i>nfds</i> argument is greater than {OPEN_MAX}.	
47210	The <i>ppoll</i> () function shall fail if:	+
47211	[EINVAL] An invalid timeout interval was specified.	+
47212 47213	EXAMPLES None.	
47214 47215	APPLICATION USAGE Other than the difference in the precision of the requested timeout, the following <i>ppoll()</i> call:	
47216	<pre>ready = ppoll(&amp;fds, nfds, tmo_p, &amp;sigmask);</pre>	
47217	is equivalent to atomically executing the following calls:	
47218 47219	<pre>sigset_t origmask; int timeout;</pre>	 
47220 47221 47222 47223 47224	<pre>timeout = (tmo_p == NULL) ? -1 :     (tmo_p-&gt;tv_sec * 1000 + tmo_p-&gt;tv_nsec / 1000000); pthread_sigmask(SIG_SETMASK, &amp;sigmask, &amp;origmask); ready = poll(&amp;fds, nfds, timeout); pthread_sigmask(SIG_SETMASK, &amp;origmask, NULL);</pre>	
47225 47226 47227 47228	<b>RATIONALE</b> The POLLHUP event does not occur for FIFOs just because the FIFO is not open for writing. It only occurs when the FIFO is closed by the last writer and persists until some process opens the FIFO for writing or until all read-only file descriptors for the FIFO are closed.	+
47229 47230	Code which wants to avoid the ambiguity of the signal mask for thread cancellation handlers can install an additional cancellation handler which resets the signal mask to the expected value:	+ +
47231	<pre>void cleanup(void *arg) </pre>	+
47232 47233	t sigset_t *ss = (sigset_t *) arg;	+

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1415

47234			<pre>pthread_sigmask(SIG_SETMASK, ss, NULL);</pre>	+
47235		}		+
47236		int	call_ppoll(struct pollid ids[], nids_t nids,	+
47237			const struct timespec restrict timeout,	+
47239		{	const sigset_t restrict signask)	+
47240		C C	sigset_t oldmask;	+
47241			int result;	+
47242			<pre>pthread_sigmask(SIG_SETMASK, NULL, &amp;oldmask);</pre>	+
47243			<pre>pthread_cleanup_push(cleanup, &amp;oldmask);</pre>	+
47244			<pre>result = ppoll(fds, nfds, timeout, sigmask);</pre>	+
47245			<pre>pthread_cleanup_pop(0);</pre>	+
47246		ı	return result;	+
47247	FUTUR	<sup>,</sup> E DIF	RECTIONS	+
47249		Non	e.	
47250	SEE AL	SO		
47251		psele	ct(), read(), write()	
47252		XBD	<poll.h></poll.h>	
47253	CHANC	GE HI	STORY	
47254		First	released in Issue 4, Version 2.	
47255	Issue 5			
47256		Mov	red from X/OPEN UNIX extension to BASE.	
47257		The	description of POLLWRBAND is updated.	
47258	Issue 6			
47259		Text	referring to sockets is added to the DESCRIPTION.	
47260		Func	ctionality relating to the XSI STREAMS Option Group is marked.	
47261 47262		The POL	Open Group Corrigendum U055/3 is applied, updating the DESCRIPTION of LWRBAND.	
47263 47264		IEEE the E	Std 1003.1-2001/Cor 2-2004, item XSH/TC2/D6/66 is applied, correcting the spacing in EXAMPLES section.	
47265	Issue 7			
47266		Aust	tin Group Interpretation 1003.1-2001 #209 is applied, clarifying the POLLHUP event.	
47267		The	poll() function is moved from the XSI option to the Base.	
47268		Func	ctionality relating to the XSI STREAMS option is marked obsolescent.	
47269 47270		POS [683]	IX.1-2008, Technical Corrigendum 2, XSH/TC2-2008/0249 [623] and XSH/TC2-2008/0250 ] are applied.	
47271	Issue 8			
47272		Aust	tin Group Defect 1263 is applied, adding <i>ppoll</i> ().	+
47273		Aust	tin Group Defect 1330 is applied, removing obsolescent interfaces.	

## posix\_getdents()

	NAME	
47780		posix_getdents — read directory entries
47782	SYNOP	SIS
47783		<pre>#include <dirent.h></dirent.h></pre>
47784		<pre>ssize_t posix_getdents(int fildes, void *buf, size_t nbyte, int flags);</pre>
47785	DESCRI	IPTION
47786		The <i>posix_getdents()</i> function shall attempt to read directory entries from the directory associated
47787		with the open file descriptor <i>fildes</i> and shall place information about the directory entries and the
47788		files they refer to in <b>posix_dent</b> structures in the buffer pointed to by <i>buf</i> , up to a maximum of
47789		number that will fit in white but a but shall be at least one if white is greater than the
47790 47791		<b>posix dent</b> structure plus {NAME MAX} and <i>fildes</i> is not currently at end-of-file.
47702		The application shall ensure that <i>huf</i> is aligned suitably to point to a <b>posix dent</b> structure. The
47792		alignment needed shall not be more restrictive than the alignment provided by <i>malloc()</i> . Strictly
47794		conforming applications shall ensure that the value of <i>flags</i> is zero: other applications can set it to
47795		a value constructed by a bitwise-inclusive OR of implementation-defined bitwise-distinct flag
47796		values.
47797		Each <b>posix_dent</b> structure returned in <i>buf</i> shall be located at an address that satisfies the
47798		implementation's alignment requirements for the <b>posix_dent</b> structure and shall be populated
47799		as follows:
47800		• The value of the <i>d_ino</i> member shall be set to the file serial number of the file named by the
47801		<i>d_name</i> member.
47802		• The value of the <i>d_reclen</i> member shall be set to the number of bytes occupied by this entry
47803		in <i>buf</i> , including any padding bytes needed before the next entry, if any. If this is the last
47804 47805		entry in <i>buf, a_recten</i> shall include any padding bytes needed to make the address of this entry plus <i>d_recten</i> bytes satisfy the alignment requirements for the <b>posix_dent</b> structure
47806		• The value of the <i>d</i> ture member shall be set to indicate the file ture of the named file if the
47807		file type can be determined without needing to use the file serial number to obtain the
47808		file's metadata; otherwise it may be set to DT UNKNOWN. If the file type is determined
47809		and it is one of the file types defined in this standard, the value of <i>d</i> type shall be DT BLK,
47810		DT_CHR, DT_DIR, DT_FIFO, DT_LNK, DT_REG, DT_SOCK, DT_MQ, DT_SEM,
47811	TYM	DT_SHM, or DT_TMO (see <dirent.h>). If it is determined but is not a standard file type,</dirent.h>
47812		the value of $d_type$ shall not equal any of those listed here.
47813		• The <i>d_name</i> member shall be a filename string, and (if not dot or dot-dot) shall contain the
47814		same byte sequence as the last pathname component of the string used to create the
47815		directory entry, plus the terminating NUL byte.
47816		If the <i>d_name</i> member names a symbolic link, the values of the <i>d_ino</i> and <i>d_type</i> members shall
47817		be set to the values for the symbolic link itself.
47818		The <i>posix_getdents()</i> function shall start reading at the current file offset in the open file
47819		description associated with <i>fildes</i> . On successful return, the file offset shall be incremented to
47820		point to the directory entry immediately following the last entry whose information was
47821		returned in <i>vuf</i> , or to point to end-of-file if there are no more directory entries. On failure, the
47822		on the open file description associated with fildes. The behavior is unspecified if least() is used
47824		to set the file offset to a value other than zero or a value returned by a previous call to $leak()$ on
47825		the same open file description.
010		

The *posix\_getdents()* function shall not return directory entries containing empty names. If 47826 entries for dot or dot-dot exist, a sequence of calls that reads from offset zero to end-of-file shall 47827 47828 return one entry for dot and one entry for dot-dot; otherwise, they shall not be returned.

- Upon successful completion, *posix\_getdents()* shall mark for update the last data access 47829 timestamp of the directory. 47830
- If *fildes* is a file descriptor associated with a directory stream opened using *fdopendir()* or 47831 *opendir()*, the behavior is unspecified. 47832
- If *posix\_getdents*() is called concurrently with an operation that adds, deletes, or modifies a 47833 47834 directory entry, the results from *posix\_getdents()* shall reflect either all of the effects of the concurrent operation or none of them. If a sequence of calls to *posix\_getdents()* is made that reads 47835 from offset zero to end-of-file and a file is removed from or added to the directory between the 47836 first and last of those calls, whether the sequence of calls returns an entry for that file is 47837 47838 unspecified.

#### **RETURN VALUE** 47839

Upon successful completion, either a non-negative integer shall be returned indicating the 47840 number of bytes occupied by the **posix\_dent** structures placed in *buf* or 0 shall be returned 47841 indicating the end of the directory was reached without any directory entries being placed in buf. 47842 Otherwise, –1 shall be returned and *errno* shall be set to indicate the error. 47843

#### ERRORS 47844

47868

The *posix\_getdents()* function shall fail if: 47845 [EBADF] The *fildes* argument is not a valid file descriptor open for reading. 47846 47847 [EINVAL] The *nbyte* argument is not large enough to contain the information to be returned about the directory entry located at the current file offset. 47848 [ENOENT] The current file offset is not located at a valid directory entry. 47849 [ENOTDIR] The *fildes* argument is associated with a non-directory file. 47850 [EOVERFLOW] 47851 One of the values in a structure to be placed in *buf* cannot be represented correctly. 47852 The *posix\_getdents()* function may fail if: 47853 [EIO] A physical I/O error has occurred. 47854 47855 [ENOMEM] Insufficient memory was available to fulfill the request. **EXAMPLES** 47856 This example function lists the files in a specified directory with their file serial number and file 47857 type. If the file type is not available from *posix\_getdents()*, it is obtained using *fstatat()*. 47858 #include <dirent.h> 47859 #include <fcntl.h> 47860 #include <stdio.h> 47861 #include <stdlib.h> 47862 47863 #include <sys/stat.h> #include <unistd.h> 47864 47865 #define ENTBUFSIZ 10240 int list\_dir(const char \*dirnam) 47866 47867 { int fd = open(dirnam, O\_RDONLY | O\_DIRECTORY);

```
if (fd == -1)
47869
47870
                      return -1;
47871
                 char *buf = malloc(ENTBUFSIZ);
                 if (buf == NULL)
47872
47873
                 {
                     close(fd);
47874
                     return -1;
47875
                 }
47876
                 ssize_t bytesinbuf;
47877
                 for(;;)
47878
                 {
47879
                      ssize_t nextent = 0;
47880
47881
                     bytesinbuf = posix_getdents(fd, buf, ENTBUFSIZ, 0);
                      if (bytesinbuf <= 0)
47882
47883
                          break;
47884
                      do {
                          const char *ftype;
47885
47886
                          struct posix_dent *entp = (void *)&buf[nextent];
                          if (entp->d_type == DT_UNKNOWN)
47887
47888
                          {
                               struct stat stbuf;
47889
                               if (fstatat(fd, entp->d_name, &stbuf,
47890
                                            AT_SYMLINK_NOFOLLOW) == -1)
47891
                                   ftype = "?";
47892
                               else
47893
                                   ftype = S_ISBLK(stbuf.st_mode) ? "b" :
47894
                                         S_ISCHR(stbuf.st_mode) ? "c" :
47895
                                         S_ISDIR(stbuf.st_mode) ? "d" :
47896
47897
                                         S_ISFIFO(stbuf.st_mode) ? "p" :
                                         S_ISLNK(stbuf.st_mode) ? "1" :
47898
                                         S_ISREG(stbuf.st_mode) ? "r" :
47899
                                         S_ISSOCK(stbuf.st_mode) ? "s" :
47900
                                         S_TYPEISMQ(&stbuf) ? "mq" :
47901
                                         S_TYPEISSEM(&stbuf) ? "sem" :
47902
                                         S TYPEISSHM(&stbuf) ? "shm" :
47903
            #ifdef S_TYPEISTMO
47904
                                         S_TYPEISTMO(&stbuf) ? "tmo" :
47905
            #endif
47906
                                         "?";
47907
47908
                          }
                          else
47909
47910
                          {
                               ftype = entp->d_type == DT_BLK ? "b" :
47911
                                        entp->d_type == DT_CHR ? "c" :
47912
                                        entp->d_type == DT_DIR ? "d" :
47913
47914
                                        entp->d_type == DT_FIFO ? "p" :
                                        entp->d_type == DT_LNK ? "1" :
47915
                                        entp->d_type == DT_REG ? "r" :
47916
                                        entp->d_type == DT_SOCK ? "s" :
47917
                                        entp->d_type == DT_MQ ? "mq" :
47918
```

1432

```
entp->d_type == DT_SEM ? "sem"
47919
                                              entp->d_type == DT_SHM ? "shm"
47920
47921
              #ifdef DT_TMO
47922
                                              entp->d type == DT TMO ? "tmo" :
              #endif
47923
                                              "?";
47924
                              }
47925
                              printf("%ld\t%s\t%s\n", (long)entp->d_ino, ftype,
47926
                                   entp->d_name);
47927
                              nextent += entp->d_reclen;
47928
                         } while (nextent < bytesinbuf);</pre>
47929
                    }
47930
47931
                    close(fd);
                    free(buf);
47932
                    return bytesinbuf;
47933
47934
              }
      APPLICATION USAGE
47935
              If an array of posix_dent structures (which is only possible on implementations where d_name is
47936
              not a flexible array member) is used to provide the storage for buf in order to satisfy the
47937
              alignment requirement, it should be noted that the number of array elements used to size the
47938
              array may bear little or no relation to the number of directory entries that can be stored in it. It is
47939
              recommended that the number of elements is calculated from the desired size in bytes, for
47940
47941
              example:
              #define DESIREDSIZE 10240
47942
              struct posix_dent buf[DESIREDSIZE / sizeof(struct posix_dent) + 1];
47943
              size_t nbyte = sizeof buf;
47944
              When posix_getdents() is called with a buf that is not type char *, it is important to note that
47945
              d_reclen is a byte count and therefore any pointer arithmetic involved in calculating the start of
47946
              the next entry needs to use a char * pointer.
47947
              On implementations where directory entries in a directory take up more space than the
47948
              corresponding posix_dent structures in buf, a call to posix_getdents() may read nbyte bytes from
47949
              the directory, resulting (in most cases) in the actual number of bytes placed in buf being less than
47950
              nbyte.
47951
              One advantage of posix_getdents() is that it provides the file type of each directory entry (if
47952
              available), whereas readdir() only does so on implementations that have the file type as a non-
47953
              standard additional member of the dirent structure. Knowing the file type can greatly reduce the
47954
              number of fstatat() calls that need to be made when traversing the file hierarchy.
47955
              Whether or not a file's type can be determined without needing to use the file serial number to
47956
47957
              obtain the file's metadata may vary across the different file system types supported by an
              implementation. Therefore applications should not assume that if d_{type} contains known file
47958
              types (i.e. not DT_UNKNOWN) for entries in a given directory then it will also contain known
47959
              file types for entries in subdirectories of that directory or in its parent.
47960
              Since the d_reclen value for the last entry in buf includes padding to satisfy alignment
47961
              requirements, applications can grow the buffer and call posix_getdents() again to append to it
47962
              without needing to perform an alignment calculation.
47963
```

#### 47964 RATIONALE

- 47965The posix\_getdents() function was derived from existing getdents() functions but the name was47966changed because the existing getdents() functions differed in various ways, in particular the type47967of the second argument (structure pointer or void \*), the members of the populated structures,47968and the error numbers used for some conditions. The name change also provided an47969opportunity to add a flags argument to provide for future extensibility.
- 47970Implementations are encouraged to include support for a DT\_FORCE\_TYPE flag which, when47971that bit is set in *flags*, causes *posix\_getdents*() to look up the file type if it can not be obtained from47972the directory entry. This will allow applications that need to know the file type of every directory47973entry to keep the cost of these lookups to the minimum needed to obtain the type at the file47974system level, without the additional overhead of making a call to *fstatat*() for every file (that has47975 $d_type$  equal to DT\_UNKNOWN).
- 47976Some existing getdents() or similar functions return directory entry structures for deleted47977directory entries in buf, marked with a special value of one of the structure members to47978distinguish them from non-deleted entries. This behavior is not allowed for posix\_getdents(),47979although the data from a deleted directory entry may be present in buf in the form of extra47980padding on the end of the previous entry.

#### 47981 FUTURE DIRECTIONS

- 47982 A future version of this standard may add a DT\_FORCE\_TYPE flag as described in 47983 RATIONALE.
- 47984 SEE ALSO
- 47985 *fdopendir(), fstatat(), lseek(), readdir()*
- 47986 XBD **<dirent.h>**
- 47987 CHANGE HISTORY

47988 First released in Issue 8.

47989

## ppoll()

	NAME
49675	ppoll — input/output multiplexing
49676	SYNOPSIS
49677	<pre>#include <poll.h></poll.h></pre>
49678	<pre>int ppoll(struct pollfd fds[], nfds_t nfds,</pre>
49679	const struct timespec *restrict timeout,
49680	<pre>const sigset_t *restrict sigmask);</pre>
49681	DESCRIPTION
49682	Refer to <i>poll</i> ().

Santis

49836 of the *pselect*() call.

**RETURN VALUE** 49837 49838 Upon successful completion, the *pselect()* and *select()* functions shall return the total number of bits set in the bit masks. Otherwise, -1 shall be returned, and errno shall be set to indicate the 49839 49840 error. FD\_CLR(), FD\_SET(), and FD\_ZERO() do not return a value. FD\_ISSET() shall return a non-49841 zero value if the bit for the file descriptor fd is set in the file descriptor set pointed to by fdset, and 49842 49843 0 otherwise. ERRORS 49844 Under the following conditions, *pselect()* and *select()* shall fail and set *errno* to: 49845 [EBADF] One or more of the file descriptor sets specified a file descriptor that is not a 49846 valid open file descriptor. 49847 [EINTR] The function was interrupted by a signal. 49848 If SA\_RESTART has been set for the interrupting signal, it is implementation-49849 49850 defined whether the function restarts or returns with [EINTR]. [EINVAL] An invalid timeout interval was specified. 49851 [EINVAL] The *nfds* argument is less than 0 or greater than FD\_SETSIZE. 49852 **EXAMPLES** 49853 None. 49854 APPLICATION USAGE 49855 The use of *select()* and *pselect()* requires that the application construct the set of file descriptors 49856 to work on each time through a polling loop, and is inherently limited from operating on file 49857 descriptors larger than FD\_SETSIZE. Also, the amount of work to perform scales as nfds 49858 increases, even if the number of file descriptors selected within the larger set remains the same. 49859 Thus, applications may wish to consider using *poll()* and *ppoll()* instead, for better scaling. 49860 RATIONALE 49861 In earlier versions of the Single UNIX Specification, the select() function was defined in the 49862 <sys/time.h> header. This is now changed to <sys/select.h>. The rationale for this change was 49863 as follows: the introduction of the *pselect()* function included the **<sys/select.h>** header and the 49864 <sys/select.h> header defines all the related definitions for the *pselect()* and *select()* functions. 49865 49866 Backwards-compatibility to existing XSI implementations is handled by allowing **<sys/time.h>** to include <sys/select.h>. 49867 Code which wants to avoid the ambiguity of the signal mask for thread cancellation handlers 49868 can install an additional cancellation handler which resets the signal mask to the expected value. 49869 void cleanup(void \*arg) 49870 { 49871 sigset\_t \*ss = (sigset\_t \*) arg; 49872 pthread\_sigmask(SIG\_SETMASK, ss, NULL); 49873 49874 } int call\_pselect(int nfds, fd\_set \*readfds, fd\_set \*writefds, 49875 49876 fd\_set errorfds, const struct timespec \*timeout, const sigset\_t \*sigmask) 49877 { 49878 sigset\_t oldmask; 49879 int result; 49880

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## pthread\_cond\_broadcast()

51397 51398	NAME pthread_cond_broadcast, pthread_cond_signal — broadcast or signal a condition
51399	SYNOPSIS
51400	<pre>#include <pthread.h></pthread.h></pre>
51401 51402	int pthread_cond_broadcast(pthread_cond_t * <i>cond</i> ); int pthread_cond_signal(pthread_cond_t * <i>cond</i> );
51403	DESCRIPTION
51403	These functions shall unblock threads blocked on a condition variable.
51405 51406	The <i>pthread_cond_broadcast()</i> function shall unblock all threads currently blocked on the specified condition variable <i>cond</i> .
51407 51408	The <i>pthread_cond_signal()</i> function shall unblock at least one of the threads that are blocked on the specified condition variable <i>cond</i> (if any threads are blocked on <i>cond</i> ).
51409 51410 51411 51412 51413 51413 51414 51415	If more than one thread is blocked on a condition variable, the scheduling policy shall determine the order in which threads are unblocked. When each thread unblocked as a result of a <i>pthread_cond_broadcast()</i> or <i>pthread_cond_signal()</i> returns from its call to <i>pthread_cond_clockwait()</i> ,   <i>pthread_cond_timedwait()</i> , or <i>pthread_cond_wait()</i> , the thread shall own the mutex with which it + called <i>pthread_cond_clockwait()</i> , <i>pthread_cond_timedwait()</i> , or <i>pthread_cond_timedwait()</i> . The thread(s)   that are unblocked shall contend for the mutex according to the scheduling policy (if applicable), and as if each had called <i>pthread_mutex_lock()</i> .
51416 51417 51418 51419 51420	The <i>pthread_cond_broadcast()</i> or <i>pthread_cond_signal()</i> functions may be called by a thread whether or not it currently owns the mutex that threads calling <i>pthread_cond_clockwait()</i> ,   <i>pthread_cond_timedwait()</i> , or <i>pthread_cond_wait()</i> have associated with the condition variable   during their waits; however, if predictable scheduling behavior is required, then that mutex shall be locked by the thread calling <i>pthread_cond_broadcast()</i> or <i>pthread_cond_signal()</i> .
51421 51422	The <i>pthread_cond_broadcast()</i> and <i>pthread_cond_signal()</i> functions shall have no effect if there are no threads currently blocked on <i>cond</i> .
51423 51424	The behavior is undefined if the value specified by the <i>cond</i> argument to <i>pthread_cond_broadcast()</i> or <i>pthread_cond_signal()</i> does not refer to an initialized condition variable.
51425	RETURN VALUE
51425 51426 51427	If successful, the <i>pthread_cond_broadcast()</i> and <i>pthread_cond_signal()</i> functions shall return zero; otherwise, an error number shall be returned to indicate the error.
<b>F1 (0</b> 0	EDBODS
51428 51429	These functions shall not return an error code of [EINTR].
51420	FYAMPIES
51430 51431	None.
51432	APPLICATION USAGE
51433	The <i>pthread_cond_broadcast()</i> function is used whenever the shared-variable state has been
51434	changed in a way that more than one thread can proceed with its task. Consider a single
51435	producer/multiple consumer problem, where the producer can insert multiple items on a list
51436	that is accessed one item at a time by the consumers. By calling the <i>pthread_cond_broadcast()</i>
51437	function, the producer would notify all consumers that might be waiting, and thereby the
51438	application would receive more throughput on a multi-processor. In addition,
51439	pthread_cond_broadcast() makes it easier to implement a read-write lock. The
51440	pthread_cond_broadcast() function is needed in order to wake up all waiting readers when a
51441	writer releases its lock. Finally, the two-phase commit algorithm can use this broadcast function

51442

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to notify all clients of an impending transaction commit.

- 51443It is not safe to use the *pthread\_cond\_signal()* function in a signal handler that is invoked51444asynchronously. Even if it were safe, there would still be a race between the test of the Boolean51445*pthread\_cond\_wait()* that could not be efficiently eliminated.
- 51446 Mutexes and condition variables are thus not suitable for releasing a waiting thread by signaling 51447 from code running in a signal handler.

#### 51448 RATIONALE

51449If an implementation detects that the value specified by the *cond* argument to51450*pthread\_cond\_broadcast()* or *pthread\_cond\_signal()* does not refer to an initialized condition51451variable, it is recommended that the function should fail and report an [EINVAL] error.

#### 51452 Multiple Awakenings by Condition Signal

51453On a multi-processor, it may be impossible for an implementation of *pthread\_cond\_signal()* to51454avoid the unblocking of more than one thread blocked on a condition variable. For example,51455consider the following partial implementation of *pthread\_cond\_wait()* and *pthread\_cond\_signal()*,51456executed by two threads in the order given. One thread is trying to wait on the condition51457variable, another is concurrently executing *pthread\_cond\_signal()*, while a third thread is already51458waiting.

```
51459
            pthread_cond_wait(mutex, cond):
                value = cond->value; /* 1 */
51460
                pthread_mutex_unlock(mutex); /* 2 */
51461
                                                            */
                pthread_mutex_lock(cond->mutex); /* 10
51462
                if (value == cond->value) { /* 11 */
51463
                     me->next_cond = cond->waiter;
51464
                     cond->waiter = me;
51465
                     pthread_mutex_unlock(cond->mutex);
51466
                     unable_to_run(me);
51467
                 } else
51468
                     pthread_mutex_unlock(cond->mutex); /* 12 */
51469
51470
                pthread_mutex_lock(mutex); /* 13 */
            pthread_cond_signal(cond):
51471
                pthread_mutex_lock(cond->mutex); /* 3 */
51472
                cond->value++; /* 4 */
51473
                if (cond->waiter) { /* 5 */
51474
51475
                     sleeper = cond->waiter; /* 6 */
                     cond->waiter = sleeper->next_cond; /* 7 */
51476
                     able_to_run(sleeper); /* 8 */
51477
51478
                 }
                pthread_mutex_unlock(cond->mutex); /* 9 */
51479
```

51480The effect is that more than one thread can return from its call to pthread\_cond\_clockwait(),51481pthread\_cond\_timedwait(), or pthread\_cond\_wait() as a result of one call to pthread\_cond\_signal().51482This effect is called ``spurious wakeup''. Note that the situation is self-correcting in that the51483number of threads that are so awakened is finite; for example, the next thread to call51484pthread\_cond\_wait() after the sequence of events above blocks.

51485While this problem could be resolved, the loss of efficiency for a fringe condition that occurs51486only rarely is unacceptable, especially given that one has to check the predicate associated with a51487condition variable anyway. Correcting this problem would unnecessarily reduce the degree of51488concurrency in this basic building block for all higher-level synchronization operations.

51489 An added benefit of allowing spurious wakeups is that applications are forced to code a

## pthread\_cond\_timedwait()

+

+ + +

#### 51637 NAME

51638 pthread\_cond\_clockwait, pthread\_cond\_timedwait, pthread\_cond\_wait — wait on a condition

#### 51639 SYNOPSIS

51640	include <pthread.h></pthread.h>
51641 51642	<pre>nt pthread_cond_clockwait(pthread_cond_t *restrict cond,</pre>
51643	<pre>const struct timespec *restrict abstime);</pre>
51644	nt pthread_cond_timedwait(pthread_cond_t *restrict <i>cond</i> ,
51645	<pre>pthread_mutex_t *restrict mutex,</pre>
51646	<pre>const struct timespec *restrict abstime);</pre>
51647	nt pthread_cond_wait(pthread_cond_t *restrict <i>cond</i> ,
51648	<pre>pthread_mutex_t *restrict mutex);</pre>

#### 51649 **DESCRIPTION**

51650The pthread\_cond\_clockwait(), pthread\_cond\_timedwait(), and pthread\_cond\_wait() functions shall51651block on a condition variable. The application shall ensure that these functions are called with51652mutex locked by the calling thread; otherwise, an error (for PTHREAD\_MUTEX\_ERRORCHECK51653and robust mutexes) or undefined behavior (for other mutexes) results.

51654These functions atomically release *mutex* and cause the calling thread to block on the condition51655variable *cond*; atomically here means ``atomically with respect to access by another thread to the51656mutex and then the condition variable". That is, if another thread is able to acquire the mutex51657after the about-to-block thread has released it, then a subsequent call to *pthread\_cond\_broadcast()*51658or *pthread\_cond\_signal()* in that thread shall behave as if it were issued after the about-to-block51659thread has blocked.

- 51660 Upon successful return, the mutex shall have been locked and shall be owned by the calling 51661 thread.
- 51662 If *mutex* is a robust mutex where an owner terminated while holding the lock and the state is 51663 recoverable, the mutex shall be acquired even though the function returns [EOWNERDEAD].

51664When using condition variables there is always a Boolean predicate involving shared variables51665associated with each condition wait that is true if the thread should proceed. Spurious wakeups51666from the *pthread\_cond\_clockwait()*, *pthread\_cond\_timedwait()*, or *pthread\_cond\_wait()* functions51667may occur. Since the return from *pthread\_cond\_clockwait()*, *pthread\_cond\_timedwait()*, or51668*pthread\_cond\_wait()* does not imply anything about the value of this predicate, the predicate51669should be re-evaluated upon such return.

When a thread waits on a condition variable, having specified a particular mutex to the 51670 pthread\_cond\_clockwait(), pthread\_cond\_timedwait(), or pthread\_cond\_wait() operation, a dynamic 51671 binding is formed between that mutex and condition variable that remains in effect as long as at 51672 least one thread is blocked on the condition variable. During this time, the effect of an attempt 51673 by any thread to wait on that condition variable using a different mutex is undefined. Once all 51674 waiting threads have been unblocked (as by the *pthread\_cond\_broadcast(*) operation), the next 51675 wait operation on that condition variable shall form a new dynamic binding with the mutex 51676 specified by that wait operation. Even though the dynamic binding between condition variable 51677 and mutex may be removed or replaced between the time a thread is unblocked from a wait on 51678 the condition variable and the time that it returns to the caller or begins cancellation cleanup, the 51679 unblocked thread shall always re-acquire the mutex specified in the condition wait operation 51680 call from which it is returning. 51681

51682A condition wait (whether timed or not) is a cancellation point. When the cancelability type of a51683thread is set to PTHREAD\_CANCEL\_DEFERRED, a side-effect of acting upon a cancellation51684request while in a condition wait is that the mutex is (in effect) re-acquired before calling the first

51685cancellation cleanup handler. The effect is as if the thread were unblocked, allowed to execute up51686to the point of returning from the call to *pthread\_cond\_clockwait(), pthread\_cond\_timedwait()*, or51687*pthread\_cond\_wait()*, but at that point notices the cancellation request and, instead of returning to51688the caller, starts the thread cancellation activities, which includes calling cancellation cleanup51689handlers.

- 51690A thread that has been unblocked because it has been canceled while blocked in a call to51691pthread\_cond\_clockwait(), pthread\_cond\_timedwait(), or pthread\_cond\_wait() shall not consume any51692condition signal that may be directed concurrently at the condition variable if there are other51693threads blocked on the condition variable.
- 51694The pthread\_cond\_timedwait() function shall be equivalent to pthread\_cond\_wait(), except that an51695error is returned if the absolute time specified by abstime passes (that is, system time equals or51696exceeds abstime) before the condition cond is signaled or broadcasted, or if the absolute time51697specified by abstime has already been passed at the time of the call. When such timeouts occur,51698pthread\_cond\_timedwait() shall nonetheless release and re-acquire the mutex referenced by mutex,51699and may consume a condition signal directed concurrently at the condition variable.
- 51700The condition variable shall have a clock attribute which specifies the clock that shall be used by51701pthread\_cond\_timedwait() to measure the time specified by the abstime argument. The51702pthread\_cond\_timedwait() function is also a cancellation point.
- 51703The pthread\_cond\_clockwait() function shall be equivalent to pthread\_cond\_timedwait(), except that +51704the absolute time specified by abstime is measured against the clock indicated by clock\_id rather +51705than the clock specified in the condition variable's clock attribute. Implementations shall +51706support passing CLOCK\_REALTIME and CLOCK\_MONOTONIC to pthread\_cond\_clockwait() as +51707the clock\_id argument.
- 51708 If a signal is delivered to a thread waiting for a condition variable, upon return from the signal 51709 handler the thread resumes waiting for the condition variable as if it was not interrupted, or it 51710 shall return zero due to spurious wakeup.
- 51711 The behavior is undefined if the value specified by the *cond* or *mutex* argument to these 51712 functions does not refer to an initialized condition variable or an initialized mutex object, 51713 respectively.

#### 51714 **RETURN VALUE**

- 51715 Except for [ETIMEDOUT], [ENOTRECOVERABLE], and [EOWNERDEAD], all these error 51716 checks shall act as if they were performed immediately at the beginning of processing for the 51717 function and shall cause an error return, in effect, prior to modifying the state of the mutex 51718 specified by *mutex* or the condition variable specified by *cond*.
- 51719 Upon successful completion, a value of zero shall be returned; otherwise, an error number shall 51720 be returned to indicate the error.

#### 51721 ERRORS

51726

- 51722 These functions shall fail if:
- 51723[EAGAIN]The mutex is a robust mutex and the system resources available for robust51724mutexes owned would be exceeded.

#### 51725 [ENOTRECOVERABLE]

The state protected by the mutex is not recoverable.

#### 51727 [EOWNERDEAD]

- 51728The mutex is a robust mutex and the process containing the previous owning51729thread terminated while holding the mutex lock. The mutex lock shall be51730acquired by the calling thread and it is up to the new owner to make the state
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## pthread\_cond\_timedwait()

51731		consistent.	
51732 51733	[EPERM]	The mutex type is PTHREAD_MUTEX_ERRORCHECK or the mutex is a robust mutex, and the current thread does not own the mutex.	
51734	The <i>pthread_cond</i>	_clockwait() and pthread_cond_timedwait() functions shall fail if:	ļ
51735	[ETIMEDOUT]	The time specified by <i>abstime</i> has passed.	-
51736 51737 51738	[EINVAL]	The <i>abstime</i> argument specified a nanosecond value less than zero or greater than or equal to 1000 million, or the <i>clock_id</i> argument passed to <i>pthread_cond_clockwait()</i> is invalid or not supported.	I
51739	These functions r	nay fail if:	
51740 51741 51742 51743	[EOWNERDEAI	)] The mutex is a robust mutex and the previous owning thread terminated while holding the mutex lock. The mutex lock shall be acquired by the calling thread and it is up to the new owner to make the state consistent.	
51744	These functions s	shall not return an error code of [EINTR].	
51745 51746	EXAMPLES None.		
51747 51748 51749 51750 51751 51752 51753	APPLICATION USAGE Applications that use with robust read a currently incom- returns, due to ru an application is values for error com-	t have assumed that non-zero return values are errors will need updating for nutexes, since a valid return for a thread acquiring a mutex which is protecting nsistent state is [EOWNERDEAD]. Applications that do not check the error aling out the possibility of such errors arising, should not use robust mutexes. If supposed to work with normal and robust mutexes, it should check all return conditions and if necessary take appropriate action.	
51754 51755 51756 51757 51758 51759	<b>RATIONALE</b> If an implementation detects that the value specified by the <i>cond</i> argument to <i>pthread_cond_clockwait()</i> , <i>pthread_cond_timedwait()</i> , or <i>pthread_cond_wait()</i> does not refer to an initialized condition variable, or detects that the value specified by the <i>mutex</i> argument does not refer to an initialized mutex object, it is recommended that the function should fail and report an [EINVAL] error.		
51760	Condition Wait	Semantics	
51761 51762 51763 51764 51765	It is important <i>pthread_cond_wai</i> . when <i>pthread_cond_wai</i> . associated predict timeout and the p	to note that when <i>pthread_cond_clockwait()</i> , <i>pthread_cond_timedwait()</i> , and <i>t()</i> return without error, the associated predicate may still be false. Similarly, <i>nd_clockwait()</i> or <i>pthread_cond_timedwait()</i> returns with the timeout error, the cate may be true due to an unavoidable race between the expiration of the predicate state change.	
51766	The application 1	needs to recheck the predicate on any return because it cannot be sure there is	

- another thread waiting on the thread to handle the signal, and if there is not then the signal is lost. The burden is on the application to check the predicate.
- 51769 Some implementations, particularly on a multi-processor, may sometimes cause multiple 51770 threads to wake up when the condition variable is signaled simultaneously on different 51771 processors.
- 51772 In general, whenever a condition wait returns, the thread has to re-evaluate the predicate 51773 associated with the condition wait to determine whether it can safely proceed, should wait 51774 again, or should declare a timeout. A return from the wait does not imply that the associated

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51775 predicate is either true or false.

51776 It is thus recommended that a condition wait be enclosed in the equivalent of a ``while loop'' 51777 that checks the predicate.

#### 51778 Timed Wait Semantics

51779An absolute time measure was chosen for specifying the timeout parameter for two reasons.51780First, a relative time measure can be easily implemented on top of a function that specifies51781absolute time, but there is a race condition associated with specifying an absolute timeout on top51782of a function that specifies relative timeouts. For example, assume that clock\_gettime() returns51783the current time and cond\_relative\_timed\_wait() uses relative timeouts:

```
51784 clock_gettime(CLOCK_REALTIME, &now)
51785 reltime = sleep_til_this_absolute_time -now;
51786 cond_relative_timed_wait(c, m, &reltime);
```

51787If the thread is preempted between the first statement and the last statement, the thread blocks51788for too long. Blocking, however, is irrelevant if an absolute timeout is used. An absolute timeout51789also need not be recomputed if it is used multiple times in a loop, such as that enclosing a51790condition wait.

51791 For cases when the system clock is advanced discontinuously by an operator, it is expected that 51792 implementations process any timed wait expiring at an intervening time as if that time had 51793 actually occurred.

#### 51794 Choice of Clock

Care should be taken to decide which clock is most appropriate when waiting with a timeout. + 51795 The system clock CLOCK\_REALTIME, as used by default with *pthread\_cond\_timedwait()*, may be + 51796 subject to jumps forwards and backwards in order to correct it against actual time. + 51797 CLOCK\_MONOTONIC is guaranteed not to jump backwards and must also advance in real 51798 + time, so using it via *pthread\_cond\_clockwait()* or *pthread\_condattr\_setclock()* may be more + 51799 51800 appropriate.

#### 51801 Cancellation and Condition Wait

A condition wait, whether timed or not, is a cancellation point. That is, the functions + *pthread\_cond\_clockwait()*, *pthread\_cond\_timedwait()*, and *pthread\_cond\_wait()* are points where a pending (or concurrent) cancellation request is noticed. The reason for this is that an indefinite wait is possible at these points—whatever event is being waited for, even if the program is totally correct, might never occur; for example, some input data being awaited might never be sent. By making condition wait a cancellation point, the thread can be canceled and perform its cancellation cleanup handler even though it may be stuck in some indefinite wait.

51809 A side-effect of acting on a cancellation request while a thread is blocked on a condition variable is to re-acquire the mutex before calling any of the cancellation cleanup handlers. This is done in 51810order to ensure that the cancellation cleanup handler is executed in the same state as the critical 51811 code that lies both before and after the call to the condition wait function. This rule is also 51812 required when interfacing to POSIX threads from languages, such as Ada or C++, which may 51813 choose to map cancellation onto a language exception; this rule ensures that each exception 51814 51815 handler guarding a critical section can always safely depend upon the fact that the associated mutex has already been locked regardless of exactly where within the critical section the 51816 exception was raised. Without this rule, there would not be a uniform rule that exception 51817 handlers could follow regarding the lock, and so coding would become very cumbersome. 51818

### pthread\_cond\_timedwait()

**Timed Condition Wait** 51859 The *pthread\_cond\_clockwait()* and *pthread\_cond\_timedwait()* functions allow an application to give 51860 51861 up waiting for a particular condition after a given amount of time. An example follows: (void) pthread\_mutex\_lock(&t.mn); 51862 t.waiters++; 51863 clock\_gettime(CLOCK\_MONOTONIC, &ts); 51864 ts.tv\_sec += 5; 51865 rc = 0;51866 while (! mypredicate(&t) && rc == 0) 51867 rc = pthread\_cond\_clockwait(&t.cond, &t.mn, 51868 CLOCK\_MONOTONIC, &ts); 51869 t.waiters--; 51870 if (rc == 0 mypredicate(&t)) 51871 setmystate(&t); 51872 (void) pthread\_mutex\_unlock(&t.mn); 51873 By making the timeout parameter absolute, it does not need to be recomputed each time the 51874 program checks its blocking predicate. If the timeout was relative, it would have to be 51875 recomputed before each call. This would be especially difficult since such code would need to 51876 51877 take into account the possibility of extra wakeups that result from extra broadcasts or signals on the condition variable that occur before either the predicate is true or the timeout is due. Using 51878 CLOCK\_MONOTONIC rather than CLOCK\_REALTIME means that the timeout is not 51879 influenced by the system clock being changed. 51880 **FUTURE DIRECTIONS** 51881 None. 51882 SEE ALSO 51883 pthread\_cond\_broadcast() 51884 XBD Section 4.13 (on page 91), <pthread.h> 51885 CHANGE HISTORY 51886 First released in Issue 5. Included for alignment with the POSIX Threads Extension. 51887 Issue 6 51888 The *pthread\_cond\_timedwait()* and *pthread\_cond\_wait()* functions are marked as part of the 51889 Threads option. 51890 The Open Group Corrigendum U021/9 is applied, correcting the prototype for the 51891 *pthread\_cond\_wait()* function. 51892 The DESCRIPTION is updated for alignment with IEEE Std 1003.1j-2000 by adding semantics 51893 for the Clock Selection option. 51894 The ERRORS section has an additional case for [EPERM] in response to IEEE PASC 51895 Interpretation 1003.1c #28. 51896 The **restrict** keyword is added to the *pthread\_cond\_timedwait()* and *pthread\_cond\_wait()* 51897 prototypes for alignment with the ISO/IEC 9899: 1999 standard. 51898 51899 IEEE Std 1003.1-2001/Cor 2-2004, item XSH/TC2/D6/89 is applied, updating the DESCRIPTION for consistency with the *pthread\_cond\_destroy()* function that states it is safe to 51900 destroy an initialized condition variable upon which no threads are currently blocked. 51901 IEEE Std 1003.1-2001/Cor 2-2004, item XSH/TC2/D6/90 is applied, updating words in the 51902 DESCRIPTION from ``the cancelability enable state" to ``the cancelability type". 51903

51993 51994 51995	NAME pthread_condattr_getclock, pthread_condattr_setclock — get and set the clock selection condition variable attribute	
51996 51997	<pre>SYNOPSIS #include <pthread.h></pthread.h></pre>	
51998 51999 52000 52001	<pre>int pthread_condattr_getclock(const pthread_condattr_t *restrict attr,</pre>	
52002 52003 52004	<b>DESCRIPTION</b> The <i>pthread_condattr_getclock()</i> function shall obtain the value of the <i>clock</i> attribute from the attributes object referenced by <i>attr</i> .	
52005 52006 52007	The <i>pthread_condattr_setclock()</i> function shall set the <i>clock</i> attribute in an initialized attributes object referenced by <i>attr</i> . If <i>pthread_condattr_setclock()</i> is called with a <i>clock_id</i> argument that refers to a CPU-time clock, the call shall fail.	
52008 52009 52010	The <i>clock</i> attribute is the clock ID of the clock that shall be used to measure the timeout service of <i>pthread_cond_timedwait()</i> . The default value of the <i>clock</i> attribute shall refer to the system clock. The <i>clock</i> attribute shall have no effect on the <i>pthread_cond_clockwait()</i> function.	
52011 52012 52013	The behavior is undefined if the value specified by the <i>attr</i> argument to <i>pthread_condattr_getclock()</i> or <i>pthread_condattr_setclock()</i> does not refer to an initialized condition variable attributes object.	
52014 52015 52016 52017	<b>RETURN VALUE</b> If successful, the <i>pthread_condattr_getclock()</i> function shall return zero and store the value of the clock attribute of <i>attr</i> into the object referenced by the <i>clock_id</i> argument. Otherwise, an error number shall be returned to indicate the error.	
52018 52019	If successful, the <i>pthread_condattr_setclock()</i> function shall return zero; otherwise, an error number shall be returned to indicate the error.	
52020 52021	ERRORS The <i>pthread_condattr_setclock()</i> function may fail if:	
52022 52023	[EINVAL] The value specified by <i>clock_id</i> does not refer to a known clock, or is a CPU-time clock.	
52024	These functions shall not return an error code of [EINTR].	
52025 52026	EXAMPLES None.	
52027 52028	APPLICATION USAGE None.	
52029 52030 52031 52032 52033	<b>RATIONALE</b> If an implementation detects that the value specified by the <i>attr</i> argument to <i>pthread_condattr_getclock()</i> or <i>pthread_condattr_setclock()</i> does not refer to an initialized condition variable attributes object, it is recommended that the function should fail and report an [EINVAL] error.	

53021	NAME							
53022		pthread_mutex_c	lestroy, pthread_mutex_init — destroy and initialize a mutex					
53023	SYNOP	SIS						
53024		<pre>#include <pth< pre=""></pth<></pre>	nread.h>					
53025		int pthread_m	<pre>int pthread_mutex_destroy(pthread_mutex_t *mutex);</pre>					
53026		int pthread_m	<pre>uutex_init(pthread_mutex_t *restrict mutex,</pre>					
53027		const pth	<pre>iread_mutexattr_t *restrict attr);     t mutox = prupead MUTEX INITIALIZED.</pre>					
33028	DECOD							
53029	DESCR	IPTION The nthread mute	w destroy() function shall destroy the mutay abject referenced by mutay the					
53030 53031		mutex object	becomes in effect uninitialized An implementation may cause					
53032		pthread_mutex_des	stroy() to set the object referenced by <i>mutex</i> to an invalid value.					
53033		A destroyed mut	ex object can be reinitialized using <i>pthread mutex init()</i> ; the results of otherwise					
53034		referencing the ol	oject after it has been destroyed are undefined.					
53035		It shall be safe to	destroy an initialized mutex that is unlocked. Attempting to destroy a locked					
53036		mutex, or a mute	ex that another thread is attempting to lock, or a mutex that is being used in a					
53037		pthread_cond_clock	<i>kwait(), pthread_cond_timedwait(), or pthread_cond_wait()</i> call by another thread,					
55038								
53039		The pthread_mute	$x_{init}$ () function shall initialize the mutex referenced by <i>mutex</i> with attributes If attr is NULL the default mutex attributes are used, the effect shall be the					
53040 53041		same as passing	the address of a default mutex attributes object. Upon successful initialization					
53042		the state of the m	utex becomes initialized and unlocked.					
53043		See Section 2.9.9	(on page 508) for further requirements.					
53044		Attempting to ini	tialize an already initialized mutex results in undefined behavior.					
53045		In cases w	here default mutex attributes are appropriate, the macro					
53046		PTHREAD_MUT	EX_INITIALIZER can be used to initialize mutexes. The effect shall be					
53047		equivalent to dynamic initialization by a call to <i>pthread_mutex_init()</i> with parameter <i>attr</i> specified as NULL except that no error checks are performed						
53048		specified as NUL	L, except that no error checks are performed.					
53049	The behavior is undefined if the value specified by the <i>mutex</i> argument to							
53050	prineau_matex_destroy() does not refer to an initialized indiex.							
53051 53052		does not refer to a	an initialized mutex attributes object.					
53053	RETUR	N VALUE						
53054		If successful, the <i>pthread_mutex_destroy()</i> and <i>pthread_mutex_init()</i> functions shall return zero;						
53055		otherwise, an erro	or number shall be returned to indicate the error.					
53056	ERROR	S						
53057		The <i>pthread_mute</i> .	$x_{init}()$ function shall fail if:					
53058 53059		[EAGAIN]	The system lacked the necessary resources (other than memory) to initialize another mutex.					
53060		[ENOMEM]	Insufficient memory exists to initialize the mutex.					
53061		[EPERM]	The caller does not have the privilege to perform the operation.					

The *pthread\_mutex\_init()* function may fail if:

53062 53063 53064

[EINVAL] The attributes object referenced by *attr* has the robust mutex attribute set

without the process-shared attribute being set.

53065 These functions shall not return an error code of [EINTR].

53066 EXAMPLES

53067 None.

#### 53068 APPLICATION USAGE

None.

53069

#### 53070 RATIONALE

- 53071If an implementation detects that the value specified by the *mutex* argument to53072*pthread\_mutex\_destroy()* does not refer to an initialized mutex, it is recommended that the53073function should fail and report an [EINVAL] error.
- If an implementation detects that the value specified by the *mutex* argument to 53074 pthread\_mutex\_destroy() or pthread\_mutex\_init() refers to a locked mutex or a mutex that is 53075 53076 referenced (for example, while being used in a pthread\_cond\_clockwait(), *pthread\_cond\_timedwait()*, or *pthread\_cond\_wait()* call) by another thread, or detects that the value 53077 specified by the *mutex* argument to *pthread\_mutex\_init(*) refers to an already initialized mutex, it 53078 is recommended that the function should fail and report an [EBUSY] error. 53079
- 53080If an implementation detects that the value specified by the *attr* argument to53081*pthread\_mutex\_init()* does not refer to an initialized mutex attributes object, it is recommended53082that the function should fail and report an [EINVAL] error.

#### 53083 Alternate Implementations Possible

This volume of POSIX.1-202x supports several alternative implementations of mutexes. An 53084 implementation may store the lock directly in the object of type **pthread\_mutex\_t**. Alternatively, 53085 an implementation may store the lock in the heap and merely store a pointer, handle, or unique 53086 53087 ID in the mutex object. Either implementation has advantages or may be required on certain hardware configurations. So that portable code can be written that is invariant to this choice, this 53088 volume of POSIX.1-202x does not define assignment or equality for this type, and it uses the 53089 term ``initialize" to reinforce the (more restrictive) notion that the lock may actually reside in the 53090 mutex object itself. 53091

- 53092Note that this precludes an over-specification of the type of the mutex or condition variable and53093motivates the opaqueness of the type.
- An implementation is permitted, but not required, to have *pthread\_mutex\_destroy()* store an illegal value into the mutex. This may help detect erroneous programs that try to lock (or otherwise reference) a mutex that has already been destroyed.
- 53097 Tradeoff Between Error Checks and Performance Supported
- 53098 Many error conditions that can occur are not required to be detected by the implementation in 53099 order to let implementations trade off performance *versus* degree of error checking according to 53100 the needs of their specific applications and execution environment. As a general rule, conditions 53101 caused by the system (such as insufficient memory) are required to be detected, but conditions 53102 caused by an erroneously coded application (such as failing to provide adequate 53103 synchronization to prevent a mutex from being deleted while in use) are specified to result in 53104 undefined behavior.
- A wide range of implementations is thus made possible. For example, an implementation

53199particular, it can happen at most as many times as there are statically allocated synchronization53200objects. Dynamically allocated objects would still be initialized via *pthread\_mutex\_init()* or53201*pthread\_cond\_init()*.

53202Finally, if none of the above optimization techniques for out-of-line allocation yields sufficient53203performance for an application on some implementation, the application can avoid static53204initialization altogether by explicitly initializing all synchronization objects with the53205corresponding *pthread\_\*\_init()* functions, which are supported by all implementations. An53206implementation can also document the tradeoffs and advise which initialization technique is53207more efficient for that particular implementation.

#### 53208 Destroying Mutexes

53209A mutex can be destroyed immediately after it is unlocked. However, since attempting to53210destroy a locked mutex, or a mutex that another thread is attempting to lock, or a mutex that is53211being used in a *pthread\_cond\_clockwait(), pthread\_cond\_timedwait()*, or *pthread\_cond\_wait()* call by53212another thread, results in undefined behavior, care must be taken to ensure that no other thread53213may be referencing the mutex.

#### 53214 Robust Mutexes

Implementations are required to provide robust mutexes for mutexes with the process-shared 53215 attribute set to PTHREAD\_PROCESS\_SHARED. Implementations are allowed, but not required, 53216 provide robust mutexes when the process-shared attribute set to is to 53217 PTHREAD\_PROCESS\_PRIVATE. 53218

#### 53219 FUTURE DIRECTIONS

53220 None.

#### 53221 SEE ALSO

- 53222 pthread\_mutex\_getprioceiling(), pthread\_mutexattr\_getrobust(), pthread\_mutex\_lock(), 53223 pthread\_mutex\_timedlock(), pthread\_mutexattr\_getpshared()
- 53224 XBD **<pthread.h>**

#### 53225 CHANGE HISTORY

53226 First released in Issue 5. Included for alignment with the POSIX Threads Extension.

#### 53227 **Issue 6**

- 53228
   The pthread\_mutex\_destroy() and pthread\_mutex\_init() functions are marked as part of the

   53229
   Threads option.
- 53230The *pthread\_mutex\_timedlock()* function is added to the SEE ALSO section for alignment with53231IEEE Std 1003.1d-1999.
- 53232 IEEE PASC Interpretation 1003.1c #34 is applied, updating the DESCRIPTION.
- 53233The restrict keyword is added to the *pthread\_mutex\_init()* prototype for alignment with the53234ISO/IEC 9899: 1999 standard.

53235 Issue 7

- 53236 Changes are made from The Open Group Technical Standard, 2006, Extended API Set Part 3.
- 53237The *pthread\_mutex\_destroy()* and *pthread\_mutex\_init()* functions are moved from the Threads53238option to the Base.
- 53239 The [EINVAL] error for an uninitialized mutex or an uninitialized mutex attributes object is 53240 removed; this condition results in undefined behavior.

53518	NAME	aroad mutay clocklock athroad mutay timedlock lock a mutay	
55519	pu	nead_indiex_clocklock, plinead_indiex_timedlock — lock a indiex	т
53520 53521	SYNOPSIS #i:	nclude <pthread.h></pthread.h>	
53522 53523 53524 53525	in <sup>.</sup>	<pre>t pthread_mutex_clocklock(pthread_mutex_t *restrict mutex, clockid_t clock_id, const struct timespec *restrict abstime); t pthread_mutex_timedlock(pthread_mutex_t *restrict mutex, const struct timespec *restrict abstime);</pre>	+ +
53526	DESCRIPT	ION	
53527 53528 53529 53530 53531	The obj mu wit spe	e <i>pthread_mutex_clocklock()</i> and <i>pthread_mutex_timedlock()</i> functions shall lock the mutex ject referenced by <i>mutex</i> . If the mutex is already locked, the calling thread shall block until the utex becomes available as in the <i>pthread_mutex_lock()</i> function. If the mutex cannot be locked thout waiting for another thread to unlock the mutex, this wait shall be terminated when the ecified timeout expires.	
53532 53533 53534 53535	The clo <i>abs</i> cal	e timeout shall expire when the absolute time specified by <i>abstime</i> passes, as measured by the ock on which timeouts are based (that is, when the value of that clock equals or exceeds <i>etime</i> ), or if the absolute time specified by <i>abstime</i> has already been passed at the time of the l.	
53536 53537 53538 53539 53540	For <i>pth</i> arg Imj <i>pth</i>	r <i>pthread_mutex_timedlock()</i> , the timeout shall be based on the CLOCK_REALTIME clock. For <i>tread_mutex_clocklock()</i> , the timeout shall be based on the clock specified by the <i>clock_id</i> gument. The resolution of the timeout shall be the resolution of the clock on which it is based. plementations shall support passing CLOCK_REALTIME and CLOCK_MONOTONIC to <i>tread_mutex_clocklock()</i> as the <i>clock_id</i> argument.	
53541 53542 53543	Un imi loc	nder no circumstance shall the function fail with a timeout if the mutex can be locked mediately. The validity of the <i>abstime</i> parameter need not be checked if the mutex can be eked immediately.	
53544 53545 53546 53547	RPI TPI As PR pri thr	a consequence of the priority inheritance rules (for mutexes initialized with the IO_INHERIT protocol), if a timed mutex wait is terminated because its timeout expires, the ority of the owner of the mutex shall be adjusted as necessary to reflect the fact that this read is no longer among the threads waiting for the mutex.	
53548 53549 53550 53551 53552 53553 53554 53555 53556 53557	If the pth	<i>mutex</i> is a robust mutex and the process containing the owning thread terminated while lding the mutex lock, a call to <i>pthread_mutex_clocklock()</i> or <i>pthread_mutex_timedlock()</i> shall urn the error value [EOWNERDEAD]. If <i>mutex</i> is a robust mutex and the owning thread minated while holding the mutex lock, a call to <i>pthread_mutex_clocklock()</i> or <i>tread_mutex_timedlock()</i> may return the error value [EOWNERDEAD] even if the process in the owning thread resides has not terminated. In these cases, the mutex is locked by the read but the state it protects is marked as inconsistent. The application should ensure that the te is made consistent for reuse and when that is complete call <i>pthread_mutex_consistent()</i> . If e application is unable to recover the state, it should unlock the mutex without a prior call to <i>tread_mutex_consistent()</i> , after which the mutex is marked permanently unusable.	+
53558	If <i>n</i>	nutex does not refer to an initialized mutex object, the behavior is undefined.	
53559 53560 53561	RETURN V If s zer	ALUE successful, the <i>pthread_mutex_clocklock()</i> and <i>pthread_mutex_timedlock()</i> functions shall return to; otherwise, an error number shall be returned to indicate the error.	

## pthread\_mutex\_timedlock()

53562	ERRORS			
53563	The <i>pthread_mutex_clocklock()</i> and <i>pthread_mutex_timedlock()</i> functions shall fail if:			
53564 53565	[EAGAIN]	The mutex could not be acquired because the maximum number of recursive locks for <i>mutex</i> has been exceeded.		
53566 53567	[EAGAIN]	The mutex is a robust mutex and the system resources available for robust mutexes owned would be exceeded.		
53568 53569	[EDEADLK]	The mutex type is PTHREAD_MUTEX_ERRORCHECK and the current thread already owns the mutex.		
53570 53571 53572	[EINVAL]	The mutex was created with the protocol attribute having the value PTHREAD_PRIO_PROTECT and the calling thread's priority is higher than the mutex' current priority ceiling.		
53573 53574 53575 53576	[EINVAL]	The process or thread would have blocked, and either the <i>abstime</i> parameter specified a nanoseconds field value less than zero or greater than or equal to 1000 million, or the <i>pthread_mutex_clocklock()</i> function was passed an invalid or unsupported <i>clock_id</i> value.		
53577 53578	[ENOTRECOVE	RABLE] The state protected by the mutex is not recoverable.		
53579 53580 53581 53582 53583	[EOWNERDEAI	D] The mutex is a robust mutex and the process containing the previous owning thread terminated while holding the mutex lock. The mutex lock shall be acquired by the calling thread and it is up to the new owner to make the state consistent.		
53584	[ETIMEDOUT]	The mutex could not be locked before the specified timeout expired.		
53585	The <i>pthread_mute</i>	ex_clocklock() and pthread_mutex_timedlock() functions may fail if:		
53586	[EDEADLK]	A deadlock condition was detected.		
53587 53588 53589 53590	[EOWNERDEAI	D] The mutex is a robust mutex and the previous owning thread terminated while holding the mutex lock. The mutex lock shall be acquired by the calling thread and it is up to the new owner to make the state consistent.		
53591	This function shall not return an error code of [EINTR].			
53592	EXAMPLES			
53593	None.			
53594 53595 53596 53597 53598 53599 53600	APPLICATION USAGE Applications that have assumed that non-zero return values are errors will need updating for use with robust mutexes, since a valid return for a thread acquiring a mutex which is protecting a currently inconsistent state is [EOWNERDEAD]. Applications that do not check the error returns, due to ruling out the possibility of such errors arising, should not use robust mutexes. If an application is supposed to work with normal and robust mutexes, it should check all return values for error conditions and if necessary take appropriate action.			
53601 53602	<b>RATIONALE</b> Refer to <i>pthread</i> _	mutex_lock().		

#### 54163 NAME

<sup>54164</sup> pthread\_mutexattr\_gettype, pthread\_mutexattr\_settype — get and set the mutex type attribute

#### 54165 SYNOPSIS

- 54166 #include <pthread.h>
- 54167 int pthread\_mutexattr\_gettype(const pthread\_mutexattr\_t \*restrict attr,
- 54168 int \*restrict type);
- 54169 int pthread\_mutexattr\_settype(pthread\_mutexattr\_t \*attr, int type);

#### 54170 **DESCRIPTION**

- 54171The pthread\_mutexattr\_gettype() and pthread\_mutexattr\_settype() functions, respectively, shall get54172and set the mutex type attribute. This attribute is set in the type parameter to these functions. The54173default value of the type attribute is PTHREAD\_MUTEX\_DEFAULT.
- 54174 The type of mutex is contained in the *type* attribute of the mutex attributes. Valid mutex types 54175 include:
- 54176 PTHREAD\_MUTEX\_NORMAL
  54177 PTHREAD\_MUTEX\_ERRORCHECK
  54178 PTHREAD\_MUTEX\_RECURSIVE
  54179 PTHREAD\_MUTEX\_DEFAULT
- 54180The mutex type affects the behavior of calls which lock and unlock the mutex. See54181*pthread\_mutex\_lock()* for details. An implementation may map PTHREAD\_MUTEX\_DEFAULT54182to one of the other mutex types.
- 54183The behavior is undefined if the value specified by the *attr* argument to54184*pthread\_mutexattr\_gettype()* or *pthread\_mutexattr\_settype()* does not refer to an initialized mutex54185attributes object.

#### 54186 **RETURN VALUE**

- 54187Upon successful completion, the *pthread\_mutexattr\_gettype()* function shall return zero and store54188the value of the *type* attribute of *attr* into the object referenced by the *type* parameter. Otherwise,54189an error shall be returned to indicate the error.
- 54190 If successful, the *pthread\_mutexattr\_settype()* function shall return zero; otherwise, an error 54191 number shall be returned to indicate the error.

#### 54192 ERRORS

- 54193 The *pthread\_mutexattr\_settype()* function shall fail if:
- 54194 [EINVAL] The value *type* is invalid.
- 54195 These functions shall not return an error code of [EINTR].

#### 54196 EXAMPLES

54197 None.

#### 54198 APPLICATION USAGE

54199It is advised that an application should not use a PTHREAD\_MUTEX\_RECURSIVE mutex with54200condition variables because the implicit unlock performed in a *pthread\_cond\_clockwait()*,54201*pthread\_cond\_timedwait()*, or *pthread\_cond\_wait()* call may not actually release the mutex (if it had54202been locked multiple times). If this happens, no other thread can satisfy the condition of the54203predicate.

#### 54555 NAME

	.1 1 1 1	1 1 11 1 .1		1 11 1	1 1 1 1	1 1 6 1.	
54556	pthread_rwlock_	_clockrdlock, pthi	ead_rwlock	_timedrdlock —	- lock a read-writ	e lock for reading	+

#### 54557 SYNOPSIS

54558 #include <pthread.h>

# 54559int pthread\_rwlock\_clockrdlock (pthread\_rwlock\_t \*restrict rwlock,+54560clockid\_t clock\_id, const struct timespec \*restrict abstime);+54561int pthread\_rwlock\_timedrdlock (pthread\_rwlock\_t \*restrict rwlock,54562const struct timespec \*restrict abstime);

#### 54563 **DESCRIPTION**

- 54564The pthread\_rwlock\_clockrdlock() and pthread\_rwlock\_timedrdlock() functions shall apply a read54565lock to the read-write lock referenced by rwlock as in the pthread\_rwlock\_rdlock() function.54566However, if the lock cannot be acquired without waiting for other threads to unlock the lock,54567this wait shall be terminated when the specified timeout expires. The timeout shall expire when54568the absolute time specified by abstime passes, as measured by the clock on which timeouts are54569based (that is, when the value of that clock equals or exceeds abstime), or if the absolute time54570specified by abstime has already been passed at the time of the call.
- 54571For pthread\_rwlock\_timedrdlock(), the timeout shall be based on the CLOCK\_REALTIME clock.54572For pthread\_rwlock\_clockrdlock(), the timeout shall be based on the clock specified by the clock\_id54573argument. The resolution of the timeout shall be the resolution of the clock on which it is based.54574Implementations shall support passing CLOCK\_REALTIME and CLOCK\_MONOTONIC to54575pthread\_rwlock\_clockrdlock() as the clock\_id argument.
- 54576 Under no circumstances shall the function fail with a timeout if the lock can be acquired 54577 immediately. The validity of the *abstime* parameter need not be checked if the lock can be 54578 immediately acquired.
- 54579If a signal that causes a signal handler to be executed is delivered to a thread blocked on a read-54580write lock via a call to *pthread\_rwlock\_clockrdlock()* or *pthread\_rwlock\_timedrdlock()*, upon return +54581from the signal handler the thread shall resume waiting for the lock as if it was not interrupted.
- The calling thread may deadlock if at the time the call is made it holds a write lock on *rwlock*. The results are undefined if this function is called with an uninitialized read-write lock.

#### 54584 **RETURN VALUE**

54585The pthread\_rwlock\_clockrdlock() and pthread\_rwlock\_timedrdlock() functions shall return zero if54586the lock for reading on the read-write lock object referenced by rwlock is acquired. Otherwise, an54587error number shall be returned to indicate the error.

#### 54588 ERRORS

1654

54589	The <i>pthread_rwloo</i>	ck_clockrdlock() and pthread_rwlock_timedrdlock() functions shall fail if:	
54590	[ETIMEDOUT]	The lock could not be acquired before the specified timeout expired.	
54591	The <i>pthread_rwloo</i>	ck_clockrdlock() and pthread_rwlock_timedrdlock() functions may fail if:	
54592 54593	[EAGAIN]	The read lock could not be acquired because the maximum number of read locks for lock would be exceeded.	
54594 54595	[EDEADLK]	A deadlock condition was detected or the calling thread already holds a write lock on <i>rwlock</i> .	
54596 54597 54598	[EINVAL]	The <i>abstime</i> nanosecond value is less than zero or greater than or equal to 1000 million, or the <i>pthread_rwlock_clockrdlock()</i> function was passed an invalid or unsupported <i>clock_id</i> value.	
54599	This function sha	ll not return an error code of [EINTR].	

## pthread\_rwlock\_timedrdlock()

54600 54601	EXAMPLES None.	
54602 54603 54604	APPLICATION USAGE Applications using this function may be subject to priority inversion, as discussed in XBD Section 3.260 (on page 66).	
54605 54606 54607 54608	<b>RATIONALE</b> If an implementation detects that the value specified by the <i>rwlock</i> argument to <i>pthread_rwlock_clockrdlock()</i> or <i>pthread_rwlock_timedrdlock()</i> does not refer to an initialized read-write lock object, it is recommended that the function should fail and report an [EINVAL] error.	+
54609 54610	FUTURE DIRECTIONS None.	
54611 54612 54613	<b>SEE ALSO</b> <i>pthread_rwlock_destroy(), pthread_rwlock_rdlock(), pthread_rwlock_timedwrlock(),</i> <i>pthread_rwlock_trywrlock(), pthread_rwlock_unlock()</i>	
54614	XBD Section 3.260 (on page 66), Section 4.13 (on page 91), <pthread.h>, <time.h></time.h></pthread.h>	
54615 54616	CHANGE HISTORY First released in Issue 6. Derived from IEEE Std 1003.1j-2000.	
54617 54618	IEEE Std 1003.1-2001/Cor 2-2004, item XSH/TC2/D6/102 is applied, updating the ERRORS section so that the [EDEADLK] error includes detection of a deadlock condition.	
54619 54620	<b>Issue 7</b> The <i>pthread_rwlock_timedrdlock()</i> function is moved from the Timeouts option to the Base.	
54621 54622	The [EINVAL] error for an uninitialized read-write lock object is removed; this condition results in undefined behavior.	
54623 54624 54625	<b>Issue 8</b> Austin Group Defect 592 is applied, removing text relating to <b><time.h></time.h></b> from the SYNOPSIS and DESCRIPTION sections.	+
54626	Austin Group Defect 1216 is applied, adding <i>pthread_rwlock_clockrdlock()</i> .	
### 54627 NAME

54628 pthread\_rwlock\_clockwrlock, pthread\_rwlock\_timedwrlock — lock a read-write lock for writing +

### 54629 SYNOPSIS

54630 #include <pthread.h>

# 54631int pthread\_rwlock\_clockwrlock (pthread\_rwlock\_t \*restrict rwlock,+54632clockid\_t clock\_id, const struct timespec \*restrict abstime);+54633int pthread\_rwlock\_timedwrlock (pthread\_rwlock\_t \*restrict rwlock,54634const struct timespec \*restrict abstime);

### 54635 **DESCRIPTION**

- 54636The pthread\_rwlock\_clockwrlock() and pthread\_rwlock\_timedwrlock() functions shall apply a write54637lock to the read-write lock referenced by rwlock as in the pthread\_rwlock\_wrlock() function.54638However, if the lock cannot be acquired without waiting for other threads to unlock the lock,54639this wait shall be terminated when the specified timeout expires. The timeout shall expire when54640the absolute time specified by abstime passes, as measured by the clock on which timeouts are54641based (that is, when the value of that clock equals or exceeds abstime), or if the absolute time54642specified by abstime has already been passed at the time of the call.
- 54643For pthread\_rwlock\_timedwrlock(), the timeout shall be based on the CLOCK\_REALTIME clock.54644For pthread\_rwlock\_clockwrlock(), the timeout shall be based on the clock specified by the clock\_id54645argument. The resolution of the timeout shall be the resolution of the clock on which it is based.54646Implementations shall support passing CLOCK\_REALTIME and CLOCK\_MONOTONIC to54647pthread\_rwlock\_clockwrlock() as the clock\_id argument.
- 54648 Under no circumstances shall the function fail with a timeout if the lock can be acquired 54649 immediately. The validity of the *abstime* parameter need not be checked if the lock can be 54650 immediately acquired.
- 54651If a signal that causes a signal handler to be executed is delivered to a thread blocked on a read-54652write lock via a call to *pthread\_rwlock\_clockwrlock()* or *pthread\_rwlock\_timedwrlock()*, upon return +54653from the signal handler the thread shall resume waiting for the lock as if it was not interrupted.
- The calling thread may deadlock if at the time the call is made it holds the read-write lock. The results are undefined if this function is called with an uninitialized read-write lock.

### 54656 **RETURN VALUE**

54657The pthread\_rwlock\_clockwrlock() and pthread\_rwlock\_timedwrlock() functions shall return zero if54658the lock for writing on the read-write lock object referenced by rwlock is acquired. Otherwise, an54659error number shall be returned to indicate the error.

### 54660 ERRORS

54661	The <i>pthread_rwloo</i>	ck_clockwrlock() and pthread_rwlock_timedwrlock() functions shall fail if:	
54662	[ETIMEDOUT]	The lock could not be acquired before the specified timeout expired.	
54663	The <i>pthread_rwloo</i>	ck_clockwrlock() and pthread_rwlock_timedwrlock() functions may fail if:	
54664 54665	[EDEADLK]	A deadlock condition was detected or the calling thread already holds the <i>rwlock</i> .	
54666 54667 54668	[EINVAL]	The <i>abstime</i> nanosecond value is less than zero or greater than or equal to 1 000 million, or the <i>pthread_rwlock_clockwrlock()</i> function was passed an invalid or unsupported <i>clock_id</i> value.	
54669	This function sha	all not return an error code of [EINTR].	

# pthread\_rwlock\_timedwrlock()

54670 54671	EXAMPLES None.
54672 54673 54674	APPLICATION USAGE Applications using this function may be subject to priority inversion, as discussed in XBD Section 3.260 (on page 66).
54675 54676 54677 54678	<b>RATIONALE</b> If an implementation detects that the value specified by the <i>rwlock</i> argument to + <i>pthread_rwlock_clockwrlock()</i> or <i>pthread_rwlock_timedwrlock()</i> does not refer to an initialized read-write lock object, it is recommended that the function should fail and report an [EINVAL] error.
54679 54680	FUTURE DIRECTIONS None.
54681 54682 54683	SEE ALSO pthread_rwlock_destroy(), pthread_rwlock_rdlock(), pthread_rwlock_timedrdlock(), pthread_rwlock_trywrlock(), pthread_rwlock_unlock()
54684	XBD Section 3.260 (on page 66), Section 4.13 (on page 91), <b><pthread.h></pthread.h></b> , <b><time.h></time.h></b>
54685 54686	CHANGE HISTORY First released in Issue 6. Derived from IEEE Std 1003.1j-2000.
54687 54688	IEEE Std 1003.1-2001/Cor 2-2004, item XSH/TC2/D6/103 is applied, updating the ERRORS section so that the [EDEADLK] error includes detection of a deadlock condition.
54689 54690	Issue 7 The <i>pthread_rwlock_timedwrlock()</i> function is moved from the Timeouts option to the Base.
54691 54692	The [EINVAL] error for an uninitialized read-write lock object is removed; this condition results in undefined behavior.
54693 54694 54695	Issue 8 Austin Group Defect 592 is applied, removing text relating to <b><time.h></time.h></b> from the SYNOPSIS and DESCRIPTION sections. +
54696	Austin Group Defect 1216 is applied, adding <i>pthread_rwlock_clockwrlock()</i> .

NAME

55818

#### qsort, qsort\_r — sort a table of data 55819 + **SYNOPSIS** 55820 #include <stdlib.h> 55821 void qsort(void \*base, size\_t nel, size\_t width, 55822 int (\*compar)(const void \*, const void \*)); 55823 void qsort\_r(void \*base, size\_t nel, size\_t width, CX 55824 int (\*compar) (const void \*, const void \*, void \*), void \*arg); 55825

### 55826 **DESCRIPTION**

- 55827 CXFor *qsort*(): The functionality described on this reference page is aligned with the ISO C +55828standard. Any conflict between the requirements described here and the ISO C standard is55829unintentional. This volume of POSIX.1-202x defers to the ISO C standard.
- 55830The *qsort*() function shall sort an array of *nel* objects, the initial element of which is pointed to by55831*base*. The size of each object, in bytes, is specified by the *width* argument. If the *nel* argument has55832the value zero, the comparison function pointed to by *compar* shall not be called and no55833rearrangement shall take place.
- 55834The application shall ensure that the comparison function pointed to by *compar* does not alter the55835contents of the array. The implementation may reorder elements of the array between calls to the55836comparison function, but shall not alter the contents of any individual element.
- 55837When the same objects (consisting of width bytes, irrespective of their current positions in the<br/>array) are passed more than once to the comparison function, the results shall be consistent with<br/>one another. That is, they shall define a total ordering on the array.
- 55840The contents of the array shall be sorted in ascending order according to a comparison function.55841The compar argument is a pointer to the comparison function, which is called with two55842arguments that point to the elements being compared. The application shall ensure that the55843function returns an integer less than, equal to, or greater than 0, if the first argument is55844considered respectively less than, equal to, or greater than the second. If two members compare55845as equal, their order in the sorted array is unspecified.
- 55846CXThe  $qsort_r()$  function shall be identical to qsort() except that the comparison function compar55847takes a third argument. The arg opaque pointer passed to  $qsort_r()$  shall in turn be passed as the55848third argument to the comparison function.
- 55849 **RETURN VALUE** 
  - These functions shall not return a value.
- 55851 ERRORS

55850

55852

No errors are defined.

### 55853 EXAMPLES

55854 None.

# 55855 APPLICATION USAGE

55856The comparison function need not compare every byte, so arbitrary data may be contained in55857the elements in addition to the values being compared.

If the *compar* callback function requires any additional state outside of the items being sorted, it + can only access this state through global variables, making it potentially unsafe to use qsort() + with the same *compar* function from separate threads at the same time. The  $qsort_r()$  function + was added with the ability to pass through arbitrary arguments to the comparator, which avoids + the need to access global variables and thus making it possible to safely share a stateful +

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55863	comparator across threads.	
55864 55865 55866 55867	<b>RATIONALE</b> The requirement that each argument (hereafter referred to as <i>p</i> ) to the comparison function is a pointer to elements of the array implies that for every call, for each argument separately, all of the following expressions are non-zero:	
55868 55869 55870	((char *)p - (char *)base) % width == 0 (char *)p >= (char *)base (char *)p < (char *)base + nel * width	
55871 55872	FUTURE DIRECTIONS None.	
55873 55874	SEE ALSO alphasort()	
55875	XBD <stdlib.h></stdlib.h>	
55876 55877	CHANGE HISTORY First released in Issue 1. Derived from Issue 1 of the SVID.	
55878 55879	Issue 6 The normative text is updated to avoid use of the term ``must'' for application requirements.	
55880 55881 55882	IEEE Std 1003.1-2001/Cor 1-2002, item XSH/TC1/D6/49 is applied, adding the last sentence to the first non-shaded paragraph in the DESCRIPTION, and the following two paragraphs. The RATIONALE is also updated. These changes are for alignment with the ISO C standard.	+
55883 55884	<b>Issue 8</b> Austin Group Defect 900 is applied, adding the <i>qsort_r()</i> function.	+

# rand()

```
keystr[len++] = c;
55971
55972
55973
                        keystr[len] = ' \setminus 0';
                        printf("%s Element%0*ld\n", keystr, elementlen, i);
55974
                        len = 0;
55975
                   }
55976
```

```
55977
```

55988

55989

### Generating the Same Sequence on Different Machines

The following code defines a pair of functions that could be incorporated into applications 55978 wishing to ensure that the same sequence of numbers is generated across different machines. 55979

```
static unsigned long next = 1;
55980
            int myrand(void) /* RAND_MAX assumed to be 32767. */
55981
             {
55982
                 next = next * 1103515245 + 12345;
55983
                 return((unsigned)(next/65536) % 32768);
55984
55985
             }
            void mysrand (unsigned seed)
55986
55987
             {
```

next = seed;

### ł APPLICATION USAGE 55990

- These functions should be avoided whenever non-trivial requirements (including safety) have to 55991 55992 be fulfilled, unless seeded using getentropy().
- The *drand48()* and *random()* functions provide much more elaborate pseudo-random number 55993 generators. 55994

#### RATIONALE 55995

- The ISO C standard rand() and srand() functions allow per-process pseudo-random streams 55996 shared by all threads. Those two functions need not change, but there has to be mutual-55997 exclusion that prevents interference between two threads concurrently accessing the random 55998 number generator. 55999
- With regard to *rand()*, there are two different behaviors that may be wanted in a multi-threaded 56000 56001 program:
- A single per-process sequence of pseudo-random numbers that is shared by all threads 56002 1. that call rand() 56003
- A different sequence of pseudo-random numbers for each thread that calls *rand()* 56004 2.
- This is provided by the modified thread-safe function based on whether the seed value is global 56005 to the entire process or local to each thread. 56006
- This does not address the known deficiencies of the *rand()* function implementations, which 56007 have been approached by maintaining more state. In effect, this specifies new thread-safe forms 56008 of a deficient function. 56009
- **FUTURE DIRECTIONS** 56010

None.

# realloc()

56680	NAME		
56681		realloc, reallocarray — memory reallocators	I
56682 56683	SYNOP	#include <stdlib.h></stdlib.h>	
56684 56685	СХ	<pre>void *realloc(void *ptr, size_t size); void *reallocarray(void *ptr, size_t nelem, size_t elsize);</pre>	+
56686 56687 56688 56689	DESCR CX	<b>IPTION</b> For <i>realloc()</i> : The functionality described on this reference page is aligned with the ISO C standard. Any conflict between the requirements described here and the ISO C standard is unintentional. This volume of POSIX.1-202x defers to the ISO C standard.	+
56690 56691 56692 56693 56694 56695 56696 56697		The <i>realloc()</i> function shall deallocate the old object pointed to by <i>ptr</i> and return a pointer to a new object that has the size specified by <i>size</i> . The contents of the new object shall be the same as that of the old object prior to deallocation, up to the lesser of the new and old sizes. Any bytes in the new object beyond the size of the old object have indeterminate values. If the size of the space requested is zero, the behavior shall be implementation-defined: either a null pointer is returned, or the behavior shall be as if the size were some non-zero value, except that the behavior is undefined if the returned pointer is used to access an object. If the space cannot be allocated, the object shall remain unchanged.	
56698 56699	СХ	The <i>reallocarray()</i> function shall be equivalent to the call realloc( <i>ptr</i> , <i>nelem</i> * <i>elsize</i> ) except that overflow in the multiplication shall be an error.	+ +
56700 56701	СХ	If <i>ptr</i> is a null pointer, <i>realloc()</i> or <i>reallocarray()</i> shall be equivalent to <i>malloc()</i> for the specified size.	
56702 56703 56704	СХ	If <i>ptr</i> does not match a pointer earlier returned by a function in POSIX.1-202x that allocates memory as if by <i>malloc()</i> , or if the space has previously been deallocated by a call to <i>free()</i> , <i>realloc()</i> , or <i>reallocarray()</i> , the behavior is undefined.	
56705 56706 56707 56708 56709 56710 56711	СХ	The order and contiguity of storage allocated by successive calls to <i>realloc()</i> or <i>reallocarray()</i> is unspecified. The pointer returned if the allocation succeeds shall be suitably aligned so that it may be assigned to a pointer to any type of object and then used to access such an object in the space allocated (until the space is explicitly freed or reallocated). Each such allocation shall yield a pointer to an object disjoint from any other object. The pointer returned shall point to the start (lowest byte address) of the allocated space. If the space cannot be allocated, a null pointer shall be returned.	
56712	RETUR	N VALUE	
56713 56714	CX CX	Upon successful completion, <i>realloc()</i> and <i>reallocarray()</i> shall return a pointer to the (possibly moved) allocated space. If <i>size</i> is 0, or either <i>nelem</i> or <i>elsize</i> is 0, then either:	
56715 56716	СХ	• A null pointer shall be returned and, if <i>ptr</i> is not a null pointer, <i>errno</i> shall be set to an implementation-defined value.	
56717 56718		• A pointer to the allocated space shall be returned, and the memory object pointed to by <i>ptr</i> shall be freed. The application shall ensure that the pointer is not used to access an object.	
56719 56720 56721	CX CX	If there is not enough available memory, <i>realloc()</i> and <i>reallocarray()</i> shall return a null pointer and set <i>errno</i> to [ENOMEM]. If <i>realloc()</i> or <i>reallocarray()</i> returns a null pointer and <i>errno</i> has been set to [ENOMEM], the memory referenced by <i>ptr</i> shall not be changed.	

### 56722 ERRORS

56723	CX	The <i>realloc()</i> and	<i>reallocarray()</i> functions shall fail if:	
56724	CX	[ENOMEM]	Insufficient memory is available.	
56725	СХ	The <i>reallocarray(</i> )	function shall fail if:	4
56726		[ENOMEM]	The calculation <i>nelem</i> * <i>elsize</i> would overflow.	+

### 56727 EXAMPLES

56728 None.

### 56729 APPLICATION USAGE

The description of *realloc()* has been modified from previous versions of this standard to align 56730 with the ISO/IEC 9899: 1999 standard. Previous versions explicitly permitted a call to realloc(p, 0)56731 56732 to free the space pointed to by p and return a null pointer. While this behavior could be interpreted as permitted by this version of the standard, the C language committee has indicated 56733 that this interpretation is incorrect. Applications should assume that if *realloc()* returns a null 56734 pointer, the space pointed to by p has not been freed. Since this could lead to double-frees, 56735 implementations should also set errno if a null pointer actually indicates a failure, and 56736 applications should only free the space if *errno* was changed. 56737

### 56738 RATIONALE

56739 None.

### 56740 FUTURE DIRECTIONS

56741This standard defers to the ISO C standard. While that standard currently has language that56742might permit realloc(p, 0), where p is not a null pointer, to free p while still returning a null56743pointer, the committee responsible for that standard is considering clarifying the language to56744explicitly prohibit that alternative.

### 56745 SEE ALSO

56746 *calloc(), free(), malloc()* 

56747 XBD **<stdlib.h>** 

### 56748 CHANGE HISTORY

56749 First released in Issue 1. Derived from Issue 1 of the SVID.

### 56750 **Issue 6**

- 56751 Extensions beyond the ISO C standard are marked.
- 56752The following new requirements on POSIX implementations derive from alignment with the56753Single UNIX Specification:
- In the RETURN VALUE section, if there is not enough available memory, the setting of *errno* to [ENOMEM] is added.
- The [ENOMEM] error condition is added.

# 56757 **Issue 7**

- 56758POSIX.1-2008, Technical Corrigendum 1, XSH/TC1-2008/0495 [400], XSH/TC1-2008/0496 [400],56759XSH/TC1-2008/0497 [400], and XSH/TC1-2008/0498 [400] are applied.
- 56760POSIX.1-2008, Technical Corrigendum 2, XSH/TC2-2008/0309 [526] and XSH/TC2-2008/031056761[526,688] are applied.

+

58787 58788	NAME	sem init—initia	lize an unnamed semaphore	
50700	SVNOP			
58790	51101	<pre>#include <sem< pre=""></sem<></pre>	aphore.h>	
58791		int sem_init(	sem_t *sem, int pshared, unsigned value);	
58792	DESCR	IPTION		
58793		The <i>sem_init(</i> ) fu	nction shall initialize the unnamed semaphore referred to by <i>sem</i> . The value of	
58794		the initialized ser	naphore shall be <i>value</i> . Following a successful call to <i>sem_init()</i> , the semaphore	
58795 58796		can be used in s	ubsequent calls to sem_clockwait(), sem_destroy(), sem_post(), sem_timedwait(), and sem_twait(). This semaphore shall remain usable until the semaphore is	I
58797		destroyed. An ur	named semaphore may be implemented using a file descriptor.	I
58798		If the <i>pshared</i> arg	ument has a non-zero value, then the semaphore is shared between processes;	
58799		in this case, any	process that can access the semaphore sem can use sem for performing	
58800 58801		operations.	sem_aestroy(), sem_post(), sem_timeawait(), sem_trywait(), and sem_wait()	I
58802		If the <i>vshared</i> argu	ment is zero, then the semaphore is shared between threads of the process; any	
58803		thread in this pr	cocess can use sem for performing sem_clockwait(), sem_destroy(), sem_post(),	
58804		<pre>sem_timedwait(),</pre>	sem_trywait(), and sem_wait() operations.	
58805		See Section 2.9.9	on page 508) for further requirements.	
58806		Attempting to ini	tialize an already initialized semaphore results in undefined behavior.	
58807	RETUR	N VALUE		
58808 58809		Upon successful return 0. Otherwi	completion, the <i>sem_init()</i> function shall initialize the semaphore in <i>sem</i> and se, it shall return –1 and set <i>errno</i> to indicate the error.	
58810	ERROR	S		
58811		The <i>sem_init(</i> ) fur	nction shall fail if:	
58812		[EINVAL]	The <i>value</i> argument exceeds {SEM_VALUE_MAX}.	
58813 58814		[ENOSPC]	A resource required to initialize the semaphore has been exhausted, or the limit on semaphores ({SEM_NSEMS_MAX}) has been reached.	
58815		[EPERM]	The process lacks appropriate privileges to initialize the semaphore.	
58816		The <i>sem_init()</i> fur	nction may fail if:	
58817		[EMFILE]	All file descriptors available to the process are currently open.	
58818		[ENFILE]	The maximum allowable number of files is currently open in the system.	
58819	EXAMP	LES		
58820		None.		
58821 58822	APPLIC	CATION USAGE None.		
58823	RATIO	NALE		
58824		None.		
58825 58826	FUTUR	E DIRECTIONS None.		

58849	NAME		
58850		sem_open –	- initialize and open a named semaphore
58851	SYNOP	SIS	
58852		#include	<semaphore.h></semaphore.h>
58853		sem_t *se	<pre>m_open(const char *name, int oflag,);</pre>
58854	DESCR	IPTION	
58855		The <i>sem_oper</i>	n() function shall establish a connection between a named semaphore and a process.
58856		A named set	maphore may be implemented using a file descriptor. Following a call to <i>sem_open()</i>
58857		with semap	hore name name, the process may reference the semaphore associated with name
58858		using the ac	ddress returned from the call. This semaphore can be used in subsequent calls to
58859		sem_clockwai	t(), sem_close(), sem_post(), sem_timedwait(), sem_trywait(), and sem_wait(). The
58860		semaphore 1	remains usable by this process until the semaphore is closed by a successful call to
58861		sem_close(),	_exit(), or one of the exec functions.
58862 58863		The oflag are sem open().	gument controls whether the semaphore is created or merely accessed by the call to The following flag bits may be set in <i>oflag</i> :
58864		O_CREAT	I his flag is used to create a semaphore if it does not already exist. If O_CKEAT is
58865			set and the semaphore already exists, then O_CKEAT has no effect, except as noted
58866			flag requires a third and a fourth argument, made which is of two mode t and
50000			ridging requires a time and a fourth argument. <i>Mole</i> , which is of type <b>mode_t</b> , and
58860			value Valid initial values for semaphores are less than or equal to
58870			{SEM_VALUE_MAX}.
58871			The user ID of the semaphore shall be set to the effective user ID of the process.
58872			The group ID of the semaphore shall be set to the effective group ID of the process;
58873			however, if the <i>name</i> argument is visible in the file system, the group ID may be set
58874			to the group ID of the containing directory. The permission bits of the semaphore
58875			are set to the value of the <i>mode</i> argument except those set in the file mode creation
58876			mask of the process. When bits in mode other than the file permission bits are
58877			specified, the effect is unspecified.
58878			After the semaphore named <i>name</i> has been created by <i>sem_open()</i> with the
58879			O_CREAT flag, other processes can connect to the semaphore by calling
58880			sem_open() with the same value of name.
58881		O_EXCL	If O_EXCL and O_CREAT are set, sem_open() fails if the semaphore name exists.
58882			The check for the existence of the semaphore and the creation of the semaphore if it
58883			does not exist are atomic with respect to other processes executing <i>sem_open()</i> with
58884			O_EXCL and O_CREAT set. If O_EXCL is set and O_CREAT is not set, the effect is
58885			undefined.
58886			If flags other than O_CREAT and O_EXCL are specified in the oflag parameter, the
58887			effect is unspecified.
58888		The name ar	gument points to a string naming a semaphore object. It is unspecified whether the
58889		name appea	rs in the file system and is visible to functions that take pathnames as arguments.
58890		The name a	rgument conforms to the construction rules for a pathname, except that the
58891		interpretatio	n of <slash> characters other than the leading <slash> character in name is</slash></slash>
58892		implementa	tion-defined, and that the length limits for the name argument are implementation-
58893		defined and	need not be the same as the pathname limits {PATH_MAX} and {NAME_MAX}. If
58894		name begins	with the <slash> character, then processes calling <i>sem_open()</i> with the same value of</slash>
58895		<i>name</i> shall r	eter to the same semaphore object, as long as that name has not been removed. If

+

+

### 59036 NAME

59037 sem\_clockwait, sem\_timedwait — lock a semaphore

### 59038 SYNOPSIS

59039 #include <semaphore.h>

```
59040 int sem_clockwait(sem_t *restrict sem, clockid_t clock_id,
59041 const struct timespec *restrict abstime);
59042 int sem_timedwait(sem_t *restrict sem,
59043 const struct timespec *restrict abstime);
```

### 59044 **DESCRIPTION**

- 59045The sem\_clockwait() and sem\_timedwait() functions shall lock the semaphore referenced by sem as59046in the sem\_wait() function. However, if the semaphore cannot be locked without waiting for59047another process or thread to unlock the semaphore by performing a sem\_post() function, this59048wait shall be terminated when the specified timeout expires.
- 59049The timeout shall expire when the absolute time specified by *abstime* passes, as measured by the59050clock on which timeouts are based (that is, when the value of that clock equals or exceeds59051*abstime*), or if the absolute time specified by *abstime* has already been passed at the time of the59052call.
- 59053For sem\_timedwait(), the timeout shall be based on the CLOCK\_REALTIME clock. For59054sem\_clockwait(), the timeout shall be based on the clock specified by the clock\_id argument. The59055resolution of the timeout shall be the resolution of the clock on which it is based.59056Implementations shall support passing CLOCK\_REALTIME and CLOCK\_MONOTONIC to59057sem\_clockwait() as the clock\_id argument.
- 59058 Under no circumstance shall the function fail with a timeout if the semaphore can be locked 59059 immediately. The validity of the *abstime* need not be checked if the semaphore can be locked 59060 immediately.

### 59061 **RETURN VALUE**

59062The sem\_clockwait() and sem\_timedwait() functions shall return zero if the calling process59063successfully performed the semaphore lock operation on the semaphore designated by sem. If59064the call was unsuccessful, the state of the semaphore shall be unchanged, and the functions shall59065return a value of -1 and set errno to indicate the error.

### 59066 ERRORS

59067	The <i>sem_clockwait()</i> and <i>sem_timedwait()</i> functions shall fail if:				
59068 59069 59070 59071	[EINVAL]	The process or thread would have blocked, and either the <i>abstime</i> parameter specified a nanoseconds field value less than zero or greater than or equal to 1000 million, or the <i>sem_clockwait()</i> function was passed an invalid or unsupported <i>clock_id</i> value.			
59072	[ETIMEDOUT]	The semaphore could not be locked before the specified timeout expired.			
59073	The <i>sem_clockwai</i>	t() and <i>sem_timedwait</i> () functions may fail if:			
59074	[EDEADLK]	A deadlock condition was detected.			
59075	[EINTR]	A signal interrupted the function.			
59076	[EINVAL]	The <i>sem</i> argument does not refer to a valid semaphore.			

# sem\_timedwait()

#### **EXAMPLES** 59077 The program shown below operates on an unnamed semaphore. The program expects two 59078 command-line arguments. The first argument specifies a seconds value that is used to set an 59079 alarm timer to generate a SIGALRM signal. This handler performs a *sem\_post()* to increment the 59080 semaphore that is being waited on in *main()* using *sem\_clockwait()*. The second command-line 59081 argument specifies the length of the timeout, in seconds, for *sem\_clockwait()*. 59082 #include <unistd.h> 59083 #include <stdio.h> 59084 #include <stdlib.h> 59085 59086 #include <semaphore.h> #include <time.h> 59087 #include <assert.h> 59088 #include <errno.h> 59089 59090 #include <signal.h> sem\_t sem; 59091 static void 59092 handler(int sig) 59093 59094 { 59095 int sav\_errno = errno; static const char info\_msg[] = "sem\_post() from handler\n"; 59096 write(STDOUT\_FILENO, info\_msg, sizeof info\_msg - 1); 59097 if $(sem_post(\&sem) == -1)$ { 59098 static const char err\_msg[] = "sem\_post() failed\n"; 59099 59100 write(STDERR\_FILENO, err\_msg, sizeof err\_msg - 1); \_exit(EXIT\_FAILURE); 59101 59102 } 59103 errno = sav\_errno; } 59104 59105 int main(int argc, char \*argv[]) 59106 59107 { struct sigaction sa; 59108 struct timespec ts; 59109 int s; 59110 if (argc != 3) { 59111 fprintf(stderr, "Usage: %s <alarm-secs> <wait-secs>\n", 59112 argv[0]); 59113 exit(EXIT\_FAILURE); 59114 } 59115 if (sem\_init(&sem, 0, 0) == -1) { 59116 59117 perror("sem\_init"); exit(EXIT\_FAILURE); 59118 } 59119 59120 /\* Establish SIGALRM handler; set alarm timer using argv[1] \*/ sa.sa\_handler = handler; 59121 sigemptyset(&sa.sa\_mask); 59122 sa.sa\_flags = 0; 59123 if (sigaction(SIGALRM, &sa, NULL) == -1) { 59124

1802

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```
perror("sigaction");
59125
59126
                       exit(EXIT_FAILURE);
59127
                  }
                  alarm(atoi(argv[1]));
59128
                  /* Calculate relative interval as current time plus
59129
                      number of seconds given argv[2] */
59130
                  if (clock_gettime(CLOCK_MONOTONIC, &ts) == -1) {
59131
                       perror("clock_gettime");
59132
                       exit(EXIT_FAILURE);
59133
                  }
59134
                  ts.tv_sec += atoi(argv[2]);
59135
                  printf("main() about to call sem_clockwait()\n");
59136
                  while ((s = sem_clockwait(&sem, CLOCK_MONOTONIC, &ts)) ==
                                                                                        -1 &&
59137
                         errno == EINTR)
59138
                                           /* Restart if interrupted by handler */
                       continue;
59139
                  /* Check what happened */
59140
                  if (s == -1) {
59141
                       if (errno == ETIMEDOUT)
59142
                            printf("sem_clockwait() timed out\n");
59143
59144
                       else
                            perror("sem_clockwait");
59145
                  } else
59146
                       printf("sem_clockwait() succeeded\n");
59147
                  exit((s == 0) ? EXIT_SUCCESS : EXIT_FAILURE);
59148
             }
59149
     APPLICATION USAGE
59150
             Applications using these functions may be subject to priority inversion, as discussed in XBD
59151
             Section 3.260 (on page 66).
59152
     RATIONALE
59153
             None.
59154
     FUTURE DIRECTIONS
59155
59156
             None.
     SEE ALSO
59157
             sem_post(), sem_trywait(), semctl(), semget(), semop(), time()
59158
             XBD Section 3.260 (on page 66), <semaphore.h>, <time.h>
59159
     CHANGE HISTORY
59160
             First released in Issue 6. Derived from IEEE Std 1003.1d-1999.
59161
             IEEE Std 1003.1-2001/Cor 2-2004, item XSH/TC2/D6/120 is applied, updating the ERRORS
59162
             section so that the [EINVAL] error becomes optional.
59163
59164
     Issue 7
             The sem_timedwait() function is moved from the Semaphores option to the Base.
59165
             Functionality relating to the Timers option is moved to the Base.
59166
59167
             An example is added.
```

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# sig2str()

	NAME	
61660 61661		sig2str, str2sig — translate between signal names and numbers
61662	SYNOP	SIS
61663	CX	<pre>#include <signal.h></signal.h></pre>
61664 61665 61666		<pre>int sig2str(int signum, char *str); int str2sig(const char *restrict str, int *restrict pnum);</pre>
61667	DESCR	<b>PTION</b>
61668		The <i>sig2str()</i> function shall translate the signal number specified by <i>signum</i> to a signal name and
61669		shall store this string in the location specified by <i>str</i> . The application shall ensure that <i>str</i> points
61670		to a location that can store the string including the terminating null byte. The symbolic constant
61671		SIG2STR_MAX defined in <b><signal.h></signal.h></b> gives the maximum number of bytes required.
61672		If <i>signum</i> is equal to 0, the behavior is unspecified.
61673		If signum is equal to one of the symbolic constants listed in the table of signal numbers in
61674		<signal.h>, the stored signal name shall be the name of the symbolic constant without the SIG</signal.h>
61675		prefix.
61676		If <i>signum</i> is equal to SIGRTMIN or SIGRTMAX, the stored string shall be "RTMIN" or "RTMAX",
61677		respectively.
61678		If signum is between SIGRTMIN+1 and (SIGRTMIN+SIGRTMAX)/2 inclusive, the stored string
61679		shall be of the form "RTMIN+ $n$ ", where $n$ is the shortest decimal representation of the value of
61680		signum–SIGRTMIN.
61681		If <i>signum</i> is between (SIGRTMIN+SIGRTMAX)/2 + 1 and SIGRTMAX–1 inclusive, the stored
61682		string shall be either of the form "RTMIN+ <i>n</i> " or of the form "RTMAX– <i>m</i> ", where <i>n</i> is the shortest
61683		decimal representation of the value of <i>signum</i> –SIGRTMIN and <i>m</i> is the shortest decimal
61684		representation of the value of SIGRTMAX-signum.
61685		If signum is a valid, supported signal number, is either less than SIGRTMIN or greater than
61686		SIGRTMAX, and is not equal to one of the symbolic constants listed in the table of signal
61687		numbers in <b><signal.h></signal.h></b> , the stored string shall uniquely identify the signal number <i>signum</i> in an
61688		unspecified manner.
61689		The <i>str2sig()</i> function shall translate the signal name in the string pointed to by <i>str</i> to a signal
61690		number and shall store this value in the location specified by <i>pnum</i> .
61691		If <i>str</i> points to a string containing the name of one of the symbolic constants listed in the table of
61692		signal numbers in <b><signal.h></signal.h></b> , without the <b>SIG</b> prefix, the stored signal number shall be equal to
61693		the value of the symbolic constant.
61694		If str points to the string "RTMIN" or "RTMAX", the stored value shall be equal to SIGRTMIN or
61695		SIGRTMAX, respectively.
61696		If <i>str</i> points to a string of the form "RTMIN+ $n$ ", where $n$ is a decimal representation of a number
61697		between 1 and SIGRTMAX-SIGRTMIN-1 inclusive, the stored value shall be equal to
61698		SIGRTMIN+n.
61699		If str points to a string of the form "RTMAX $-n$ ", where <i>n</i> is a decimal representation of a number
61700		between 1 and SIGRTMAX-SIGRTMIN-1 inclusive, the stored value shall be equal to
61701		SIGRTMAX-n.
(1702		If str points to a string containing a desimal representation of a valid supported sized number
01/02 61703		the value stored in the location pointed to by <i>nnum</i> shall be equal to that number
51705		are value stored in the location pointed to by prain shall be equal to that humber.

- 61704 If *str* points to a string containing a decimal representation of the value 0 and the string was not 61705 returned by a previous successful call to *sig2str()* with a *signum* argument of 0, the behavior is 61706 unspecified.
- 61707 If *str* points to a string returned by a previous successful call to *sig2str(signum,str)*, the value 61708 stored in the location pointed to by *pnum* shall be equal to *signum*.
- 61709 If *str* points to a string that does not meet any of the above criteria, *str2sig*() shall store a value in 61710 the location pointed to by *pnum* if and only if it recognizes the string as an additional 61711 implementation-dependent form of signal name.

### 61712 **RETURN VALUE**

- 61713If signum is a valid, supported signal number (that is, one for which kill() does not return -161714with errno set to [EINVAL]), the sig2str() function shall return 0; otherwise, if signum is not equal61715to 0, it shall return -1.
- 61716 If *str2sig*() stores a value in the location pointed to by *pnum*, it shall return 0; otherwise, it shall 61717 return –1.

### 61718 ERRORS

61719 No errors are defined.

### 61720 EXAMPLES

61721 None.

### 61722 APPLICATION USAGE

61723 None.

### 61724 RATIONALE

61725Historical versions of these functions translated a *signum* value 0 to "EXIT" (and vice versa), so61726that they could be used by the shell for the *trap* utility. When adding the functions to this61727standard, the standard developers felt that they should be aimed at more general-purpose use,61728and consequently requiring this behavior did not seem appropriate and so the behavior in this61729case has been made unspecified.

### 61730 FUTURE DIRECTIONS

61731 None.

### 61732 SEE ALSO

61733 *kill(), sigaction(), strsignal()* 

### 61734 XBD <signal.h>

### 61735 CHANGE HISTORY

61736 First released in Issue 8.

# str2sig()

## NAME

str2sig — translate between signal names and numbers

### 63653 63654 **SYNOPSIS**

63655	СХ	#inc	clude <signal.< th=""><th>1&gt;</th><th></th><th></th><th></th><th></th><th></th></signal.<>	1>					
63656		int	str2sig(const	char	*restrict	str,	int	*restrict	pnum);

63657

63658 **DESCRIPTION** 

Refer to *sig2str()*.

63659

Sauth

# strlcat()

( 10 71	NAME
64871 64872	strlcat, strlcpy — size-bounded string concatenation and copying
64873	SYNOPSIS
64874	CX #include <string.h></string.h>
64075	size t strlast(abor *restrict det const abor *restrict are
64876	size t dstsize):
64877	size_t strlcpy(char *restrict <i>dst</i> , const char *restrict <i>src</i> ,
64878	size_t dstsize);
64879	
64880	DESCRIPTION
64881	The <i>strlcpy()</i> and <i>strlcat()</i> functions copy and concatenate strings, stopping when either a NUL
64882	terminator in the source string is encountered or the specified full size of the destination buffer is
64883	reached. They NUL terminate the result if there is room. The application should ensure that
64884	room for the NUL terminator is included in <i>dstsize</i> .
64885	The <i>strlcpy()</i> function shall copy not more than <i>dstsize</i> – 1 bytes from the string pointed to by <i>src</i>
64886	to the array pointed to by <i>dst</i> ; a NUL byte in <i>src</i> and bytes that follow it shall not be copied. A
64887	terminating NUL byte shall be appended to the result, unless <i>dstsize</i> is 0. If copying takes place
64888	between objects that overlap, the benavior is undefined.
64889	The $strlcat()$ function shall append not more than $dstsize - strlen(dst) - 1$ bytes from the string
64890	pointed to by <i>src</i> to the end of the string pointed to by <i>dst</i> ; a NUL byte in <i>src</i> and bytes that
64891	follow it shall not be appended. The initial byte of <i>src</i> shall overwrite the NUL byte at the end of det. A terminating NUL byte shall be appended to the result upless its location would be at or
64892 64893	dst + $dstsize$ If conving takes place between objects that overlap, the behavior is
64894	undefined.
64895	The <i>strlcpy()</i> and <i>strlcat()</i> functions shall not change the setting of <i>errno</i> on valid input.
(190)	
64897	Upon successful completion, the $strlcpu()$ function shall return the length of the string pointed to
64898	by <i>src</i> ; that is, the number of bytes in the string, not including the terminating NUL byte.
64899	Upon successful completion, the <i>strlcat()</i> function shall return the initial length of the string
64900	pointed to by usi plus the length of the string pointed to by src.
64901	No return values are reserved to indicate an error.
64902	ERRORS
64903	No errors are defined.
64904	EXAMPLES
64905	The following example detects truncation while combining a path prefix (including trailing
64906	<slash>) and a filename to produce a portable pathname:</slash>
64907	<pre>char *prefix, *filenam, pathnam[_POSIX_PATH_MAX];</pre>
64908	if (strlcpy(pathnam, prefix, sizeof pathnam) >= sizeof pathnam
64909	<pre>stricat(pathnam, filenam, sizeof pathnam) &gt;= sizeof pathnam) </pre>
04910 64911	1 // truncation occurred
64912	···
64913	}
64914	This code ensures there is room for the NUL terminator by
~ -/ 17	

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strlcat()

64915	• Calling <i>strlcpy()</i> with a non-zero <i>dstsize</i> argument.
64916	• Only calling <i>strlcat()</i> if the return value of <i>strlcpy()</i> indicated that truncation did not occur.
64917 64918 64919 64920	APPLICATION USAGE The return value of the <i>strlcpy()</i> and <i>strlcat()</i> functions follows the same convention as <i>snprintf()</i> ; that is, they return the total length of the string they tried to create. If the return value is greater than or equal to <i>dstsize</i> , the output string has been truncated.
64921 64922	RATIONALE None.
64923 64924	FUTURE DIRECTIONS None.
64925 64926	SEE ALSO fprintf(), strlen(), strncat(), strncpy(), wcslcat()
64927	XBD <string.h></string.h>
64928 64929 64930	CHANGE HISTORY First released in Issue 8.

### NAME 71081

wcslcat, wcslcpy — size-bounded wide string concatenation and copying

### 71082 **SYNOPSIS** 71083

71084
71085
71086

CX #include <wchar.h> size\_t wcslcat(wchar\_t \*restrict *dst*, const wchar\_t \*restrict *src*, 35 size\_t dstsize); 36 size\_t wcslcpy(wchar\_t \*restrict dst, const wchar\_t \*restrict src, 71087 size\_t dstsize); 71088

71089

#### DESCRIPTION 71090

The *wcslcpy()* and *wcslcat()* functions copy and concatenate wide strings, stopping when either a 71091 terminating null wide-character code in the source wide string is encountered or the specified 71092 full size (in wide-character codes) of the destination buffer is reached. They null terminate the 71093 result if there is room. The application should ensure that room for the terminating null wide-71094 character code is included in dstsize. 71095

The wcslcpy() function shall copy not more than dstsize - 1 wide-character codes from the wide 71096 string pointed to by *src* to the array pointed to by *dst*; a terminating null wide-character code in 71097 src and wide-character codes that follow it shall not be copied. A terminating null wide-71098 character code shall be appended to the result, unless *dstsize* is 0. If copying takes place between 71099 objects that overlap, the behavior is undefined. 71100

The wcslcat() function shall append not more than dstsize - wcslen(dst) - 1 wide-character codes 71101 from the wide string pointed to by src to the end of the wide string pointed to by dst; a 71102 terminating null wide-character code in src and wide-character codes that follow it shall not be 71103 appended. The initial wide-character code of *src* shall overwrite the null wide-character code at 71104 the end of *dst*. A terminating null wide-character code shall be appended to the result, unless its 71105 location would be at or beyond dst + dstsize. If copying takes place between objects that overlap, 71106 the behavior is undefined. 71107

The *wcslcpy()* and *wcslcat()* functions shall not change the setting of *errno* on valid input. 71108

#### **RETURN VALUE** 71109

- Upon successful completion, the *wcslcpy()* function shall return the length of the wide string 71110 pointed to by src; that is, the number of wide-character codes in the wide string, not including 71111 the terminating null wide-character code. 71112
- 71113 Upon successful completion, the *wcslcat()* function shall return the initial length of the wide string pointed to by *dst* plus the length of the wide string pointed to by *src*. 71114
- No return values are reserved to indicate an error. 71115

#### ERRORS 71116

No errors are defined. 71117

#### **EXAMPLES** 71118

71119 None.

#### APPLICATION USAGE 71120

The return value of the wcslcpy() and wcslcat() functions follows the same convention as 71121 71122 *snprintf()*; that is, they return the total length (in wide-character codes) of the wide string they tried to create. If the return value is greater than or equal to *dstsize*, the output wide string has 71123 been truncated. 71124

# wcslcat()

System Interfaces

71125 71126	RATIONALE None.
71127 71128	FUTURE DIRECTIONS None.
71129 71130	<b>SEE ALSO</b> <i>fprintf(), strlcat(), wcslen(), wcsncat(), wcsncpy()</i>
71131	XBD <b><wchar.h></wchar.h></b>
71132 71133 71134	CHANGE HISTORY First released in Issue 8.

121068POSIX.1b is a software, source-level standard and most of the benefits of the alternate121069representation are enjoyed by hardware implementations of clocks and algorithms. It was121070felt that mandating this format for POSIX.1b clocks and timers would unnecessarily121071burden the application developer with writing, possibly non-portable, multiple precision121072arithmetic packages to perform conversion between binary fractions and integral units121073such as nanoseconds, milliseconds, and so on.

### 121074 Rationale for the Monotonic Clock

For those applications that use time services to achieve realtime behavior, changing the value of 121075 the clock on which these services rely may cause erroneous timing behavior. For these 121076 121077 applications, it is necessary to have a monotonic clock which cannot run backwards, and which has a maximum clock jump that is required to be documented by the implementation. 121078 Additionally, it is desirable (but not required by POSIX.1-202x) that the monotonic clock 121079 increases its value uniformly. This clock should not be affected by changes to the system time; 121080 for example, to synchronize the clock with an external source or to account for leap seconds. 121081 Such changes would cause errors in the measurement of time intervals for those time services 121082 that use the absolute value of the clock. 121083

- 121084 One could argue that by defining the behavior of time services when the value of a clock is changed, deterministic realtime behavior can be achieved. For example, one could specify that 121085 relative time services should be unaffected by changes in the value of a clock. However, there are 121086 time services that are based upon an absolute time, but that are essentially intended as relative 121087 time services. For example, *pthread\_cond\_timedwait()* uses an absolute time to allow it to wake 121088 after the required interval despite spurious wakeups. Although sometimes the up 121089 *pthread\_cond\_timedwait()* timeouts are absolute in nature, there are many occasions in which they 121090 are relative, and their absolute value is determined from the current time plus a relative time 121091 interval. In this latter case, if the clock changes while the thread is waiting, the wait interval will 121092 not be the expected length. If a *pthread\_cond\_timedwait()* function were created that would take a 121093 relative time, it would not solve the problem because to retain the intended "deadline" a thread 121094 would need to compensate for latency due to the spurious wakeup, and preemption between 121095 wakeup and the next wait. 121096
- 121097The solution is to create a new monotonic clock, whose value does not change except for the121098regular ticking of the clock, and use this clock for implementing the various relative timeouts121099that appear in the different POSIX interfaces, as well as allow *pthread\_cond\_timedwait()* to choose121100this new clock for its timeout. A new *clock\_nanosleep()* function is created to allow an application121101to take advantage of this newly defined clock. Notice that the monotonic clock may be121102implemented using the same hardware clock as the system clock.
- 121103Relative timeouts for *sigtimedwait()* and *aio\_suspend()* have been redefined to use the monotonic121104clock, if present. The *alarm()* function has not been redefined, because the same effect but with121105better resolution can be achieved by creating a timer (for which the appropriate clock may be121106chosen).
- The *pthread\_cond\_timedwait()* function has been treated in a different way, compared to other 121107 functions with absolute timeouts, because it is used to wait for an event, and thus it may have a 121108 deadline, while the other timeouts are generally used as an error recovery mechanism, and for 121109 121110 them the use of the monotonic clock is not so important. Since the desired timeout for the *pthread\_cond\_timedwait()* function may either be a relative interval or an absolute time of day 121111 121112 deadline, a new initialization attribute has been created for condition variables to specify the clock that is used for measuring the timeout in a call to *pthread\_cond\_timedwait()*. In this way, if 121113 a relative timeout is desired, the monotonic clock will be used; if an absolute deadline is 121114 required instead, the CLOCK\_REALTIME or another appropriate clock may be used. For 121115 condition variables, this capability is also available by passing CLOCK\_MONOTONIC to the 121116

121117pthread\_cond\_clockwait() function. Similarly, CLOCK\_MONOTONIC can be specified when121118calling pthread\_mutex\_clocklock(), pthread\_rwlock\_clockrdlock(), pthread\_rwlock\_clockwrlock(), and121119sem\_clockwait().

121120 It was later found necessary to add variants of almost all interfaces that accept absolute timeouts that allow the clock to be specified. This is because, despite the claim in the previous paragraph, 121121 it is not possible to safely use a CLOCK\_REALTIME absolute timeout even to prevent errors 121122 when the system clock is warped by a potentially large amount. A "safety timeout" of a minute 121123 on a call to *pthread\_mutex\_timedlock(*) could actually mean that the call would return 121124 ETIMEDOUT early without acquiring the lock if the system clock is warped forwards 121125 immediately prior to or during the call. On the other hand, a short timeout could end up being 121126 arbitrarily long if the system clock is warped backwards immediately prior to or during the call. 121127 These problems are solved by the new *clockwait* and *clocklock* variants of the existing *timedwait* 121128 and *timedlock* functions. These variants accept an extra **clockid\_t** parameter to indicate the clock 121129 to be used for the wait. The clock ID is passed rather than using attributes as previously for 121130 *pthread\_cond\_timedwait()* in order to allow the ISO/IEC 14882: 2011 standard (C++11) and later to 121131 be implemented correctly. C++ requires that the clock to use for the wait is not known until the 121132 time of the wait call, so it cannot be supplied during creation. The new functions are 121133 pthread\_cond\_clockwait(), pthread\_mutex\_clocklock(), pthread\_mutex\_clockrdlock(), 121134 pthread mutex\_clockwrlock(), and sem\_clockwait(). It is expected that mq\_clockreceive() and 121135 *mq\_clocksend()* functions will be added in a future version of this standard. 121136

121137The nanosleep() function has not been modified with the introduction of the monotonic clock.121138Instead, a new clock\_nanosleep() function has been created, in which the desired clock may be121139specified in the function call.

• History of Resolution Issues

121141Due to the shift from relative to absolute timeouts in IEEE Std 1003.1d-1999, the121142amendments to the sem\_timedwait(), pthread\_mutex\_timedlock(), mq\_timedreceive(), and121143mq\_timedsend() functions of that standard have been removed. Those amendments121144specified that CLOCK\_MONOTONIC would be used for the (relative) timeouts if the121145Monotonic Clock option was supported.

121146Having these functions continue to be tied solely to CLOCK\_MONOTONIC would not121147work. Since the absolute value of a time value obtained from CLOCK\_MONOTONIC is121148unspecified, under the absolute timeouts interface, applications would behave differently121149depending on whether the Monotonic Clock option was supported or not (because the121150absolute value of the clock would have different meanings in either case).

- 121151 Two options were considered:
  - 1. Leave the current behavior unchanged, which specifies the CLOCK\_REALTIME clock for these (absolute) timeouts, to allow portability of applications between implementations supporting or not the Monotonic Clock option.
- 1211552. Modify these functions in the way that *pthread\_cond\_timedwait()* was modified to121156allow a choice of clock, so that an application could use CLOCK\_REALTIME when121157it is trying to achieve an absolute timeout and CLOCK\_MONOTONIC when it is121158trying to achieve a relative timeout.
- 121159It was decided that the features of CLOCK\_MONOTONIC are not as critical to these121160functions as they are to *pthread\_cond\_timedwait()*. The *pthread\_cond\_timedwait()* function is121161given a relative timeout; the timeout may represent a deadline for an event. When these121162functions are given relative timeouts, the timeouts are typically for error recovery121163purposes and need not be so precise.
- 121164

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121154

Therefore, it was decided that these functions should be tied to CLOCK\_REALTIME and

122004	Supported Threads Functions
122005	On POSIX-conforming systems, the following symbolic constants are always conforming:
122006	_POSIX_READER_WRITER_LOCKS
122007	_POSIX_THREADS

122008 Therefore, the following threads functions are always supported:

122009	pthread_atfork()	pthread_kill()
122010	pthread_attr_destroy()	pthread_mutex_destroy()
122011	pthread_attr_getdetachstate()	pthread_mutex_init()
122012	pthread_attr_getguardsize()	pthread_mutex_lock()
122013	pthread_attr_getschedparam()	pthread_mutex_trylock()
122014	pthread_attr_init()	pthread_mutex_unlock()
122015	pthread_attr_setdetachstate()	pthread_mutexattr_destroy()
122016	pthread_attr_setguardsize()	pthread_mutexattr_getpshared()
122017	pthread_attr_setschedparam()	pthread_mutexattr_gettype()
122018	pthread_cancel()	pthread_mutexattr_init()
122019	pthread_cleanup_pop()	pthread_mutexattr_setpshared()
122020	pthread_cleanup_push()	pthread_mutexattr_settype()
122021	pthread_cond_broadcast()	pthread_once()
122022	pthread_cond_clockwait()	pthread_rwlock_destroy()
122023	pthread_cond_destroy()	pthread_rwlock_init()
122024	pthread_cond_init()	pthread_rwlock_rdlock()
122025	pthread_cond_signal()	pthread_rwlock_tryrdlock()
122026	pthread_cond_timedwait()	pthread_rwlock_trywrlock()
122027	pthread_cond_wait()	pthread_rwlock_unlock()
122028	pthread_condattr_destroy()	pthread_rwlock_wrlock()
122029	pthread_condattr_getpshared()	pthread_rwlockattr_destroy()
122030	pthread_condattr_init()	pthread_rwlockattr_getpshared()
122031	pthread_condattr_setpshared()	pthread_rwlockattr_init()
122032	pthread_create()	pthread_rwlockattr_setpshared()
122033	pthread_detach()	pthread_self()
122034	pthread_equal()	pthread_setcancelstate()
122035	pthread_exit()	pthread_setcanceltype()
122036	pthread_getspecific()	pthread_setspecific()
122037	pthread_join()	pthread_sigmask()
122038	pthread_key_create()	pthread_testcancel()
122039	pthread_key_delete()	sigwait()
		-

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122084	pthread_mutex_lock()
122085	pthread_mutex_trylock()
122086	pthread_mutex_unlock()
122087 122088	to take account of the new mutex attribute type and to specify behavior which was declared as undefined in POSIX.1c. How a calling thread acquires or releases a mutex now
122089	The ture attribute can have the following values:
122090	The type attribute can have the following values.
122091	PTHREAD_MUTEX_NORMAL
122092	Basic mutex with no specific error checking built in. Does not report a deadlock error.
122093	PTHREAD_MUTEX_RECURSIVE
122094	Allows any thread to recursively lock a mutex. The mutex must be unlocked an equal
122095	number of times to release the mutex.
122096	PTHREAD_MUTEX_ERRORCHECK
122097	Detects and reports simple usage errors; that is, an attempt to unlock a mutex that is
122098	not locked by the calling thread or that is not locked at all, or an attempt to relock a
122099	mutex the thread already owns.
122100	PTHREAD_MUTEX_DEFAULT
122101	The default mutex type. May be mapped to any of the above mutex types or may be
122102	an implementation-defined type.
122103 122104 122105 122106 122107	<i>Normal</i> mutexes do not detect deadlock conditions; for example, a thread will hang if it tries to relock a normal mutex that it already owns. Attempting to unlock a mutex locked by another thread, or unlocking an unlocked mutex, results in undefined behavior. Normal mutexes will usually be the fastest type of mutex available on a platform but provide the least error checking.
122108	<i>Recursive</i> mutexes are useful for converting old code where it is difficult to establish clear
122109	boundaries of synchronization. A thread can relock a recursive mutex without first
122110	unlocking it. The relocking deadlock which can occur with normal mutexes cannot occur
122111	with this type of mutex. However, multiple locks of a recursive mutex require the same
122112	number of unlocks to release the mutex before another thread can acquire the mutex.
122113	Furthermore, this type of mutex maintains the concept of an owner. Thus, a thread
122114	attempting to unlock a recursive mutex which another thread has locked returns with an
122115	error. A thread attempting to unlock a recursive mutex that is not locked returns with an
122116	error. Never use a recursive mutex with condition variables because the implicit unlock
122117	performed by <i>pthread_cond_clockwait()</i> , <i>pthread_cond_timedwait()</i> , or <i>pthread_cond_wait()</i>
122118	will not actually release the mutex if it had been locked multiple times.
122119 122120 122121 122122 122123 122124 122125	<i>Errorcheck</i> mutexes provide error checking and are useful primarily as a debugging aid. A thread attempting to relock an errorcheck mutex without first unlocking it returns with an error. Again, this type of mutex maintains the concept of an owner. Thus, a thread attempting to unlock an errorcheck mutex which another thread has locked returns with an error. A thread attempting to unlock an errorcheck mutex that is not locked also returns with an error. It should be noted that errorcheck mutexes will almost always be much slower than normal mutexes due to the extra state checks performed.
122126	The default mutex type provides implementation-defined error checking. The default
122127	mutex may be mapped to one of the other defined types or may be something entirely
122128	different. This enables each vendor to provide the mutex semantics which the vendor feels
122129	will be most useful to their target users. Most vendors will probably choose to make
122130	normal mutexes the default so as to give applications the benefit of the fastest type of

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### 126192 Unsatisfied Requirements

Detailed control of common device classes, specifically magnetic tape, is not provided.

### 126194 **D.2.5** Bounded (Realtime) Response

- 126195The realtime signal functions sigqueue(), sigtimedwait(), and sigwaitinfo() provide queued signals126196and the prioritization of the handling of signals.
- 126197The SCHED\_FIFO, SCHED\_SPORADIC, and SCHED\_RR scheduling policies provide control126198over processor allocation.
- 126199The semaphore functions sem\_clockwait(), sem\_close(), sem\_destroy(), sem\_getvalue(), sem\_init(), +126200sem\_open(), sem\_post(), sem\_timedwait(), sem\_trywait(), sem\_unlink(), and sem\_wait() provide126201high-performance synchronization.
- 126202The memory management functions provide memory locking for control of memory allocation,126203file mapping for high performance, and shared memory for high-performance interprocess126204communication. The Message Passing option provides for interprocess communication without126205being dependent on shared memory.
- 126206The timers functions clock\_getres(), clock\_gettime(), clock\_settime(), nanosleep(), timer\_create(),126207timer\_delete(), timer\_getoverrun(), timer\_gettime(), and timer\_settime() provide functionality to126208manipulate clocks and timers and include a high resolution function called nanosleep() with a126209finer resolution than the sleep() function.
- The timeout functions pthread\_mutex\_clocklock(), pthread\_mutex\_timedlock(), 126210 pthread\_rwlock\_clockrdlock(), pthread\_rwlock\_clockwrlock(), pthread\_rwlock\_timedrdlock(), 126211 pthread\_rwlock\_timedwrlock(), sem\_clockwait(), and sem\_timedwait() — the Typed Memory + 126212 126213 Objects option and the Monotonic Clock option provide further facilities for applications to use to obtain predictable bounded response. 126214

### 126215 D.2.6 Operating System-Dependent Profile

POSIX.1-202x makes no distinction between text and binary files. The values of EXIT\_SUCCESSand EXIT\_FAILURE are further defined.

### 126218 Unsatisfied Requirements

126219None known, but the ISO C standard may contain some additional options that could be<br/>specified.126220specified.

### 126221 D.2.7 I/O Interaction

126222POSIX.1-202x defines how each of the ISO C standard *stdio* functions interact with the POSIX.1126223operations, typically specifying the behavior in terms of POSIX.1 operations.

126793 126794 126795	<pre>rintl(), round(), roundf(), roundl(), scalbln(), scalblnf(), scalblnl(), scalbn(), scalbnf(), scalbnl(), signbit(), sin(), sinf(), sinh(), sinhf(), sinhl(), sinl(), sqrt(), sqrtf(), sqrtl(), tan(), tanf(), tanh(), tanhf(), tanhl(), tanl(), tgamma(), tgammaf(), tgammal(), trunc(), truncf(),</pre>	
126796	truncl()	
126707	POSIX C I ANG SUPPORT Conceal ISO C Library	
126797	abs() asctime() atof() atoi() atol() atol() bearch() calloc() ctime() difftime() div()	
126799	feclearexcent() feaetenzi() feaetenzicentflag() feaetround() feholdexcent() feraiseexcent()	
126799	fesetenzy(), fesetexcentflag(), fesetround(), fetestexcent(), feundateenzy(), free(), ontime()	
126801	imarahs() imardin() isalnum() isalnha() ishlank() iscntrl() isdioit() isoranh() islomer()	
126802	ismint() ismunct() issnace() isumer() isrdigit() labs() ldig(() llabs() lldig() localecong()	
126803	localtime() malloc() memchr() memcmn() memcmu() memmoze() memset() mktime()	
126804	asort() rand() realloc() setlocale() snnrintf() snrintf() srand() sscanf() strcat() strchr()	
126805	strcmn() strcoll() strcmu() strcsnn() strerror() strftime() strlen() strncat() strncmn()	
126806	strncny() strphrk() strrchr() strsnn() strstr() strtod() strtof() strtoimax() strtok() strtol()	
126807	strtold() strtoll() strtoul() strtoull() strtoumax() strxfrm() time() tolozoer() toumer()	
126808	tzname, tzset(), va arg(), va conv(), va end(), va start(), vsnrintf(), vsnrintf(), vsscanf()	
120000		
126809	POSIX_C_LANG_SUPPORT_R: Thread-Safe General ISO C Library	
126810	asctime_r(), ctime_r(), gmtime_r(), localtime_r(), qsort_r(), strerror_r(), strtok_r()	+
126811	POSIX_C_LANG_WIDE_CHAR: Wide-Character ISO C Library	
126812	<pre>btowc(), iswalnum(), iswalpha(), iswblank(), iswcntrl(), iswctype(), iswdigit(), iswgraph(),</pre>	
126813	iswlower(), iswprint(), iswpunct(), iswspace(), iswupper(), iswxdigit(), mblen(), mbrlen(),	
126814	mbrtowc(), mbsinit(), mbsrtowcs(), mbstowcs(), mbtowc(), swprintf(), swscanf(), towctrans(),	
126815	towlower(), towupper(), vswprintf(), vswscanf(), wcrtomb(), wcscat(), wcschr(), wcscmp(),	
126816	wcscoll(), wcscpu(), wcscspn(), wcsftime(), wcslen(), wcsncat(), wcsncmp(), wcsncpu(),	
126817	wcspbrk(), wcsrchr(), wcsrtombs(), wcspn(), wcstr(), wcstod(), wcstof(), wcstof(), wcstoimax(),	
126818	westok() $westol()$ , $westold()$ $westoll()$ $westoll()$ $westoul()$ $westoul()$ , $westoul()$ , $westoumax()$	
126819	wester(), wester	
126820	wmemmove(), wereme(), wereme(), wereme(), werege(), wmemeni(), wmememp(), wmemepg(), wmemmove(). wmemset()	
	$PO(D) \subset LANC WIDE CHAD EVEE ( 1 1 WELCH ) ( 100 CL')$	
126821	POSIX_C_LANG_WIDE_CHAR_EX1: Extended Wide-Character ISO C Library	
126822	mbsnrtowcs(), wcpcpy(), wcpncpy(), wcscasecmp(), wcsdup(), wcslcat(), wcslcpy(),	+
126823	wcsncasecmp(), wcsnlen(), wcsnrtombs()	
126824	POSIX_C_LIB_EXT: General C Library Extension	
126825	fnmatch(), getentropy(), getopt(), getsubopt(), memmem(), optarg, opterr, optind, optopt,	+
126826	reallocarray(), stpcpy(), stpncpy(), strcasecmp(), strdup(), strfmon(), strlcat(), strlcpy(),	+
126827	strncasecmp(), strndup(), strnlen()	
10/000	POSIX CLOCK SELECTION: Clock Selection	
126828	clock nanoclam() nthread condattr getclock() nthread condattr setclock()	
126829	clock_nunosleep(), plnteuu_conuult1_gelclock(), plnteuu_conuult1_selclock()	
126830	POSIX_DEVICE_IO: Device Input and Output	
126831	FD_CLR(), FD_ISSET(), FD_SET(), FD_ZERO(), clearerr(), close(), fclose(), fdopen(), feof(),	
126832	ferror(), fflush(), fgetc(), fgets(), fileno(), fopen(), fprintf(), fputc(), fputs(), fread(), freopen(),	
126833	fscanf(), fwrite(), getc(), getchar(), open(), perror(), poll(), ppoll(), printf(), pread(), pselect(),	+
126834	<pre>putc(), putchar(), puts(), pwrite(), read(), scanf(), select(), setbuf(), setvbuf(), stderr, stdin,</pre>	
126835	<pre>stdout, ungetc(), vfprintf(), vfscanf(), vprintf(), vscanf(), write()</pre>	
126836	POSIX DEVICE IO EXT: Extended Device Input and Output	
120030	dnrintf() fmemonen() onen memstream() vdnrintf()	
12003/	$u_{p}(m_{j}(j,j))$ $(j,j)$	
126838	POSIX_DEVICE_SPECIFIC: General Terminal	
126839	cfgetispeed(), cfgetospeed(), cfsetispeed(), cfsetospeed(), ctermid(), isatty(), tcdrain(), tcflow(),	
126840	tcflush(), tcgetattr(), tcgetwinsize(), tcsendbreak(), tcsetattr(), tcsetwinsize(), ttyname()	

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126841 126842	POSIX_DEVICE_SPECIFIC_R: Thread-Safe General Terminal <i>ttyname_r()</i>	
126843 126844	<pre>POSIX_DYNAMIC_LINKING: Dynamic Linking</pre>	+
126845 126846 126847	<pre>POSIX_FD_MGMT: File Descriptor Management     dup(), dup2(), dup3(), fcntl(), fgetpos(), fseek(), fseeko(), fsetpos(), ftell(), ftello(), ftruncate(),     lseek(), rewind()</pre>	
126848 126849	POSIX_FIFO: FIFO mkfifo()	
126850 126851	POSIX_FIFO_FD: FIFO File Descriptor Routines <i>mkfifoat(), mknodat()</i>	
126852 126853	POSIX_FILE_ATTRIBUTES: File Attributes <i>chmod(), chown(), fchmod(), fchown(), umask()</i>	
126854 126855	POSIX_FILE_ATTRIBUTES_FD: File Attributes File Descriptor Routines fchmodat(), fchownat()	
126856 126857 126858	POSIX_FILE_LOCKING: Thread-Safe Stdio Locking flockfile(), ftrylockfile(), funlockfile(), getc_unlocked(), getchar_unlocked(), putc_unlocked(), putchar_unlocked()	
126859 126860 126861 126862	<pre>POSIX_FILE_SYSTEM: File System     access(), chdir(), closedir(), creat(), fchdir(), fpathconf(), fstat(), fstatvfs(), getcwd(), link(),     mkdir(), mkostemp(), mkstemp(), opendir(), pathconf(), posix_getdents(), readdir(), remove(),     rename(), rewinddir(), rmdir(), stat(), statvfs(), tmpfile(), tmpnam(), truncate(), unlink()</pre>	+
126863 126864	<pre>POSIX_FILE_SYSTEM_EXT: File System Extensions</pre>	
126865 126866 126867	<pre>POSIX_FILE_SYSTEM_FD: File System File Descriptor Routines faccessat(), fdopendir(), fstatat(), linkat(), mkdirat(), openat(), renameat(), unlinkat(), utimensat()</pre>	
126868 126869	POSIX_FILE_SYSTEM_GLOB: File System Glob Expansion glob(), globfree()	
126870 126871	POSIX_FILE_SYSTEM_R: Thread-Safe File System readdir_r()	
126872 126873	POSIX_I18N: Internationalization catclose(), catgets(), catopen(), iconv(), iconv_close(), iconv_open(), nl_langinfo()	
126874 126875	<pre>POSIX_JOB_CONTROL: Job Control     setpgid(), tcgetpgrp(), tcgetsid()</pre>	
126876 126877	POSIX_MAPPED_FILES: Memory Mapped Files <i>mmap(), munmap()</i>	
126878 126879	POSIX_MEMORY_PROTECTION: Memory Protection <i>mprotect()</i>	
126880 126881 126882 126883 126884	<pre>POSIX_MULTI_CONCURRENT_LOCALES: Multiple Concurrent Locales duplocale(), freelocale(), getlocalename_l(), isalnum_l(), isalpha_l(), isblank_l(), iscntrl_l(), isdigit_l(), isgraph_l(), islower_l(), isprint_l(), ispunct_l(), isspace_l(), isupper_l(), iswalnum_l(), iswalpha_l(), iswblank_l(), iswcntrl_l(), iswctype_l(), iswdigit_l(), iswgraph_l(), iswlower_l(), iswprint_l(), iswpunct_l(), iswspace_l(), iswupper_l(), iswxdigit_l(), isxdigit_l(),</pre>	+

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126885 126886 126887	<pre>newlocale(), strcasecmp_l(), strcoll_l(), strfmon_l(), strncasecmp_l(), strxfrm_l(), tolower_l(), toupper_l(), towctrans_l(), towlower(), towupper(), uselocale(), wcscasecmp_l(), wcscoll_l(), wcsncasecmp_l(), wcsxfrm_l(), wctrans_l(), wctype_l()</pre>	
126888 126889 126890 126891	POSIX_MULTI_PROCESS: Multiple Processes _Exit(), _exit(), assert(), atexit(), clock(), execl(), execle(), execlp(), execv(), execvp(), exit(), fork(), getpgrp(), getpgid(), getpid(), getppid(), getsid(), setsid(), sleep(), times(), wait(), waitid(), waitpid()	
126892 126893	POSIX_MULTI_PROCESS_FD: Multiple Processes File Descriptor Routines fexecve()	
126894 126895 126896 126897 126898 126899 126900 126901 126901 126902 126903	<pre>POSIX_NETWORKING: Networking     accept(), accept4(), bind(), connect(), endhostent(), endnetent(), endprotoent(), endservent(),     freeaddrinfo(), gai_strerror(), getaddrinfo(), gethostent(), gethostname(), getnameinfo(),     getnetbyaddr(), getnetbyname(), getnetent(), getpeername(), getprotobyname(),     getprotobynumber(), getprotoent(), getservbyname(), getservbyport(), getservent(),     getsockname(), getsockopt(), htonl(), htons(), if_freenameindex(), if_indextoname(),     if_nameindex(), if_nametoindex(), inet_addr(), inet_ntoa(), inet_pton(), listen(),     ntohl(), ntohs(), recv(), recvfrom(), recvmsg(), send(), sendmsg(), sendto(), sethostent(),     setnetent(), setservent(), setsockopt(), shutdown(), socket(), sockatmark(),     socketpair()</pre>	
126904 126905	POSIX_PIPE: Pipe pipe(), pipe2()	
126906 126907	POSIX_ROBUST_MUTEXES: Robust Mutexes pthread_mutex_consistent(), pthread_mutexattr_getrobust(), pthread_mutexattr_setrobust()	
126908 126909	POSIX_REALTIME_SIGNALS: Realtime Signals sigqueue(), sigtimedwait(), sigwaitinfo()	
126910 126911	POSIX_REGEXP: Regular Expressions regcomp(), regerror(), regexec(), regfree()	
126912 126913 126914 126915 126916 126917	POSIX_RW_LOCKS: Reader Writer Locks pthread_rwlock_clockrdlock(), pthread_rwlock_clockwrlock(), pthread_rwlock_destroy(), pthread_rwlock_init(), pthread_rwlock_rdlock(), pthread_rwlock_timedrdlock(), pthread_rwlock_timedwrlock(), pthread_rwlock_tryrdlock(), pthread_rwlock_trywrlock(), pthread_rwlock_unlock(), pthread_rwlock_wrlock(), pthread_rwlockattr_destroy(), pthread_rwlockattr_init(), pthread_rwlockattr_getpshared(), pthread_rwlockattr_setpshared()	+
126918 126919 126920	<pre>POSIX_SEMAPHORES: Semaphores     sem_clockwait(), sem_close(), sem_destroy(), sem_getvalue(), sem_init(), sem_open(),     sem_post(), sem_timedwait(), sem_trywait(), sem_unlink(), sem_wait()</pre>	+
126921 126922	<pre>POSIX_SHELL_FUNC: Shell and Utilities pclose(), popen(), system(), wordexp(), wordfree()</pre>	
126923 126924	POSIX_SIGNAL_JUMP: Signal Jump Functions siglongjmp(), sigsetjmp()	
126925 126926 126927	POSIX_SIGNALS: Signals abort(), alarm(), kill(), pause(), raise(), sigaction(), sigaddset(), sigdelset(), sigemptyset(), sigfillset(), sigismember(), signal(), sigpending(), sigprocmask(), sigsuspend(), sigwait()	
126928 126929	<pre>POSIX_SIGNALS_EXT: Extended Signals     psignal(), psiginfo(), sig2str(), str2sig(), strsignal()</pre>	+

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126930 126931	<pre>POSIX_SINGLE_PROCESS: Single Process confstr(), environ, errno, getenv(), setenv(), sysconf(), uname(), unsetenv()</pre>
126932	POSIX_SPIN_LOCKS: Spin Locks
126933	pthread_spin_destroy(), pthread_spin_init(), pthread_spin_lock(), pthread_spin_trylock(),
126934	pthread_spin_unlock()
126935	POSIX_SYMBOLIC_LINKS: Symbolic Links
126936	lchown(), <sup>11</sup> lstat(), readlink(), symlink()
126937	POSIX_SYMBOLIC_LINKS_FD: Symbolic Links File Descriptor Routines
126938	readlinkat(), symlinkat()
126939	POSIX_SYSTEM_DATABASE: System Database
126940	getgrgid(), getgrnam(), getpwnam(), getpwuid()
126941 126942	<pre>POSIX_SYSTEM_DATABASE_R: Thread-Safe System Database     getgrgid_r(), getgrnam_r(), getpwnam_r(), getpwuid_r()</pre>
126943 126944 126945 126946 126947 126948 126949 126950 126951 126952 126953 126954 126955	<pre>POSIX_THREADS_BASE: Base Threads     pthread_attr_getschedparam(), pthread_attr_getdetachstate(),     pthread_attr_getschedparam(), pthread_attr_init(), pthread_attr_setdetachstate(),     pthread_attr_setschedparam(), pthread_cancel(), pthread_cleanup_pop(), pthread_cleanup_push(),     pthread_cond_broadcast(), pthread_cond_clockwait(), pthread_cond_destroy(),     pthread_cond_init(), pthread_cond_signal(), pthread_cond_timedwait(), pthread_cond_wait(),     pthread_condattr_destroy(), pthread_condattr_init(), pthread_create(), pthread_detach(),     pthread_equal(), pthread_exit(), pthread_getspecific(), pthread_join(), pthread_key_create(),     pthread_key_delete(), pthread_kill(), pthread_mutex_clocklock(), pthread_mutex_destroy(),     pthread_mutex_init(), pthread_mutex_lock(), pthread_mutex_timedlock(),     pthread_mutex_trylock(), pthread_mutex_unlock(), pthread_setcancelstate(),     pthread_setcanceltype(), pthread_setspecific(), pthread_sigmask(), pthread_testcancel()</pre>
126956	POSIX_THREADS_EXT: Extended Threads
126957	pthread_attr_getguardsize(), pthread_attr_setguardsize(), pthread_mutexattr_gettype(),
126958	pthread_mutexattr_settype()
126959	POSIX_TIMERS: Timers
126960	clock_getres(), clock_gettime(), clock_settime(), nanosleep(), timer_create(), timer_delete(),
126961	timer_getoverrun(), timer_gettime(), timer_settime()
126962 126963 126964	<pre>POSIX_USER_GROUPS: User and Group getegid(), geteuid(), getgid(), getgroups(), getlogin(), getuid(), seteuid(), setgid(), setuid()</pre>
126965	POSIX_USER_GROUPS_R: Thread-Safe User and Group
126966	getlogin_r()
126967 126968 126969 126970	<pre>POSIX_WIDE_CHAR_DEVICE_IO: Device Input and Output fgetwc(), fgetws(), fputwc(), fputws(), fwide(), fwprintf(), fwscanf(), getwc(), getwchar(), putwc(), putwchar(), ungetwc(), vfwprintf(), vfwscanf(), vwprintf(), vwscanf(), wprintf(), wscanf()</pre>
126971	<pre>XSI_C_LANG_SUPPORT: XSI General C Library</pre>
126972	a641(), daylight, drand48(), erand48(), ffs(), ffsl(), ffsll(), getdate(), hcreate(), hdestroy(),
126973	hsearch(), initstate(), insque(), jrand48(), l64a(), lcong48(), lfind(), lrand48(), lsearch(),
126974	memccpy(), mrand48(), nrand48(), random(), remque(), seed48(), setstate(), signgam,

126975 11. The *lchown()* function also depends on POSIX\_FILE\_ATTRIBUTES.

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